

Quantum CCA-Secure PKE, Revisited

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SandboxAQ



: <https://varun-maram.github.io/>



: varun-maram-pqc

Joint work with Navid Alapati

ETH zürich



SANDBOXAQ

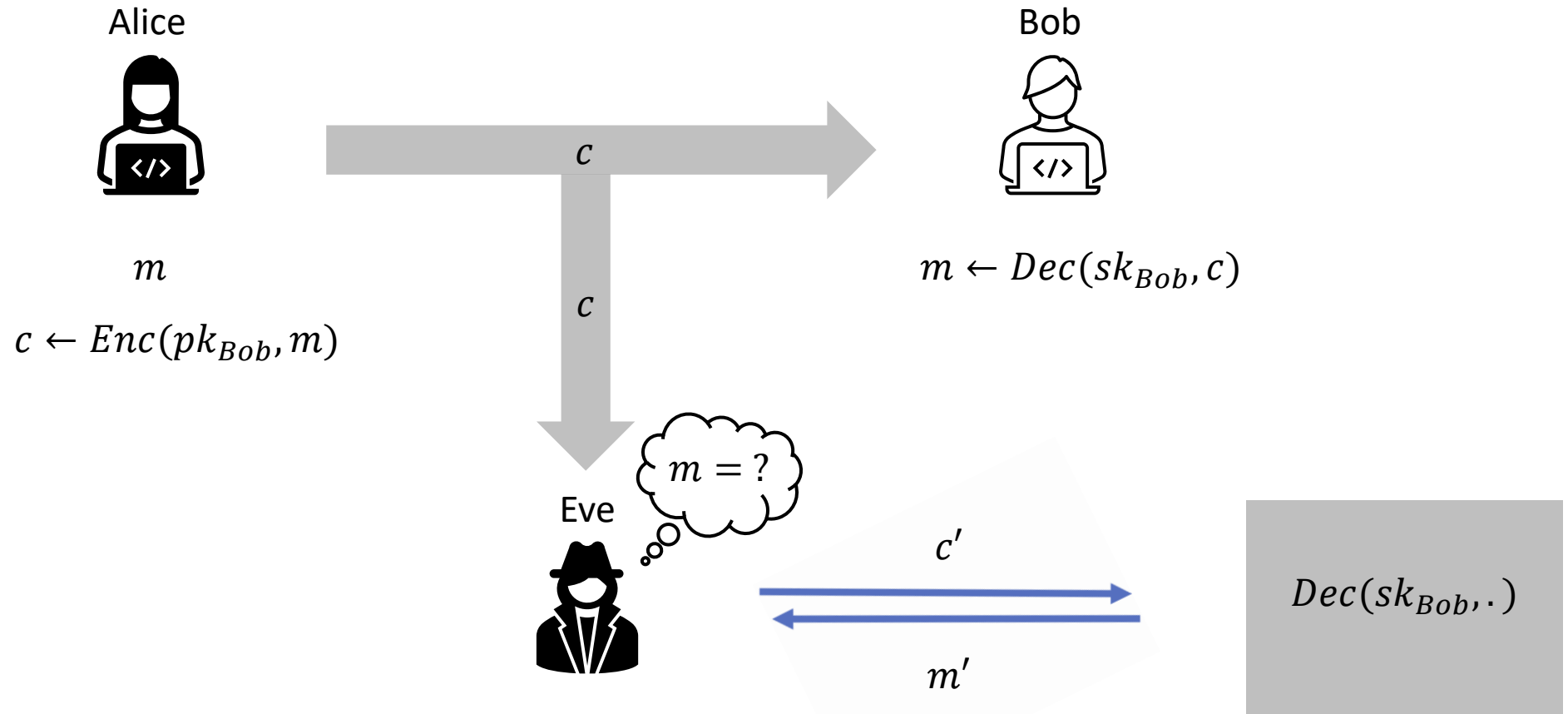
VISA

Quantum CCA-Secure PKE, Revisited

Quantum CCA-Secure PKE, Revisited

IND-CCA Security

$$PKE = (KGen, Enc, Dec)$$

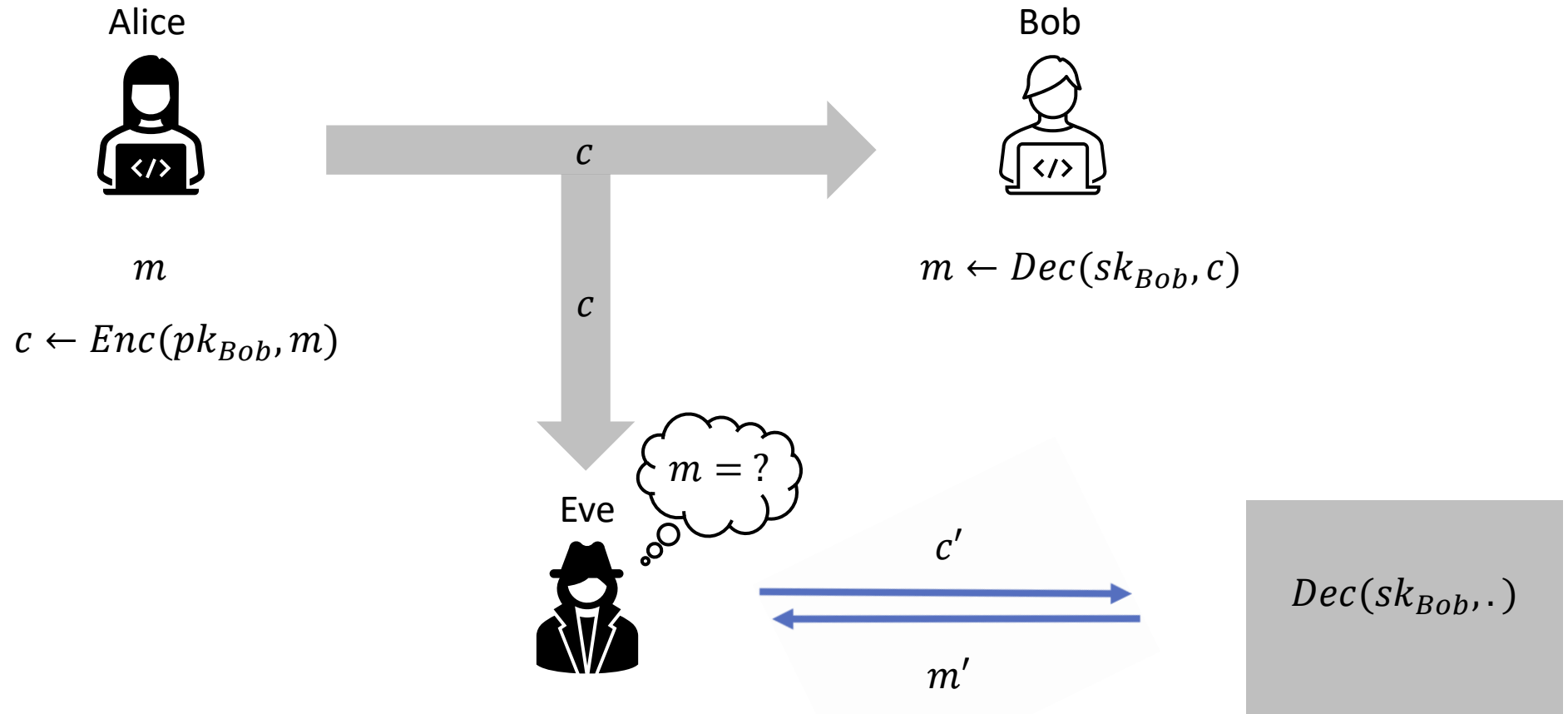


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IND-CCA Security

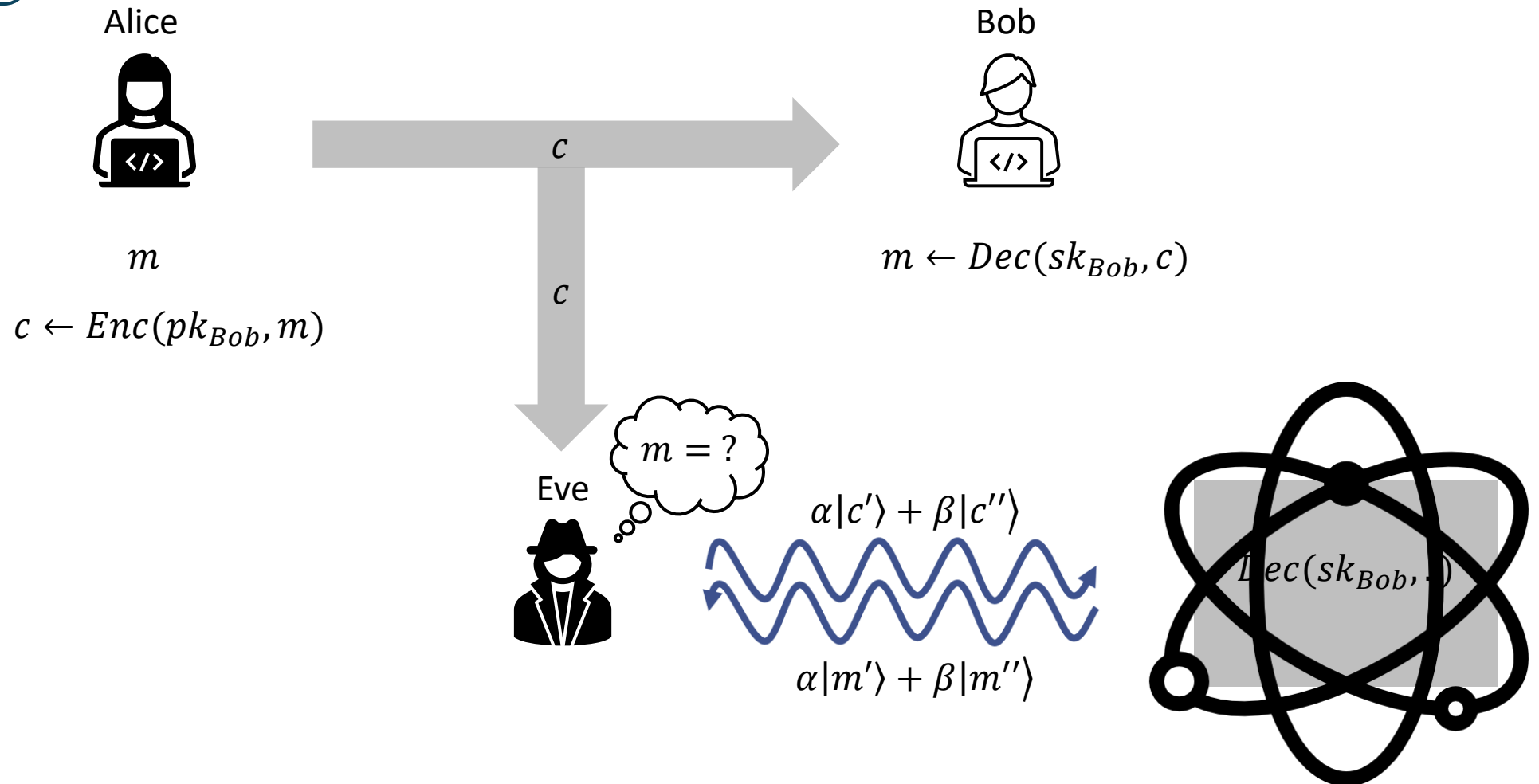
$$PKE = (KGen, Enc, Dec)$$



IND-**q**CCA Security

Introduced by [Boneh-Zhandry'13].

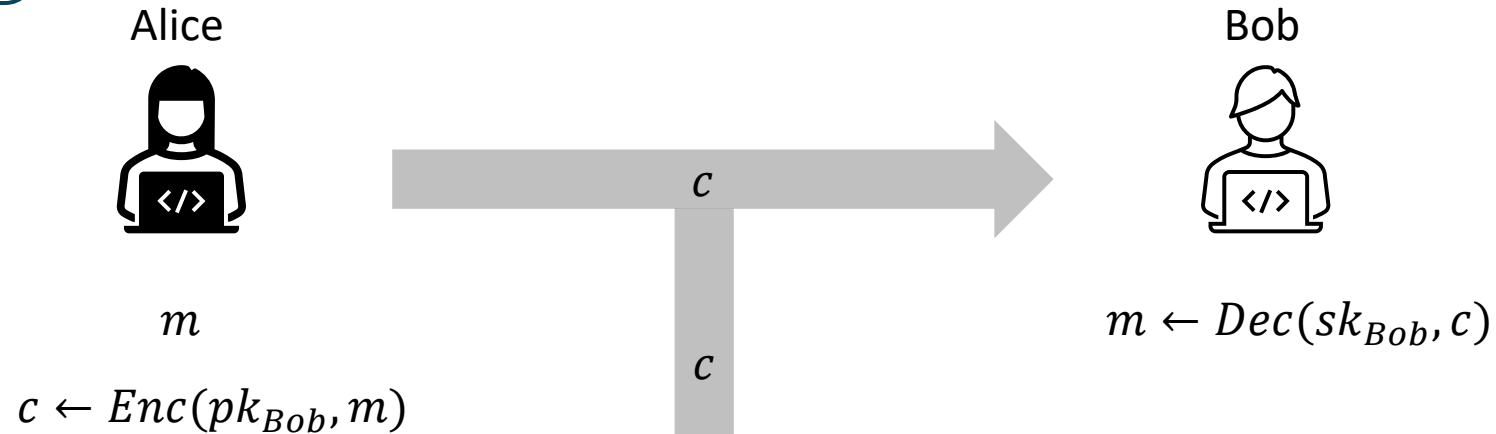
$$PKE = (KGen, Enc, Dec)$$



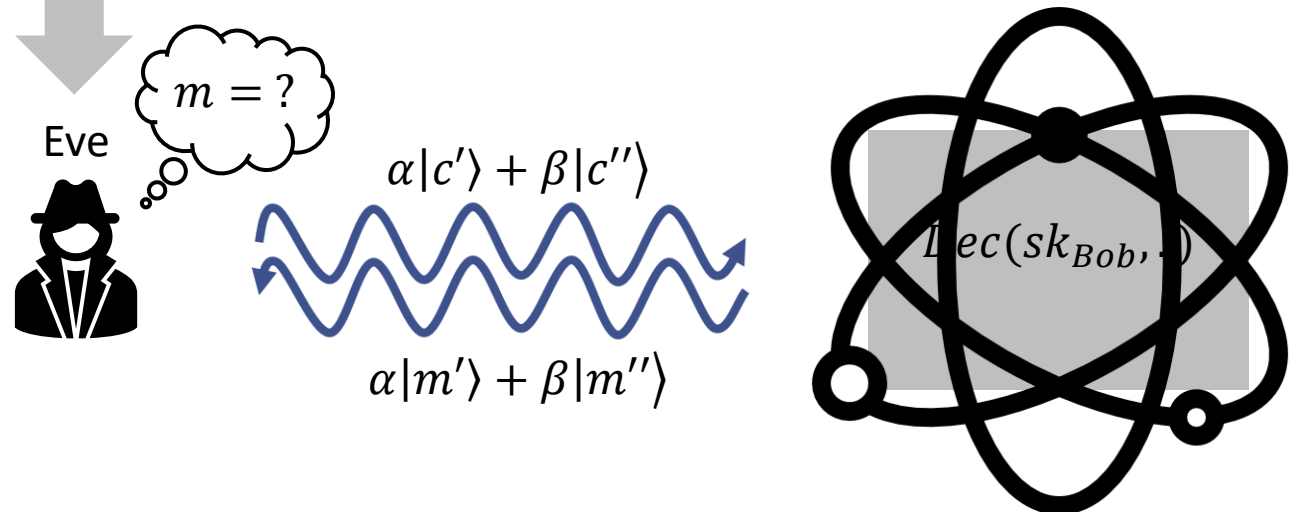
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$$PKE = (KGen, Enc, Dec)$$



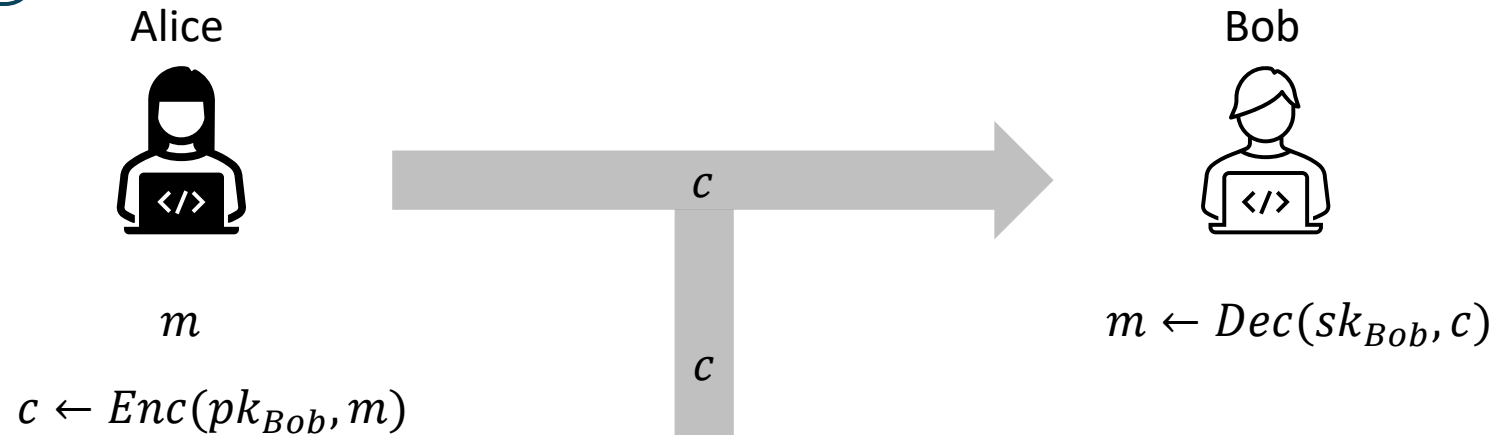
- Relevance in future when quantum computing becomes ubiquitous.



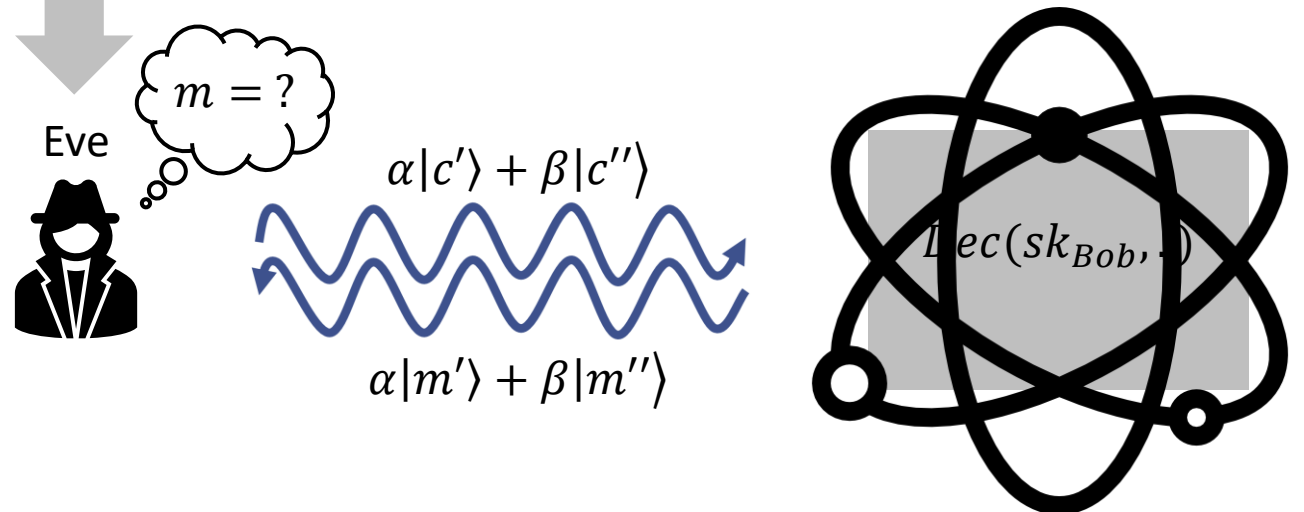
IND-**q**CCA Security

Introduced by [Boneh-Zhandry'13].

$$PKE = (KGen, Enc, Dec)$$



- Relevance in future when quantum computing becomes ubiquitous.
- Also, in not-so-far future when adversaries can trick classical devices to behave “quantumly” (e.g., “frozen smart-card attacks”).



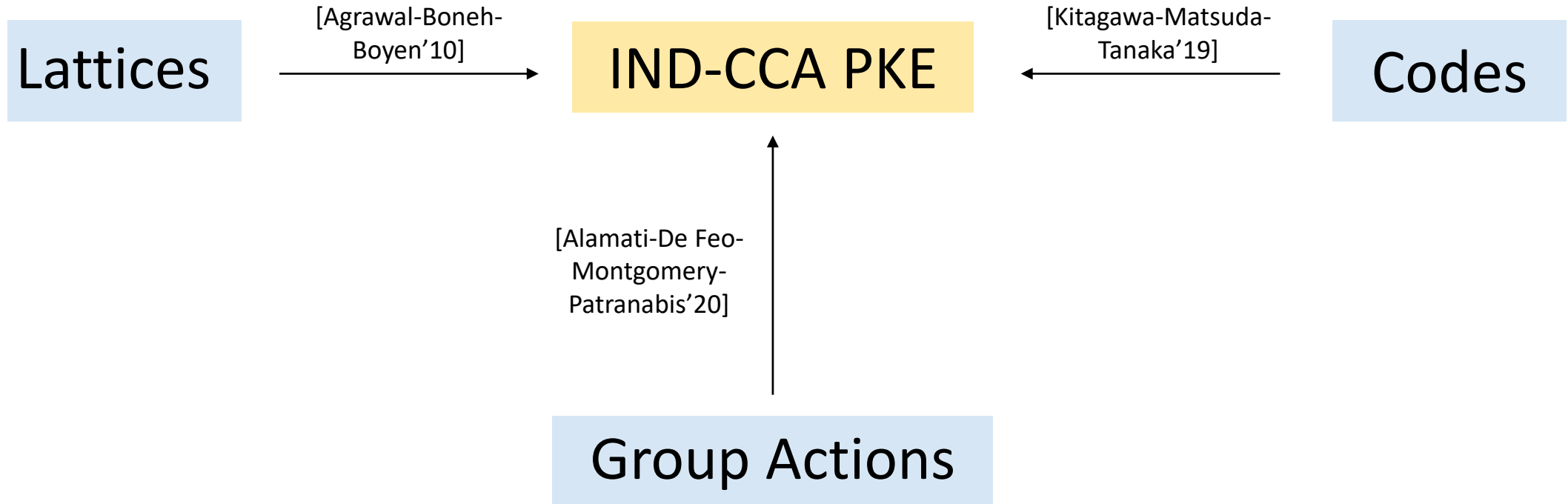
Quantum CCA-Secure PKE, Revisited

Quantum CCA-Secure PKE, Revisited

Motivation

IND-CCA PKE

Motivation



Motivation

IND-**q**CCA PKE

Motivation

Lattices

[Boneh-Zhandry'13]

IND-**q**CCA PKE

Only prior construction in
the **standard** model.*

Motivation

Lattices

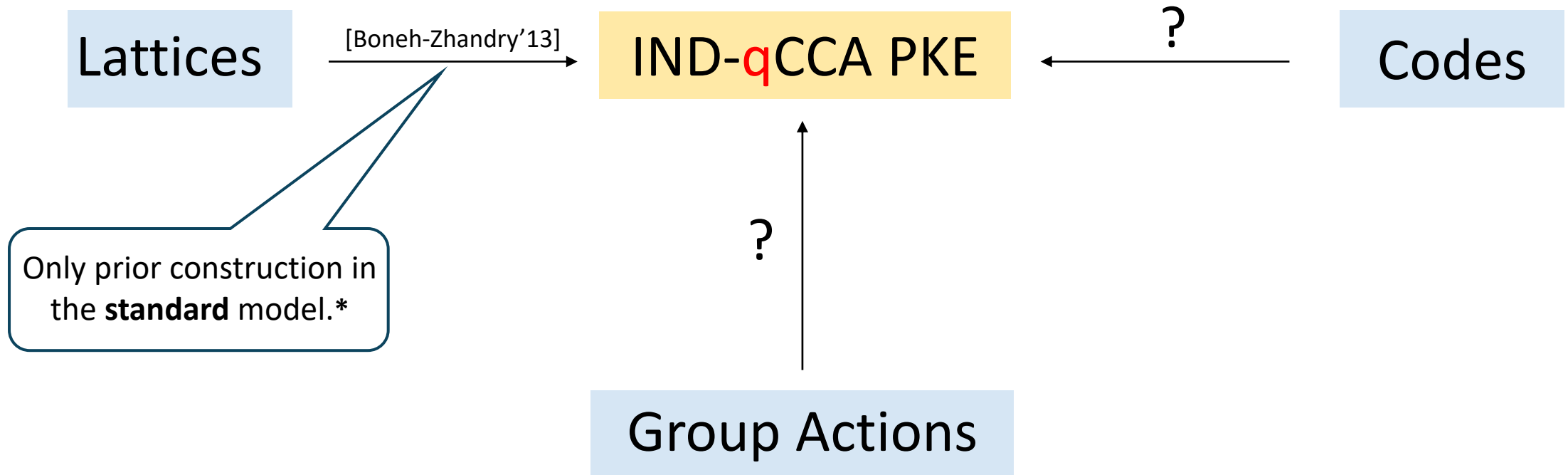
[Boneh-Zhandry'13]

IND-**q**CCA PKE

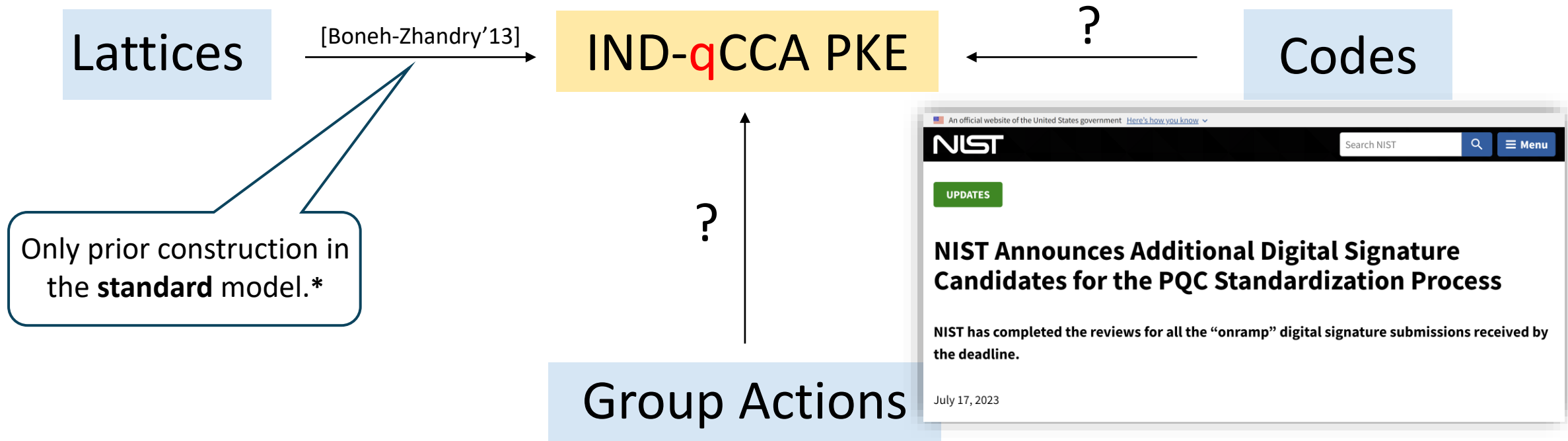
Only prior construction in
the **standard** model.*

*There exist richer constructions
of qCCA-secure PKE in the
idealized **quantum ROM** – e.g.,
by [Xagawa-Yamakawa'19].

Motivation



Motivation



Motivation

Quantum Algorithms for Lattice Problems

Yilei Chen*

April 10, 2024

Lattices

[Boneh-Zhandry'13]

IND-**q**CCA PKE

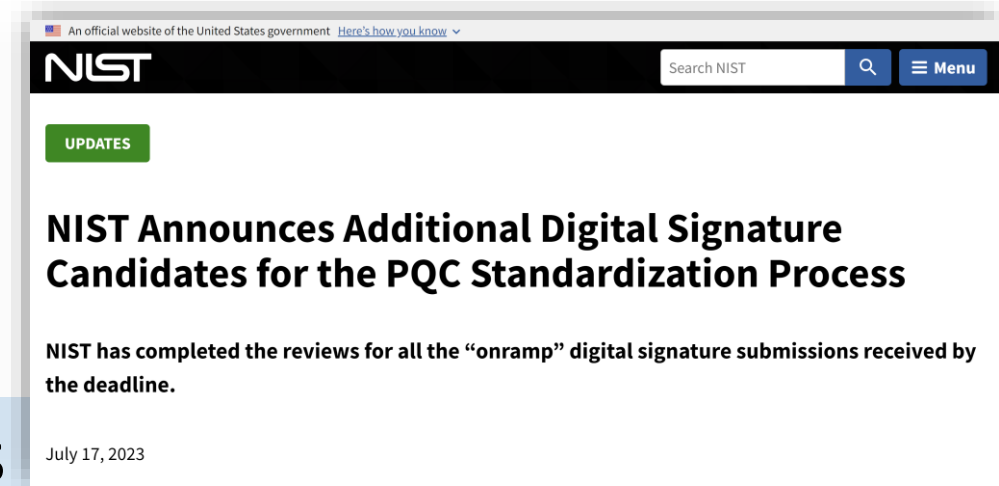
?

Codes

Only prior construction in
the **standard** model.*

?

Group Actions



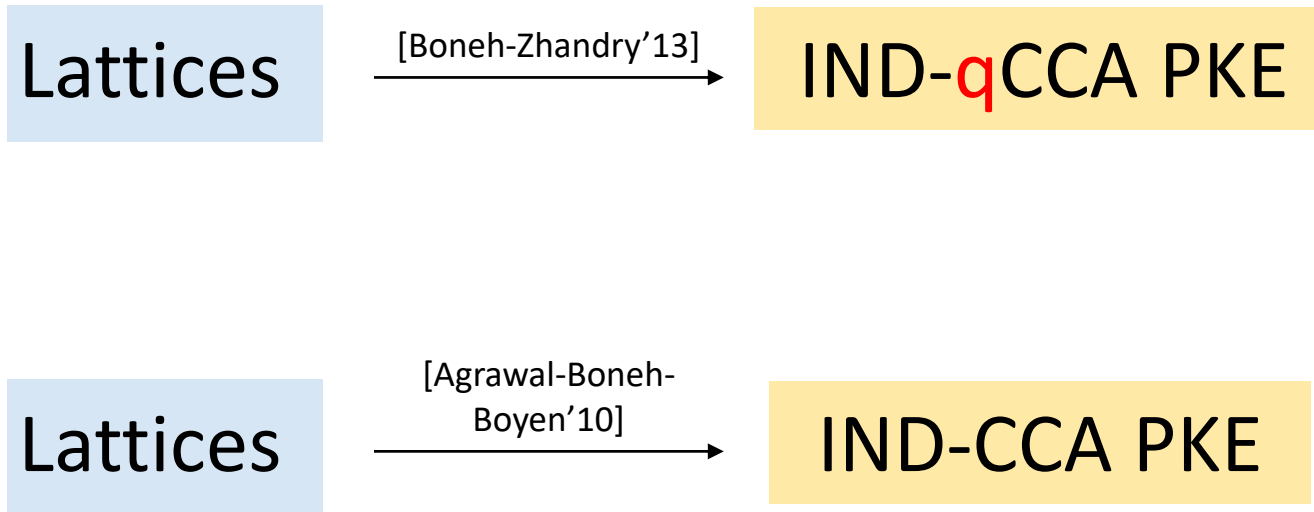
Motivation

Lattices

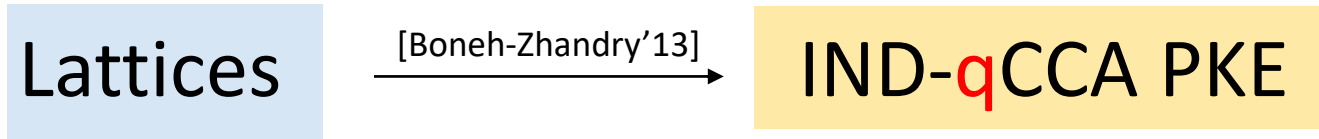
[Boneh-Zhandry'13]

IND-**q**CCA PKE

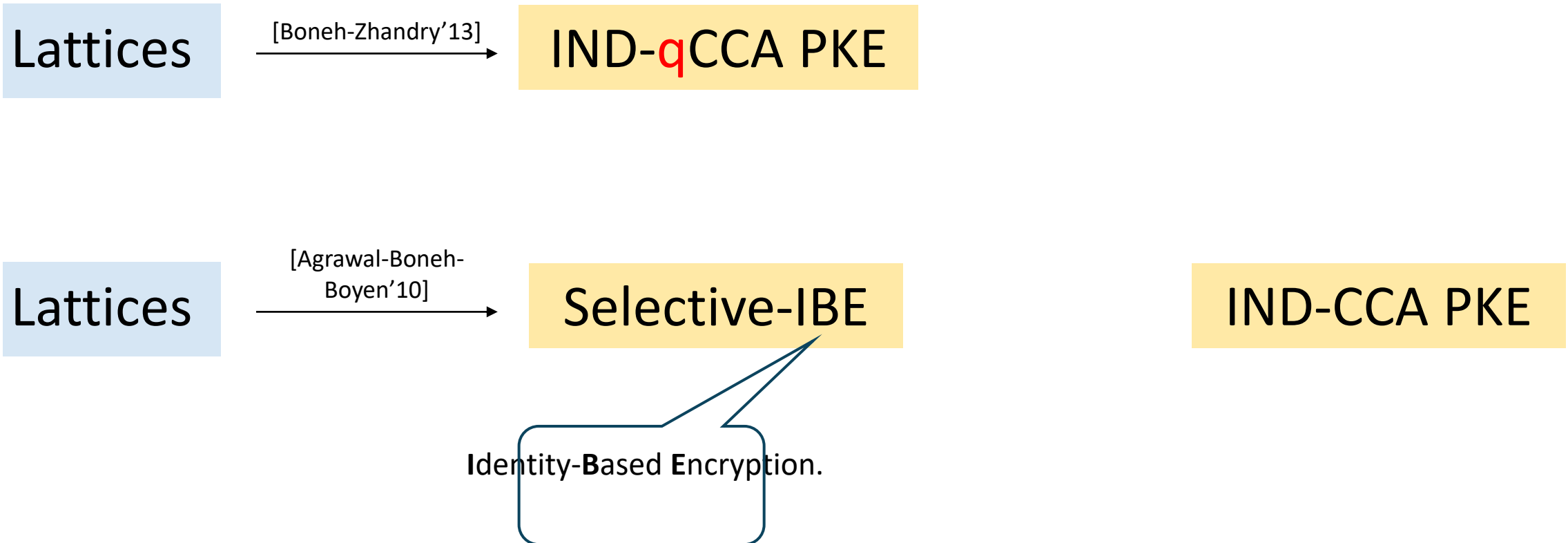
Motivation



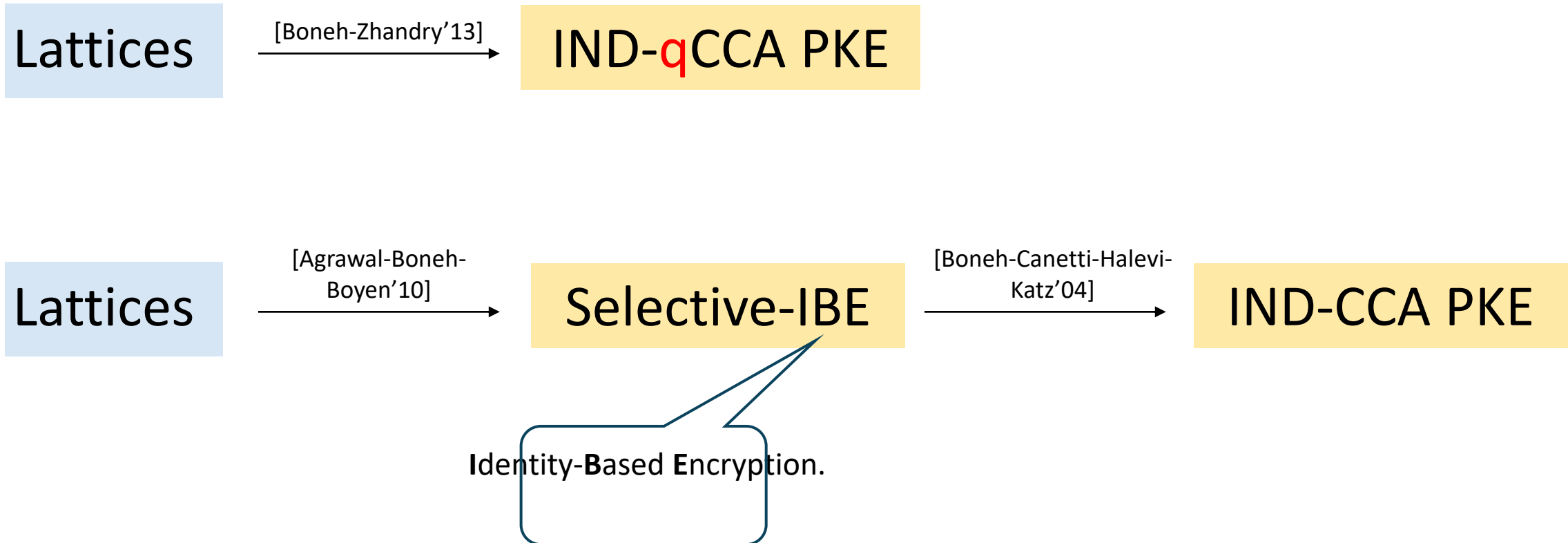
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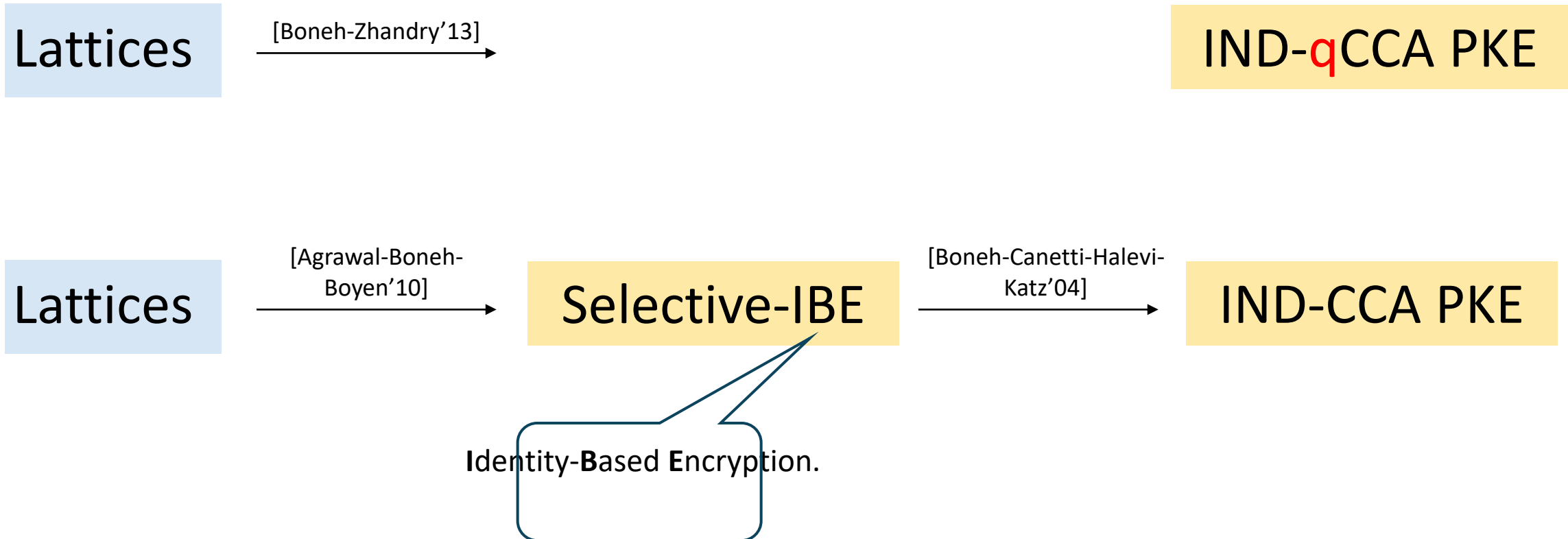
Motivation



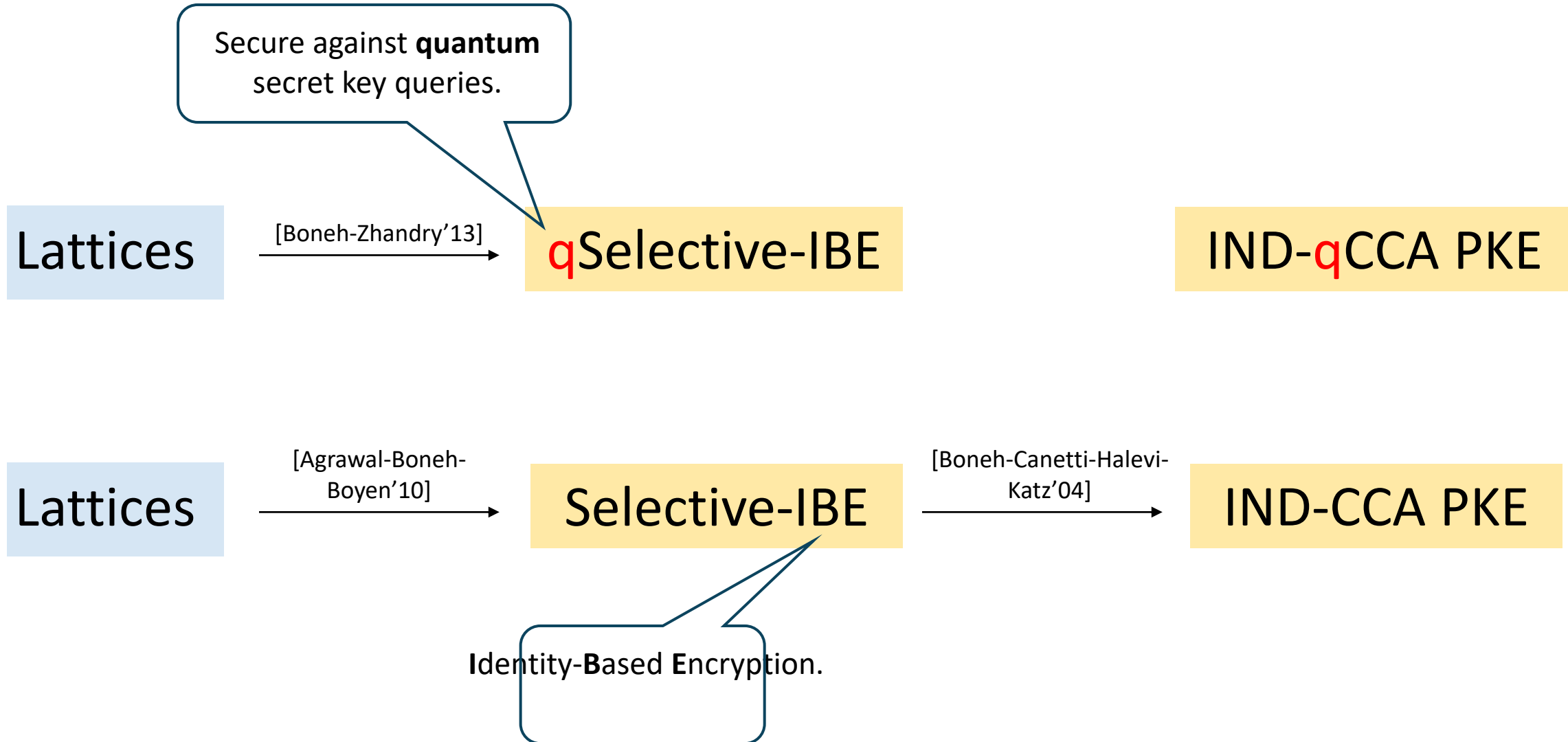
Motivation



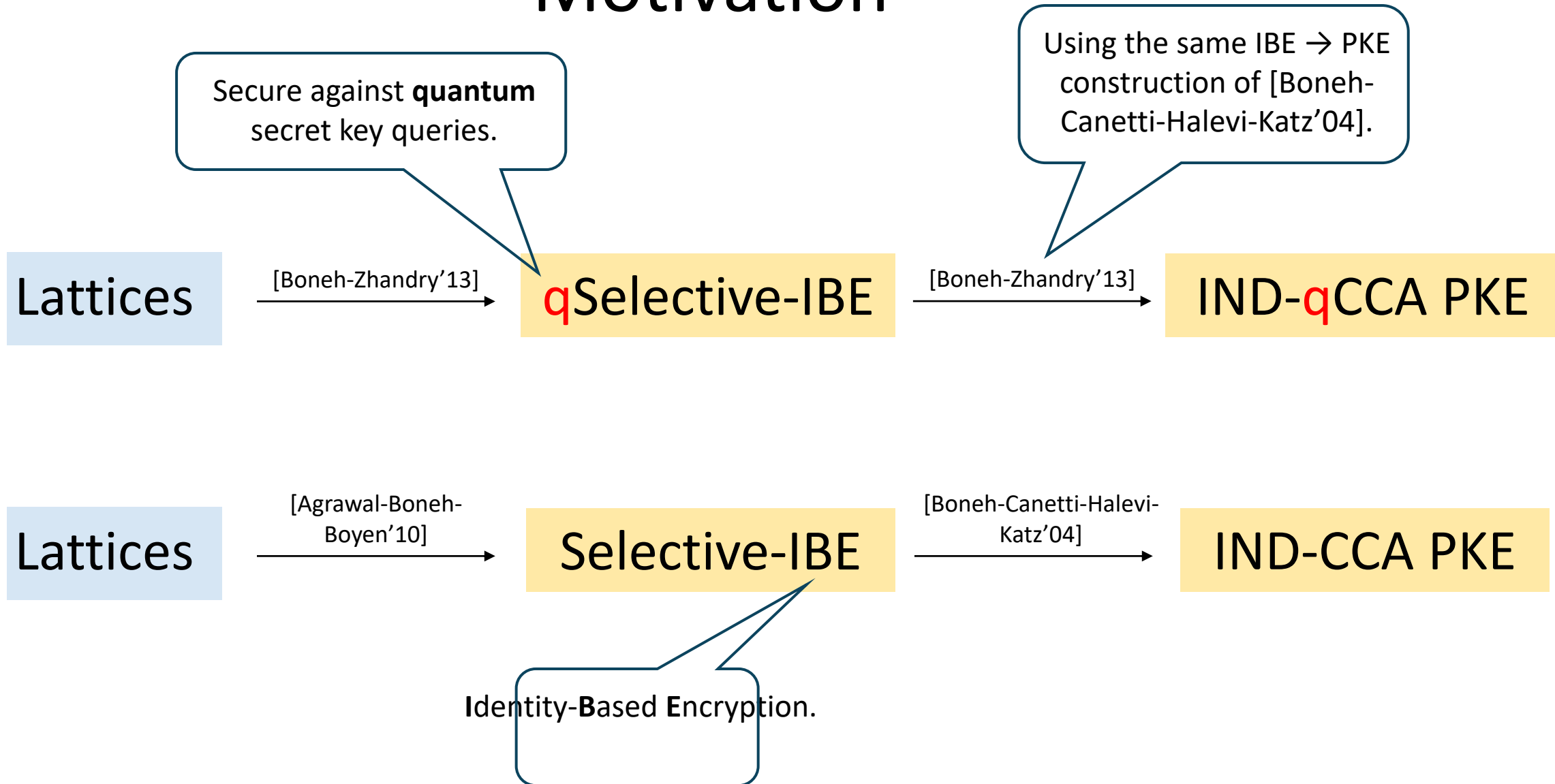
Motivation



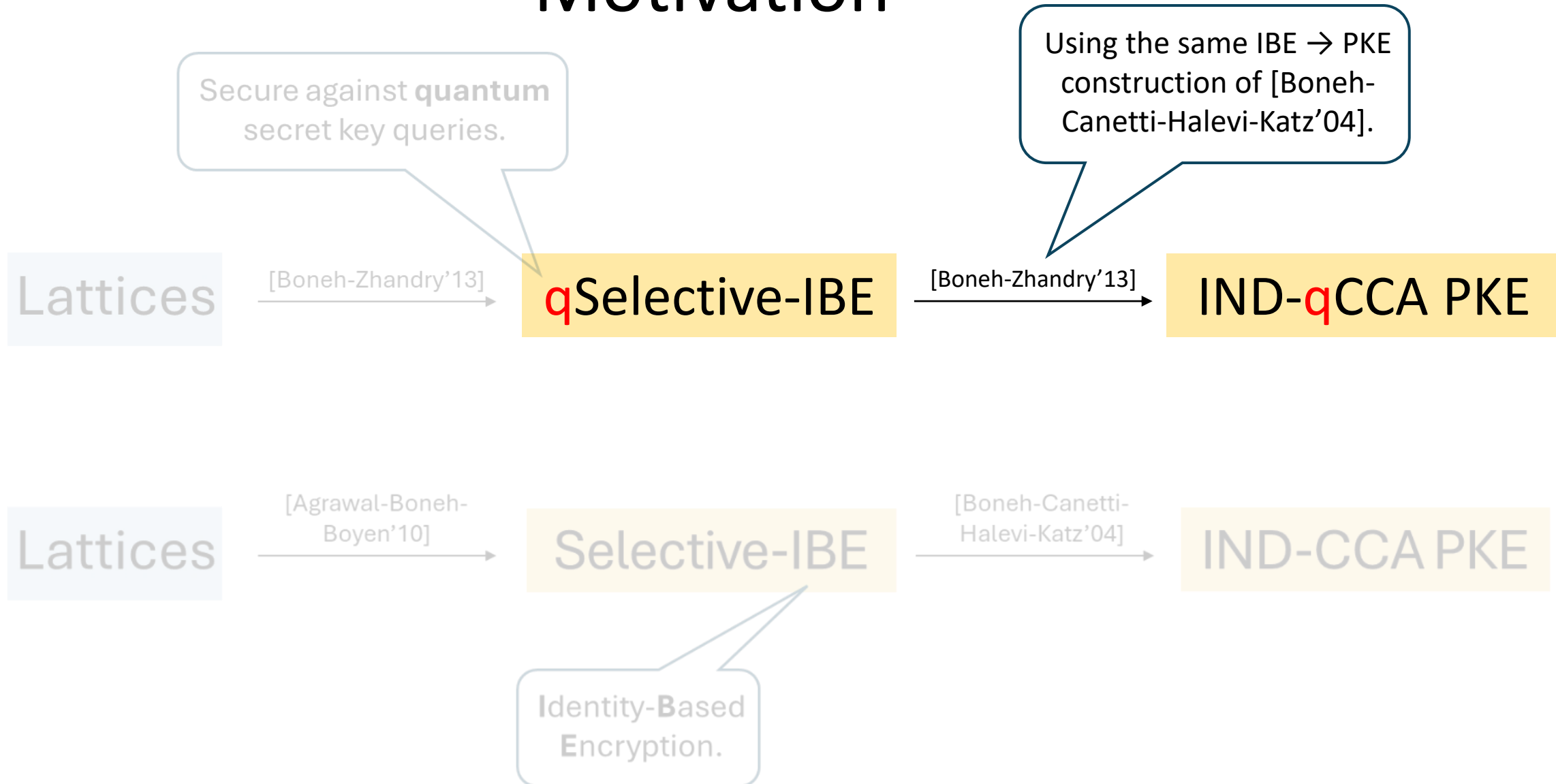
Motivation



Motivation



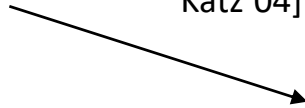
Motivation



Motivation

Selective-IBE

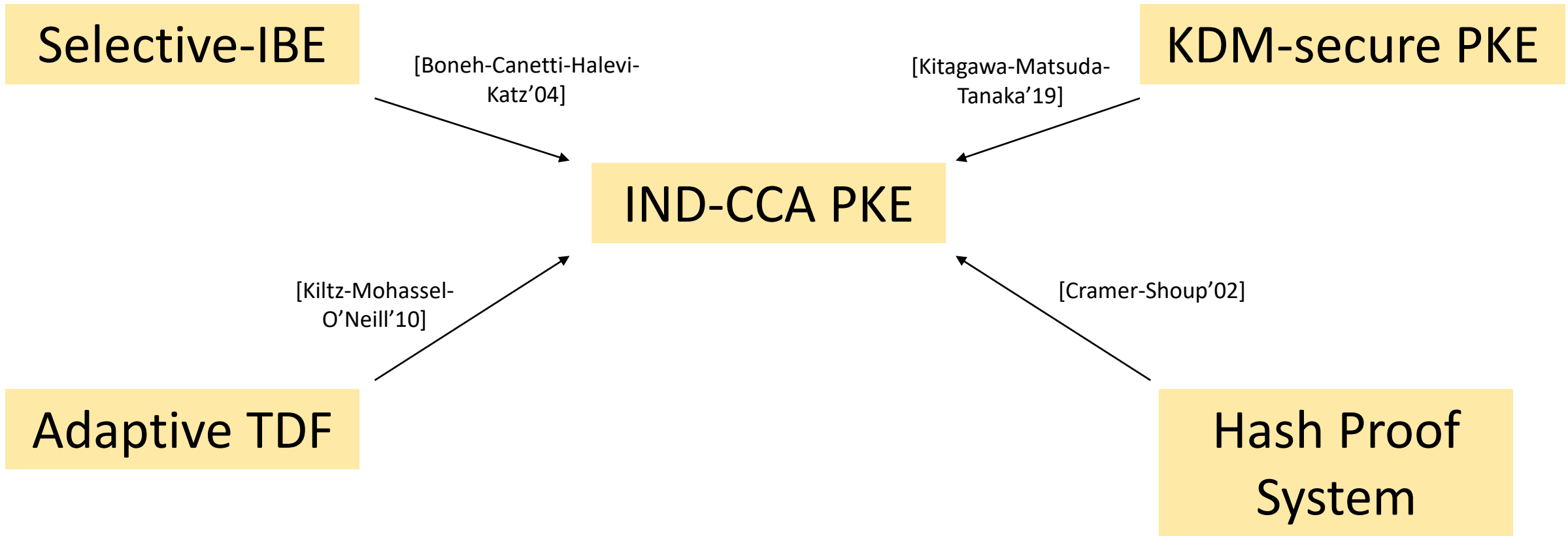
[Boneh-Canetti-Halevi-
Katz'04]



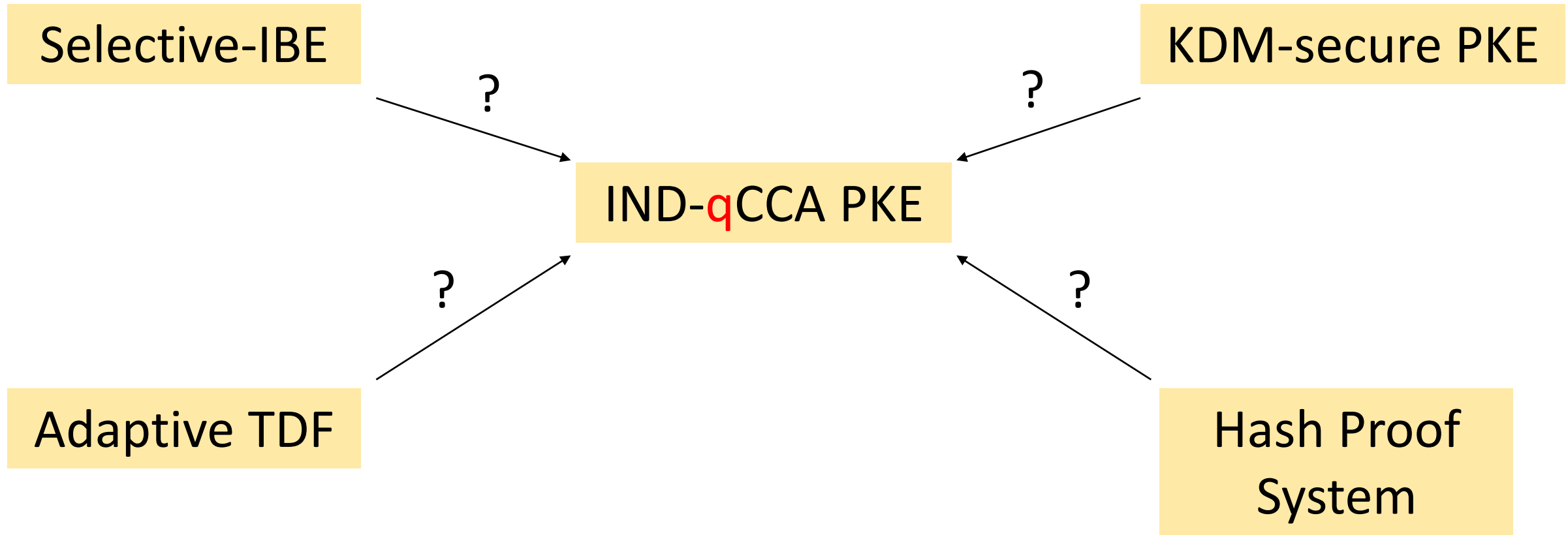
IND-CCA PKE

Motivation

Key-Dependent Message.

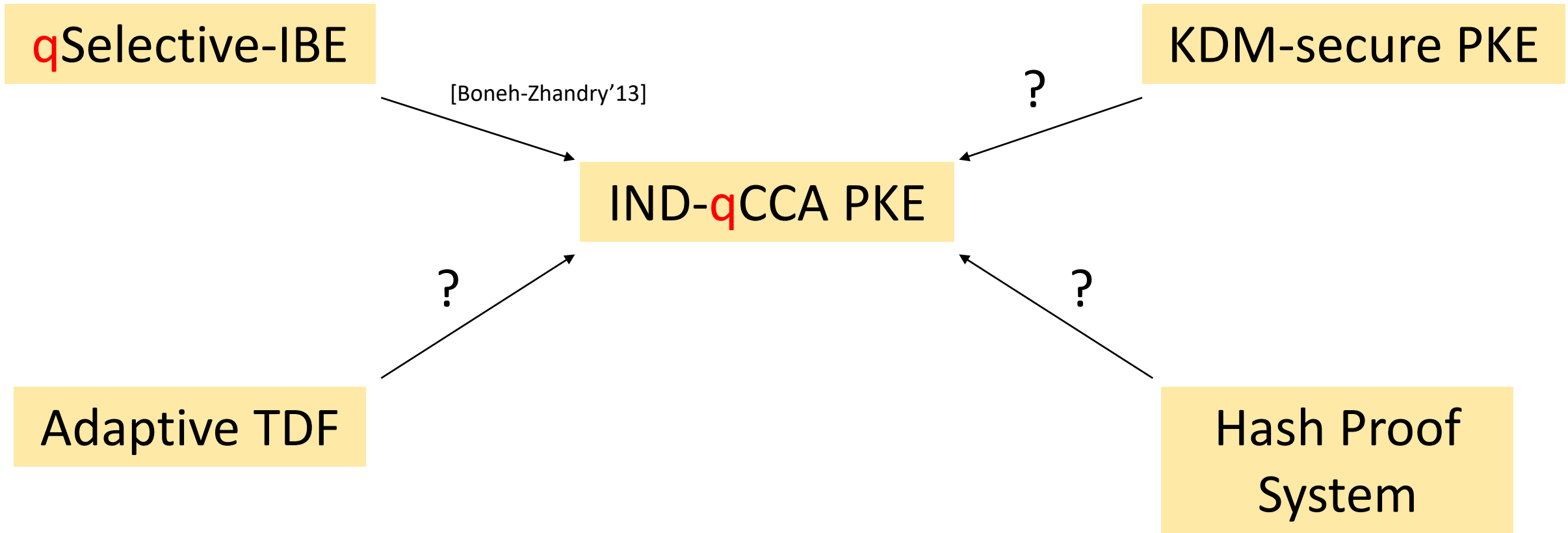


Motivation



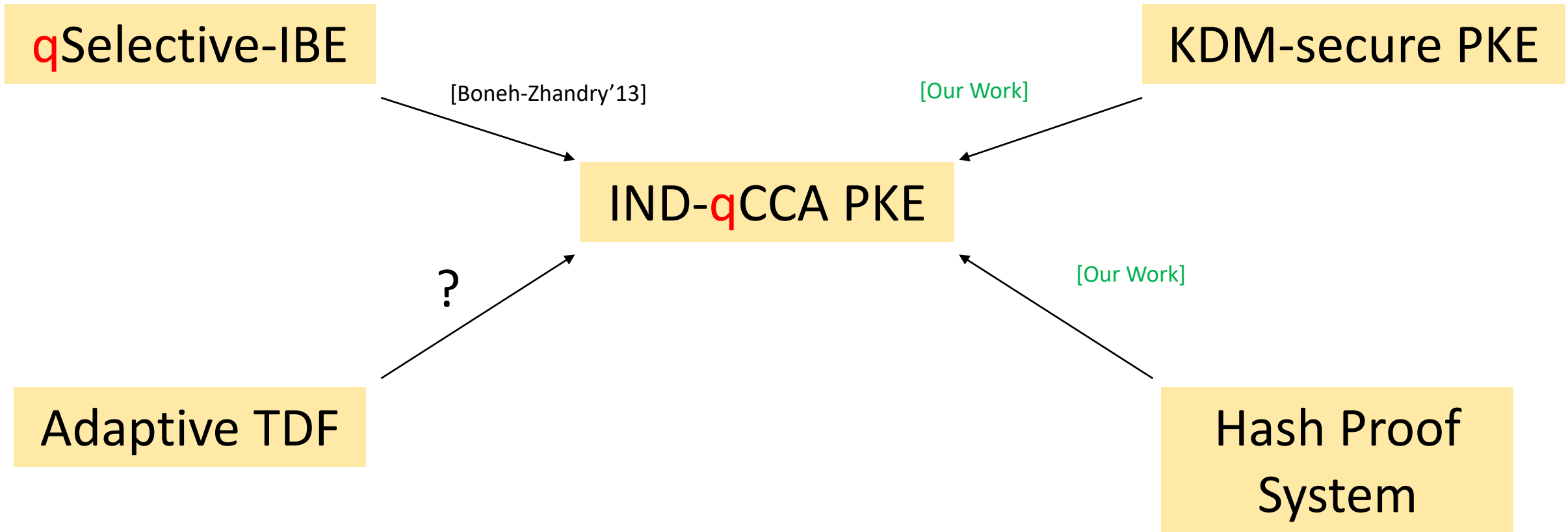
Secure against **post-quantum**
adversaries with **quantum**
access to secret oracles.

Motivation



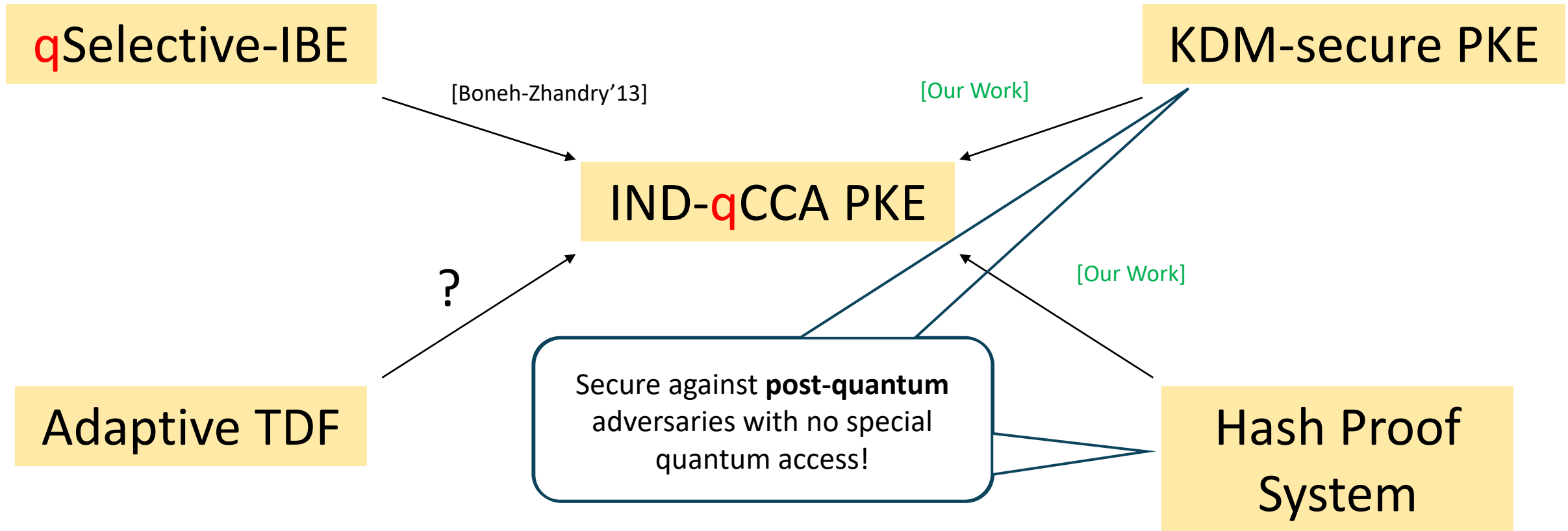
Secure against **post-quantum**
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Overview: Results

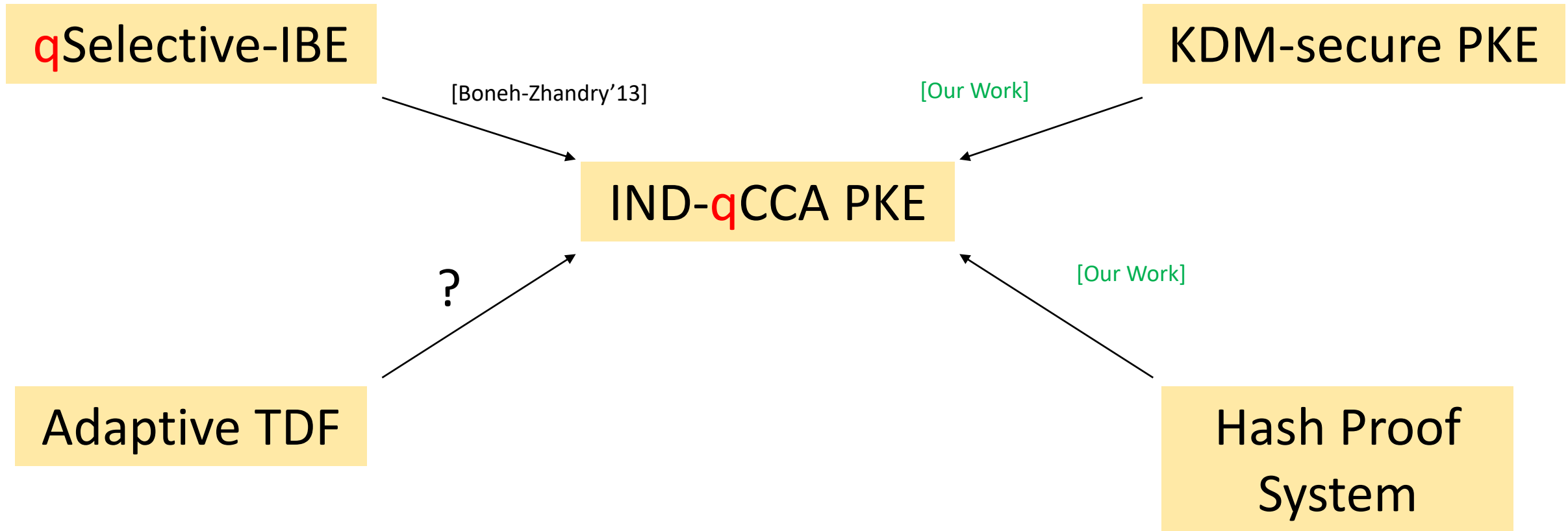


Secure against **post-quantum** adversaries with **quantum access** to secret oracles.

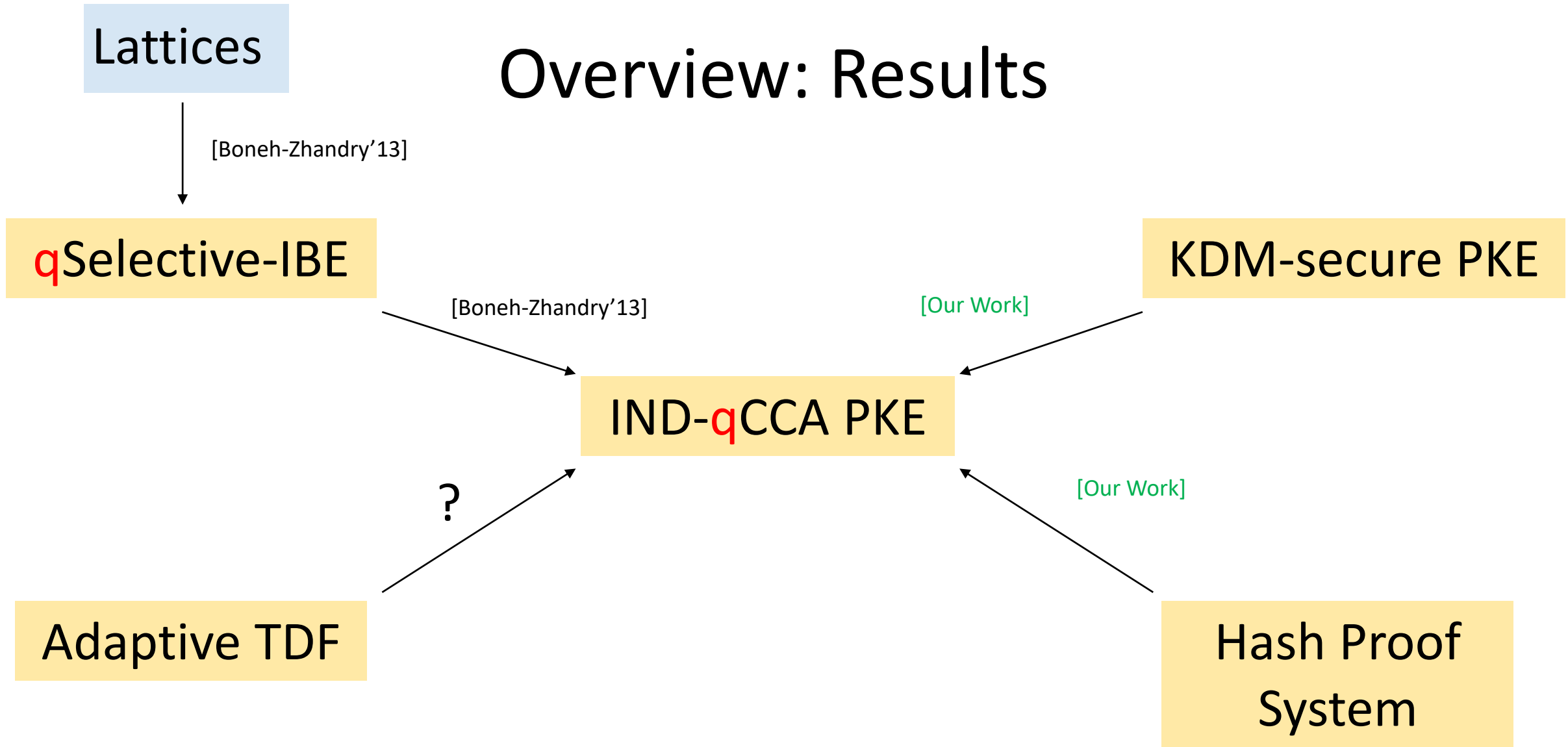
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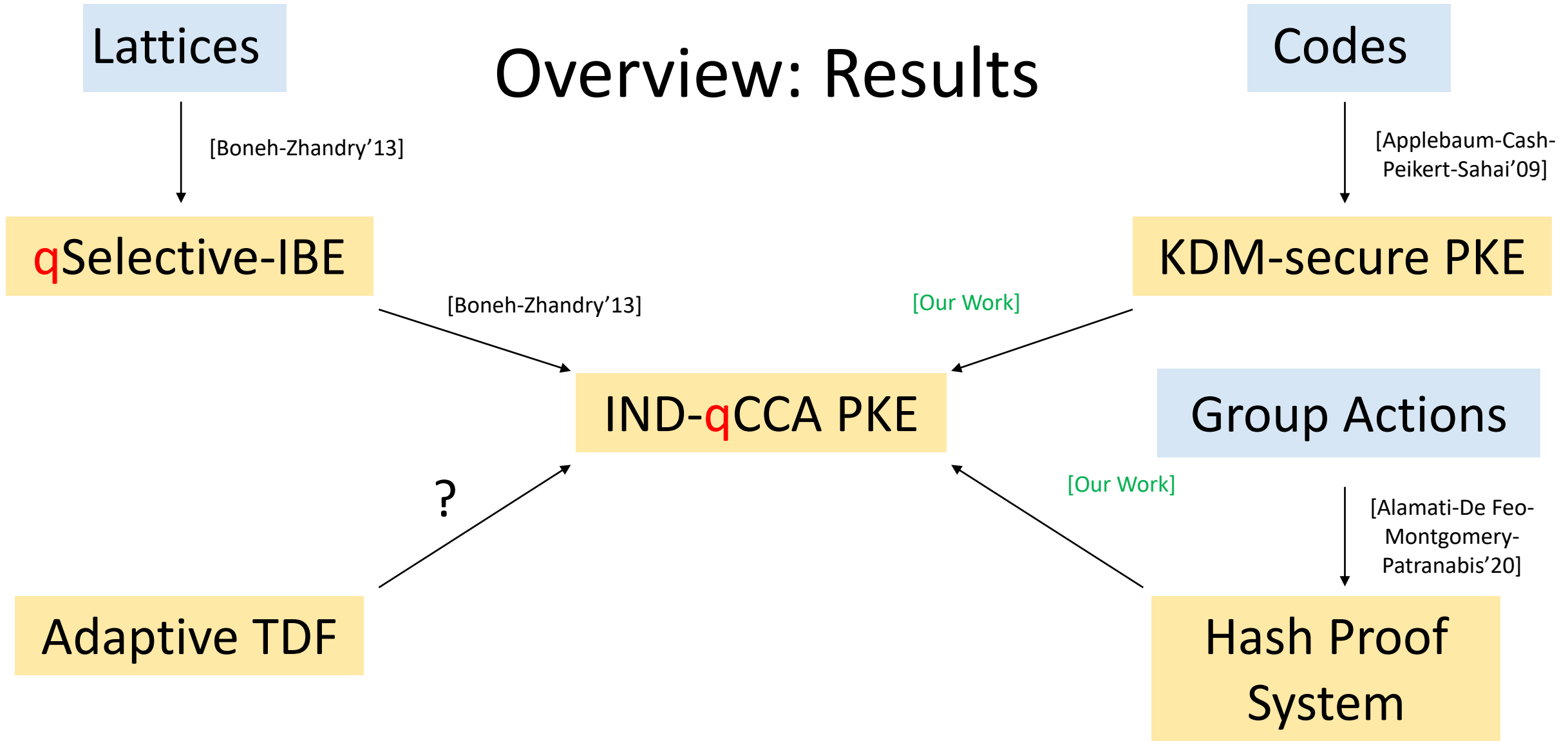
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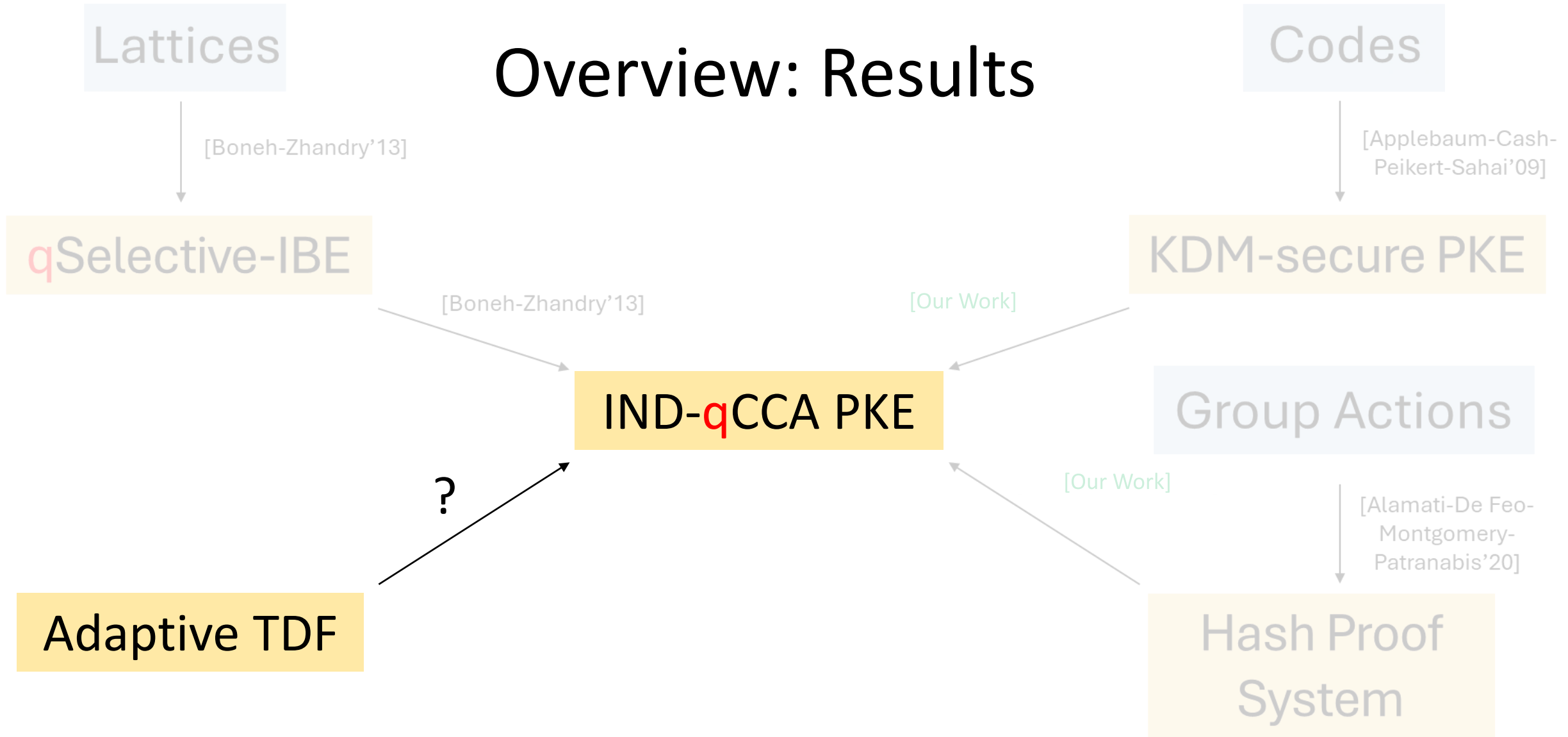
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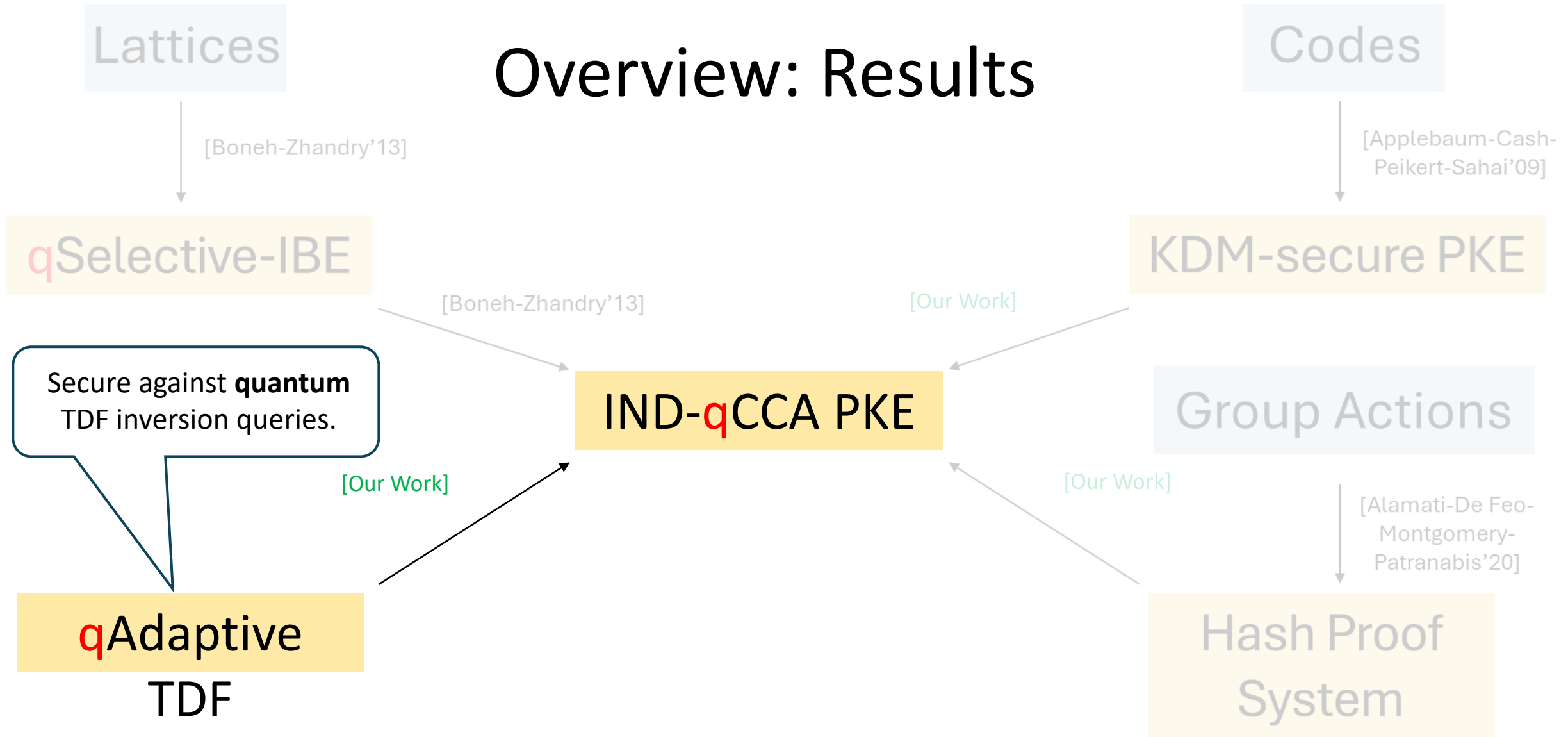
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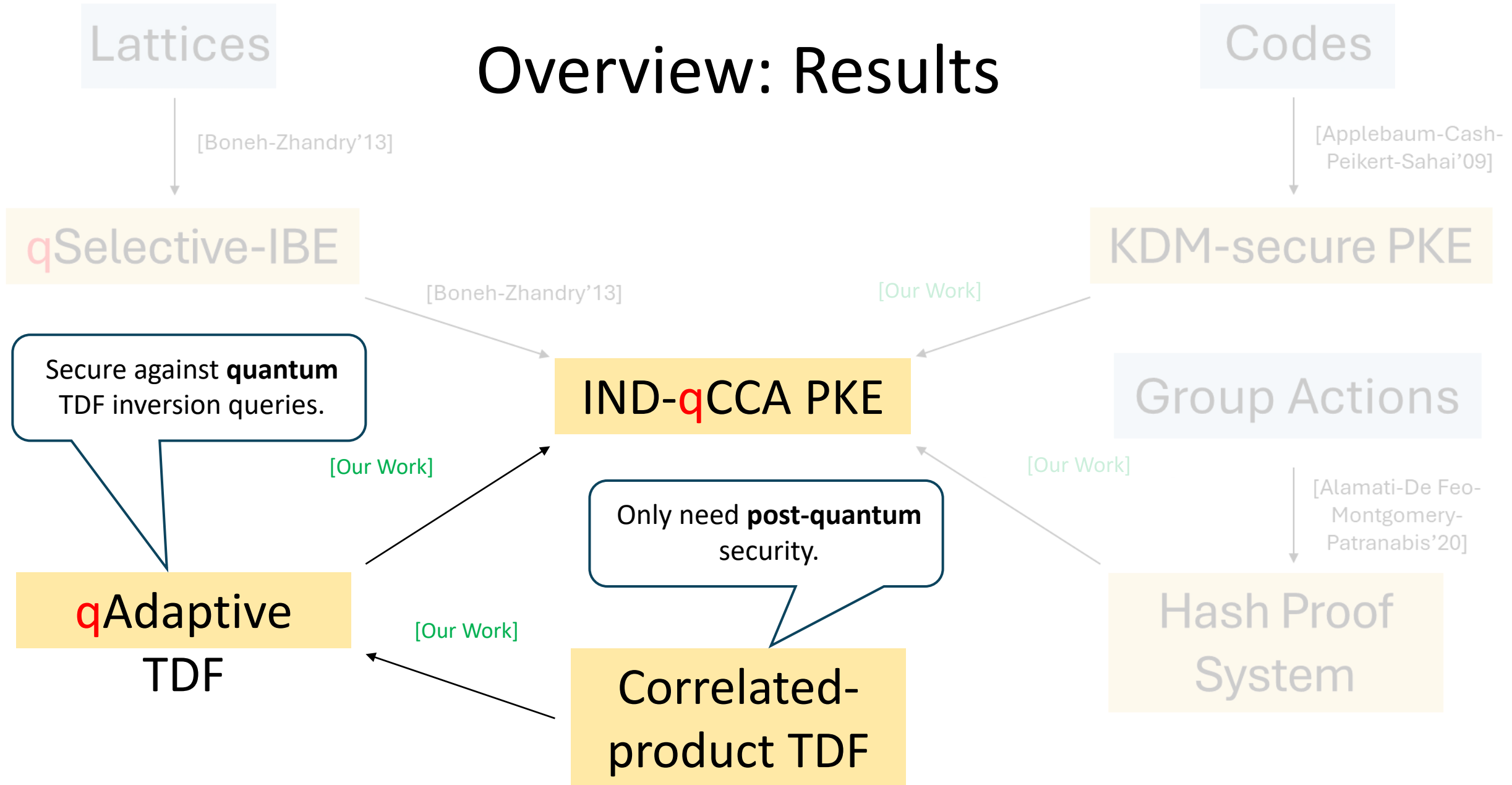
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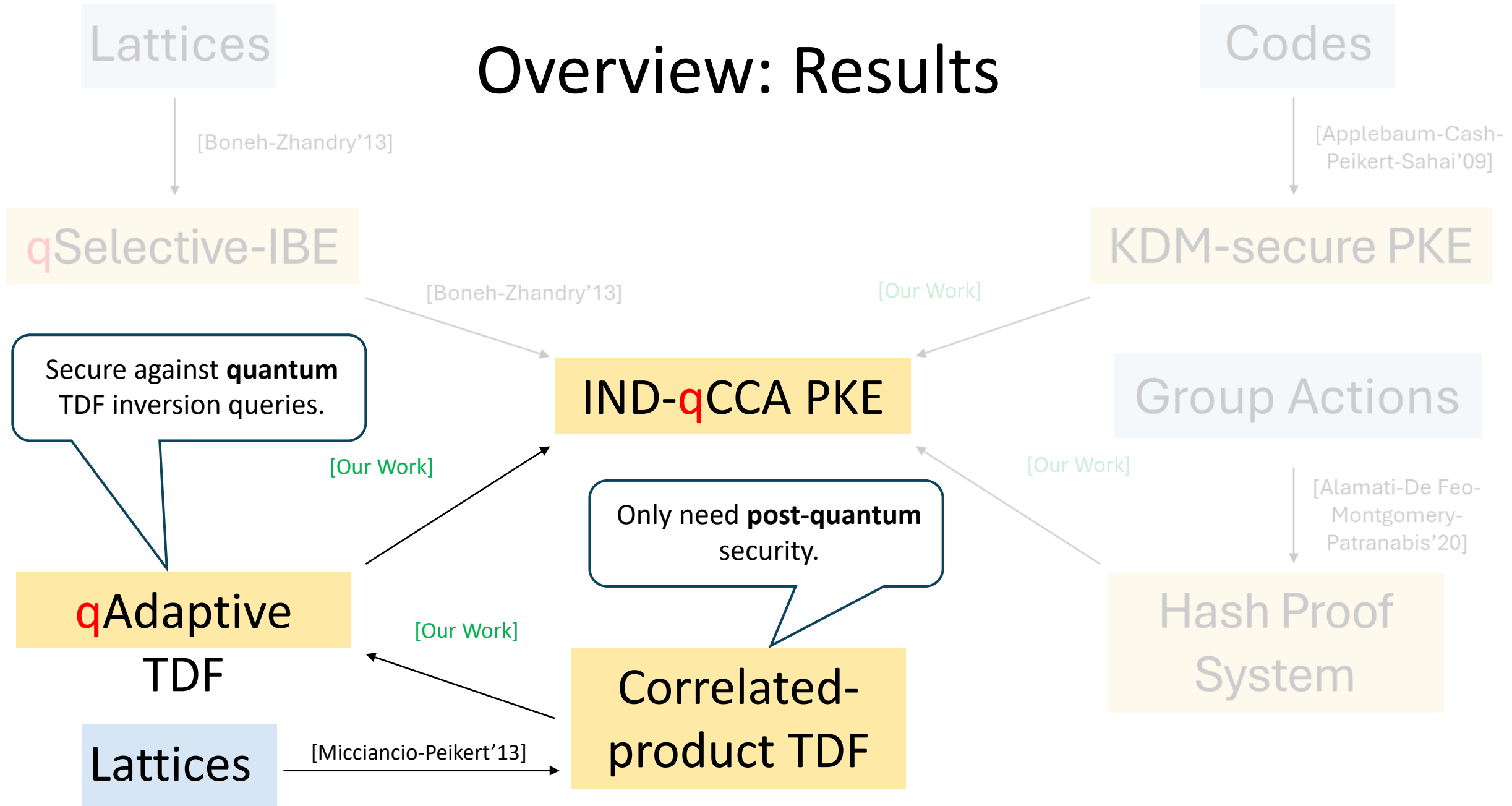
Overview: Results



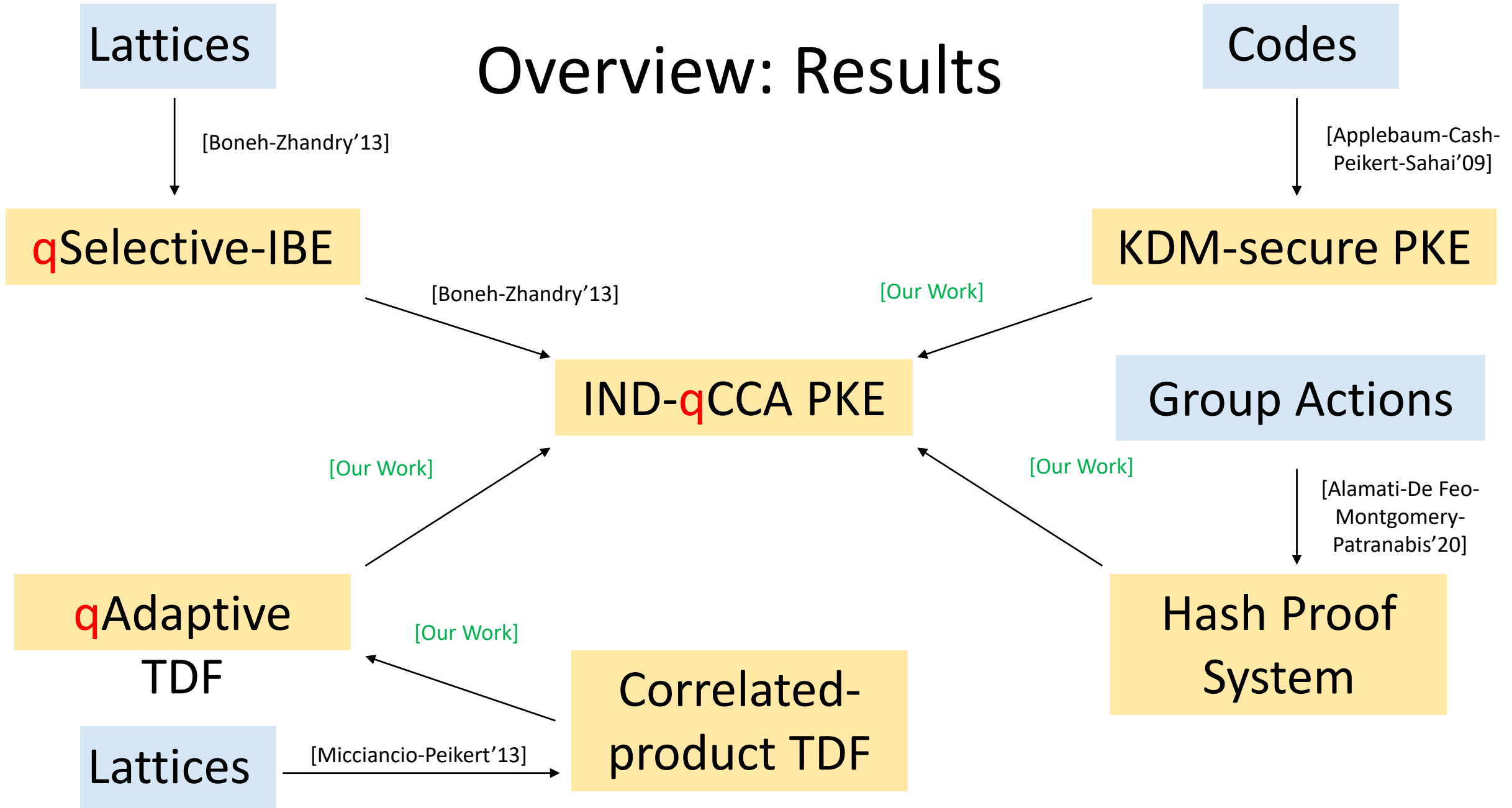
Overview: Results



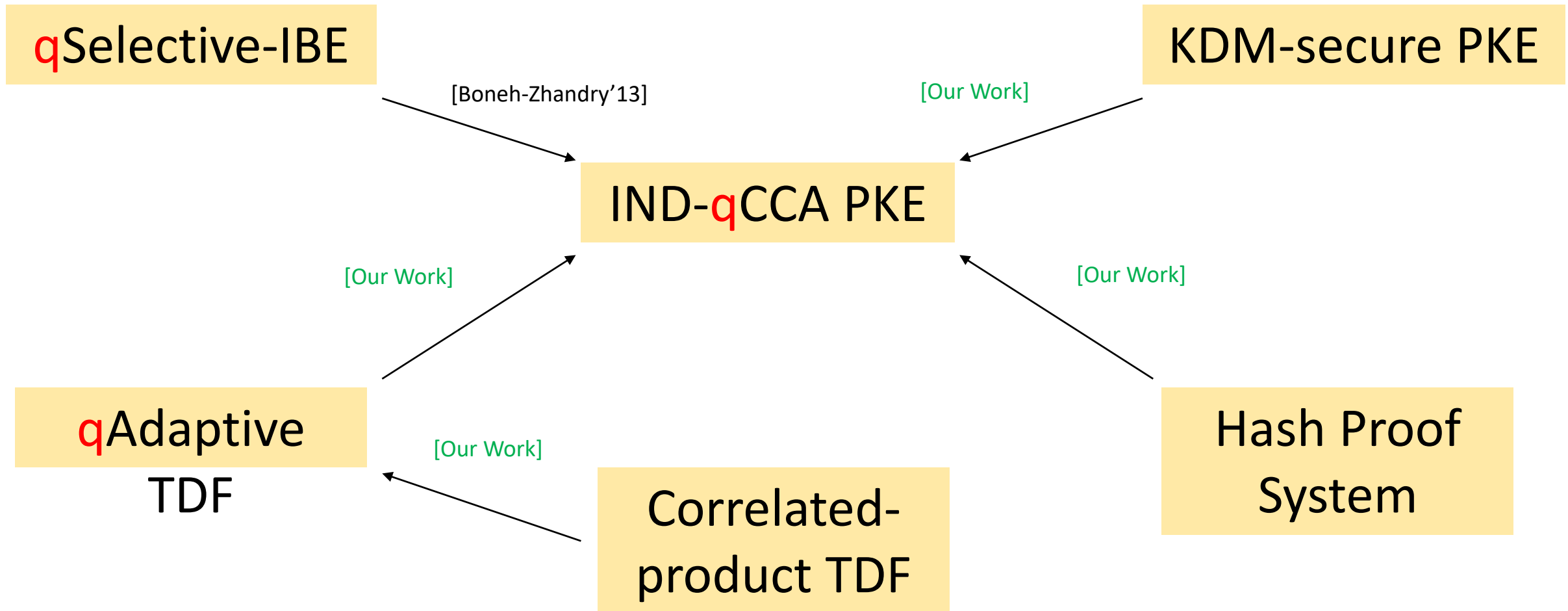
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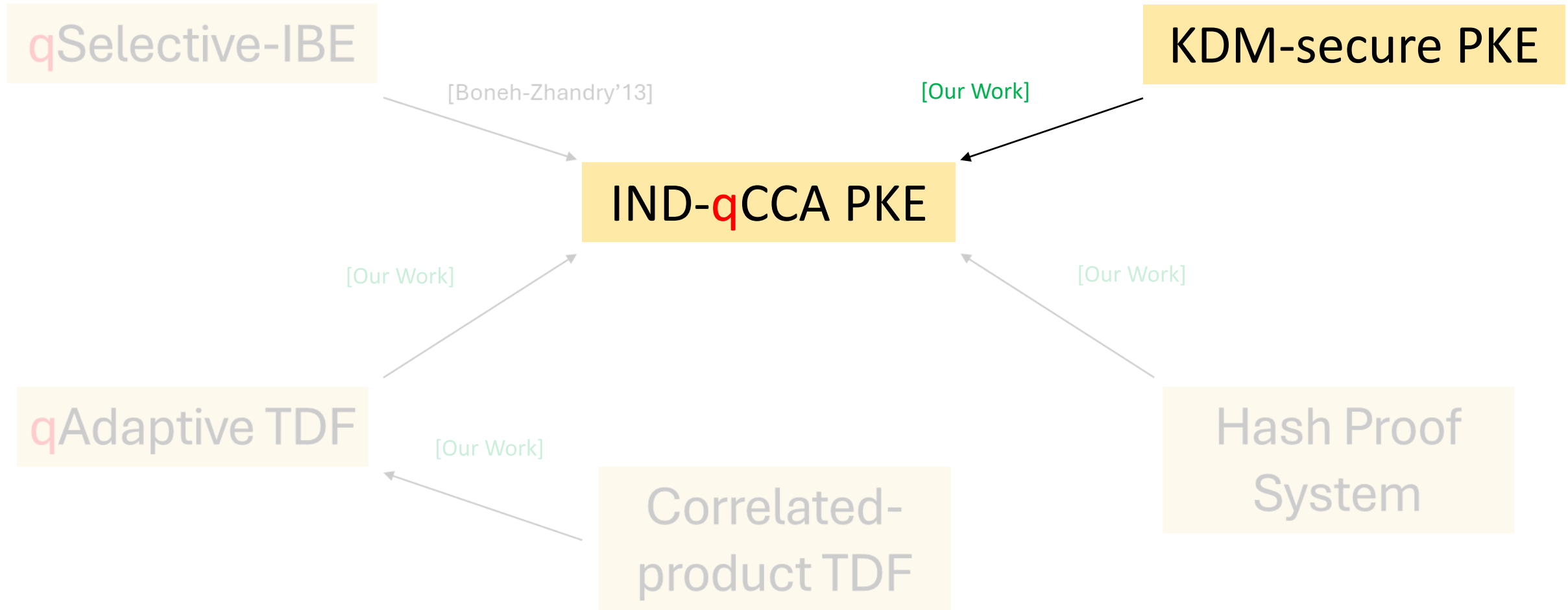
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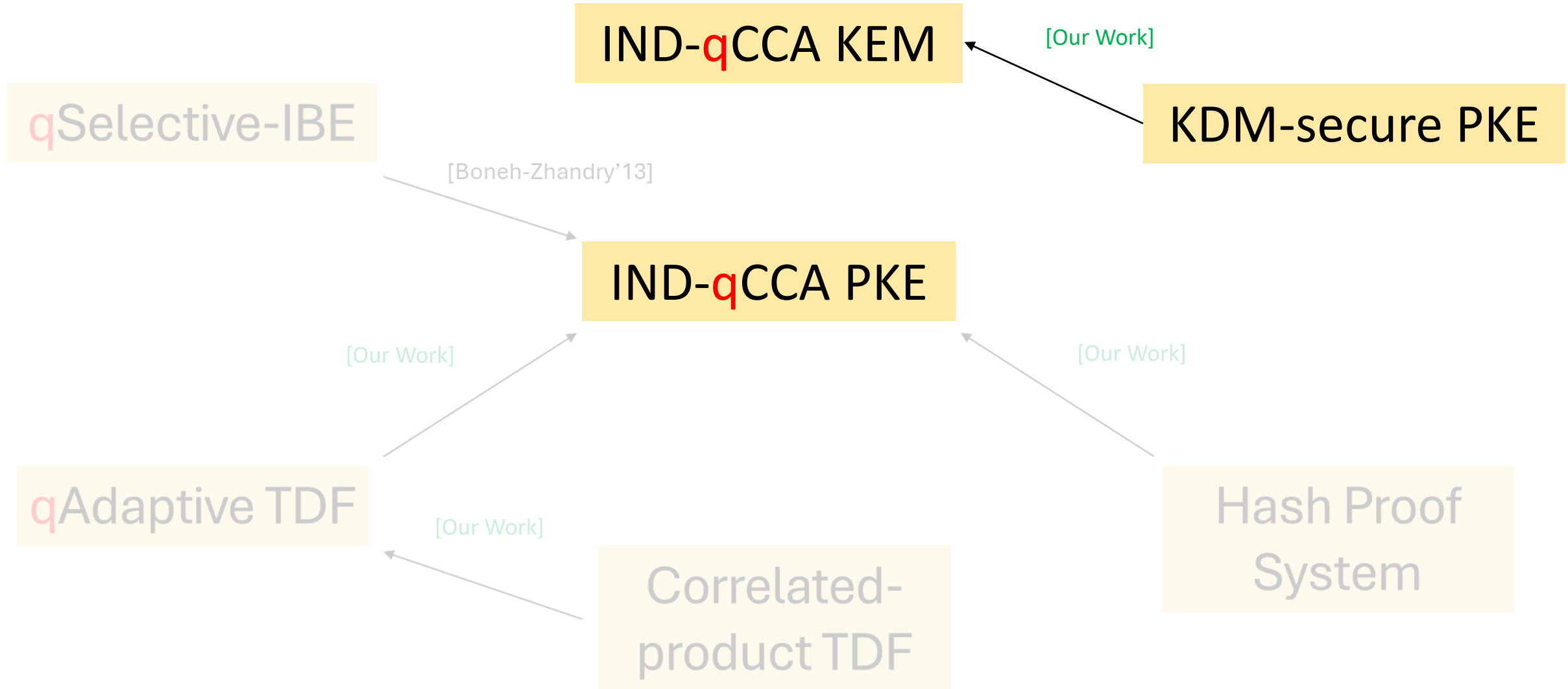
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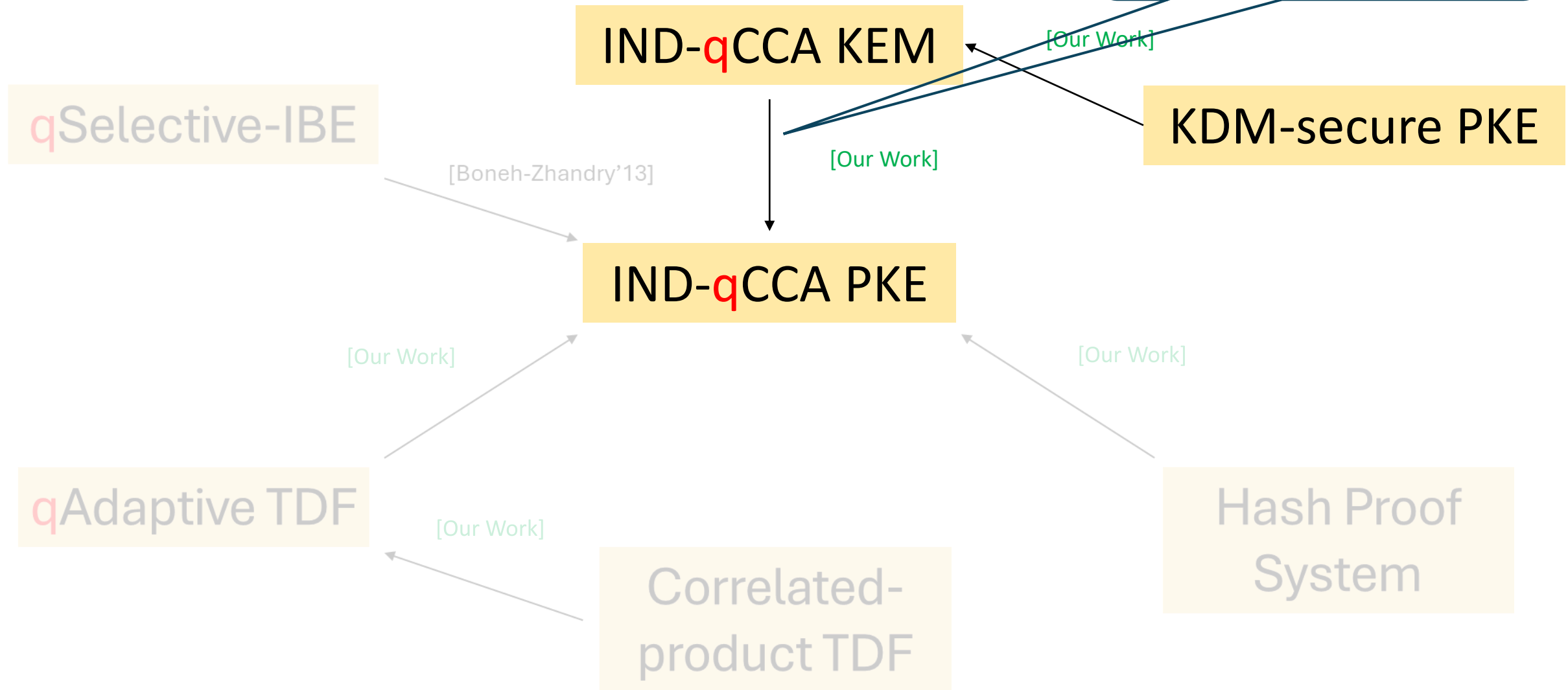


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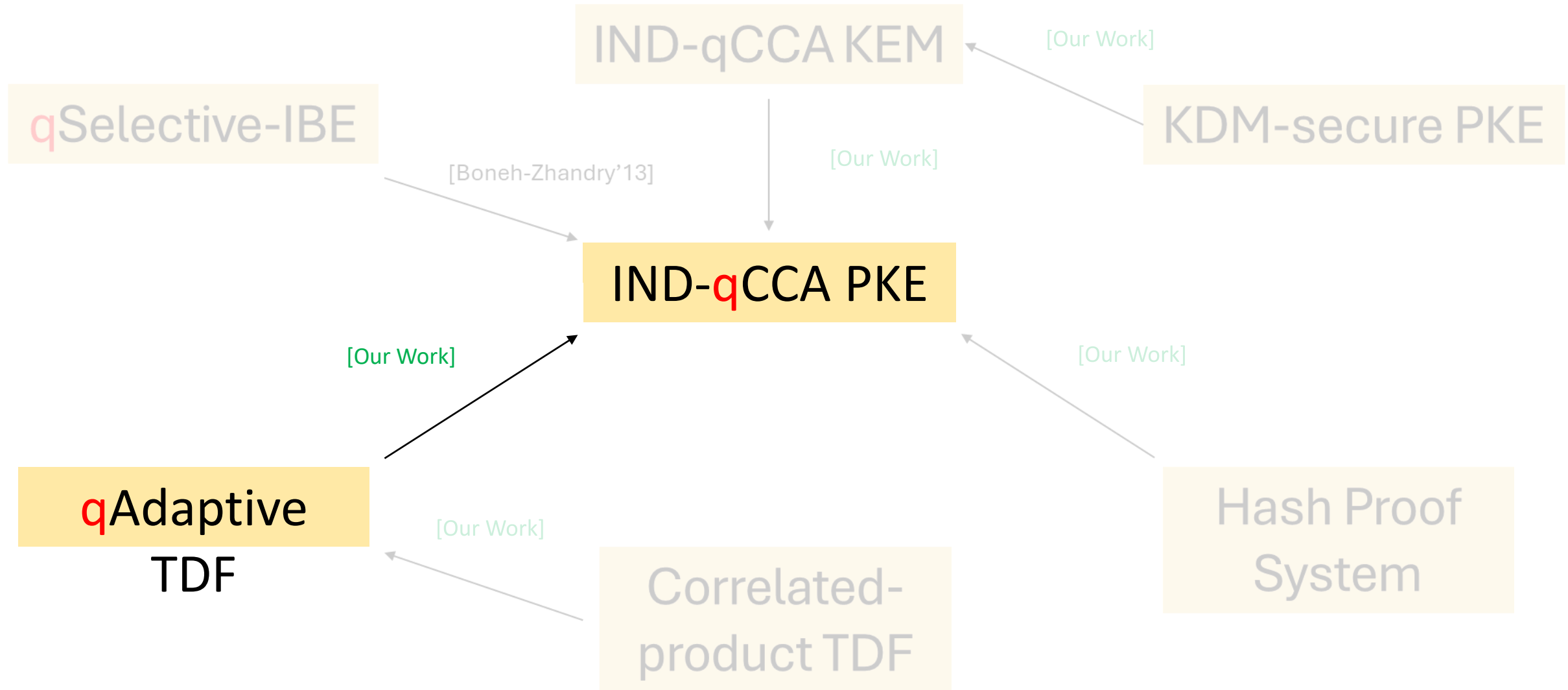


Overview: Results

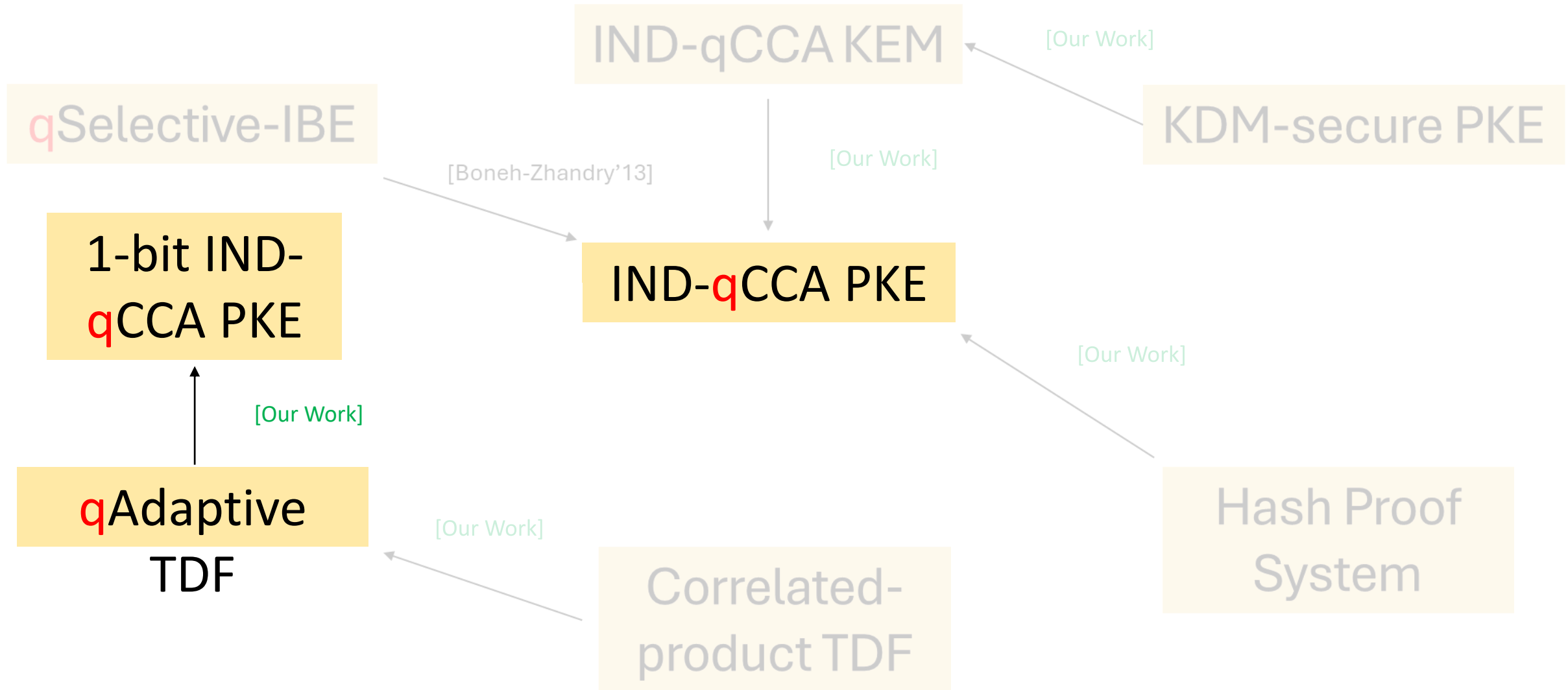
Using the KEM-DEM paradigm of [Cramer-Shoup'03], with only a **post-quantum** secure DEM!



Overview: Results

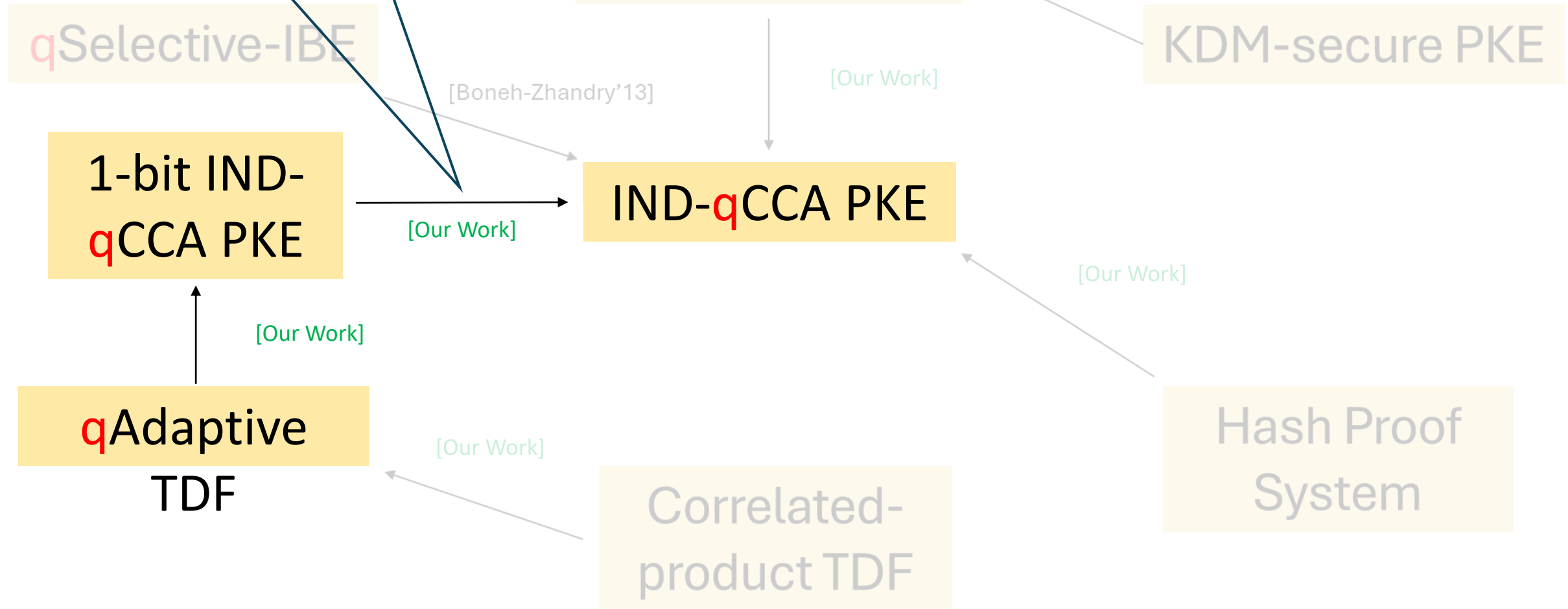


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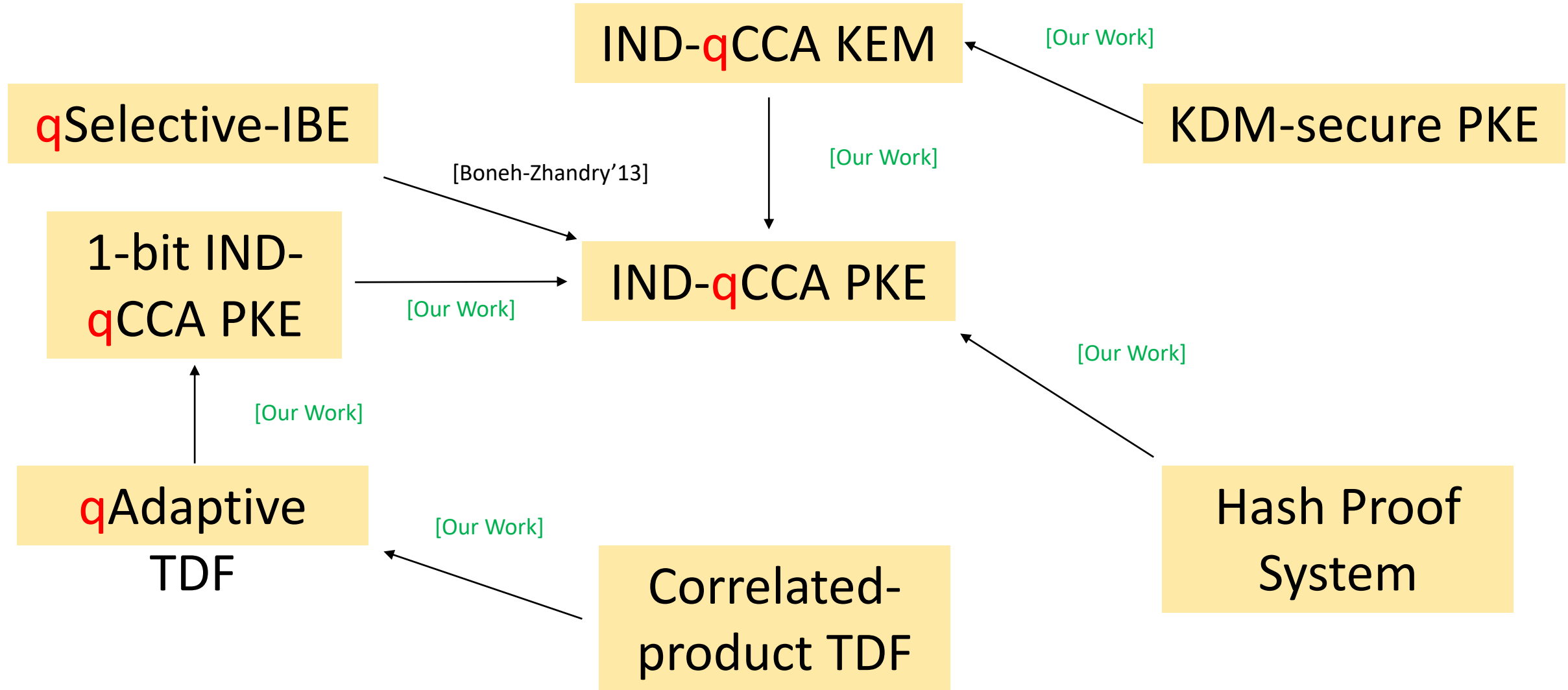


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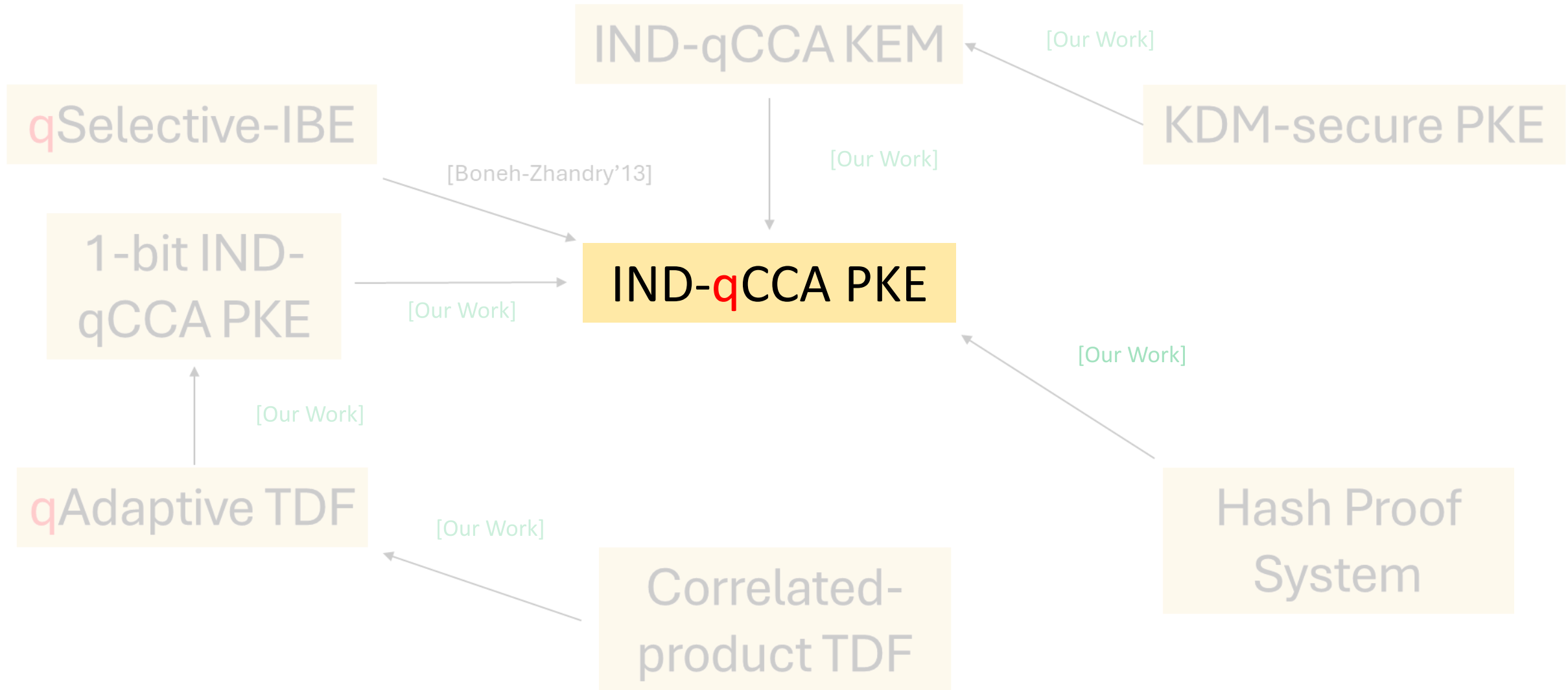
Extending **bit completeness** of CCA-secure PKE by [Hohenberger-Lewko-Waters'12] to the quantum setting.



Overview: Results



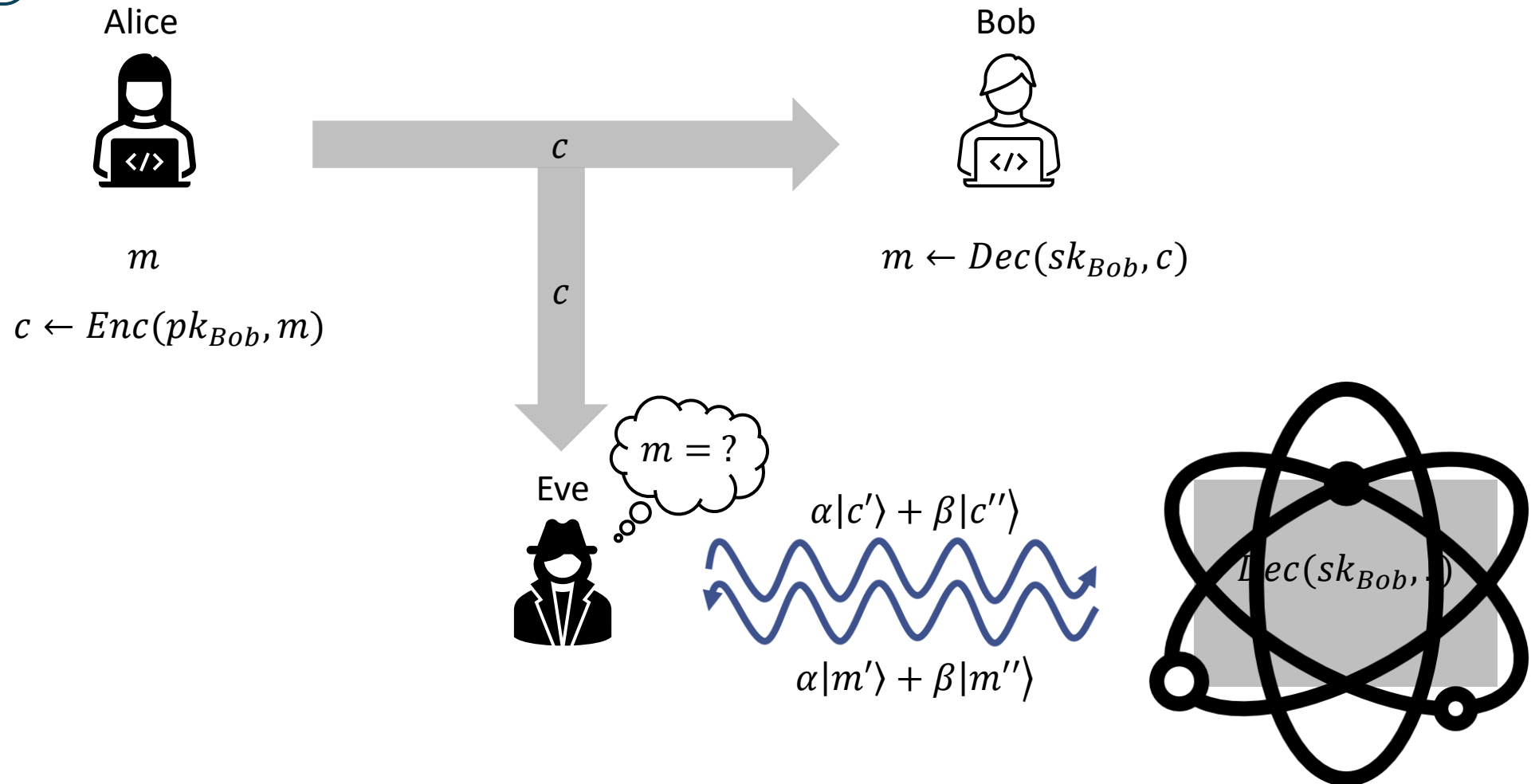
Overview: Results



IND-**q**CCA Security

Introduced by [Boneh-Zhandry'13].

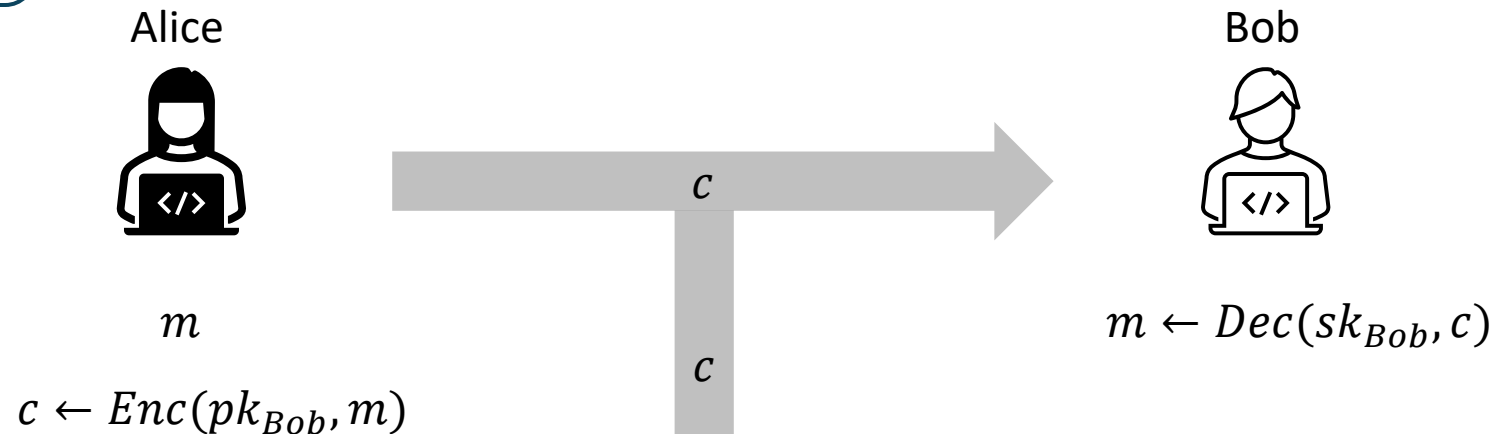
$$PKE = (KGen, Enc, Dec)$$



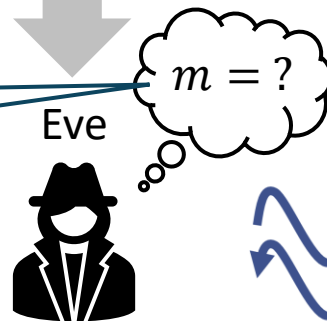
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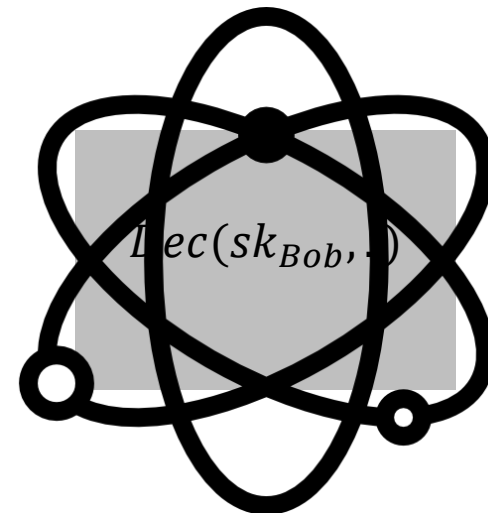
$$PKE = (KGen, Enc, Dec)$$



Notion restricted to **classical** challenge messages.



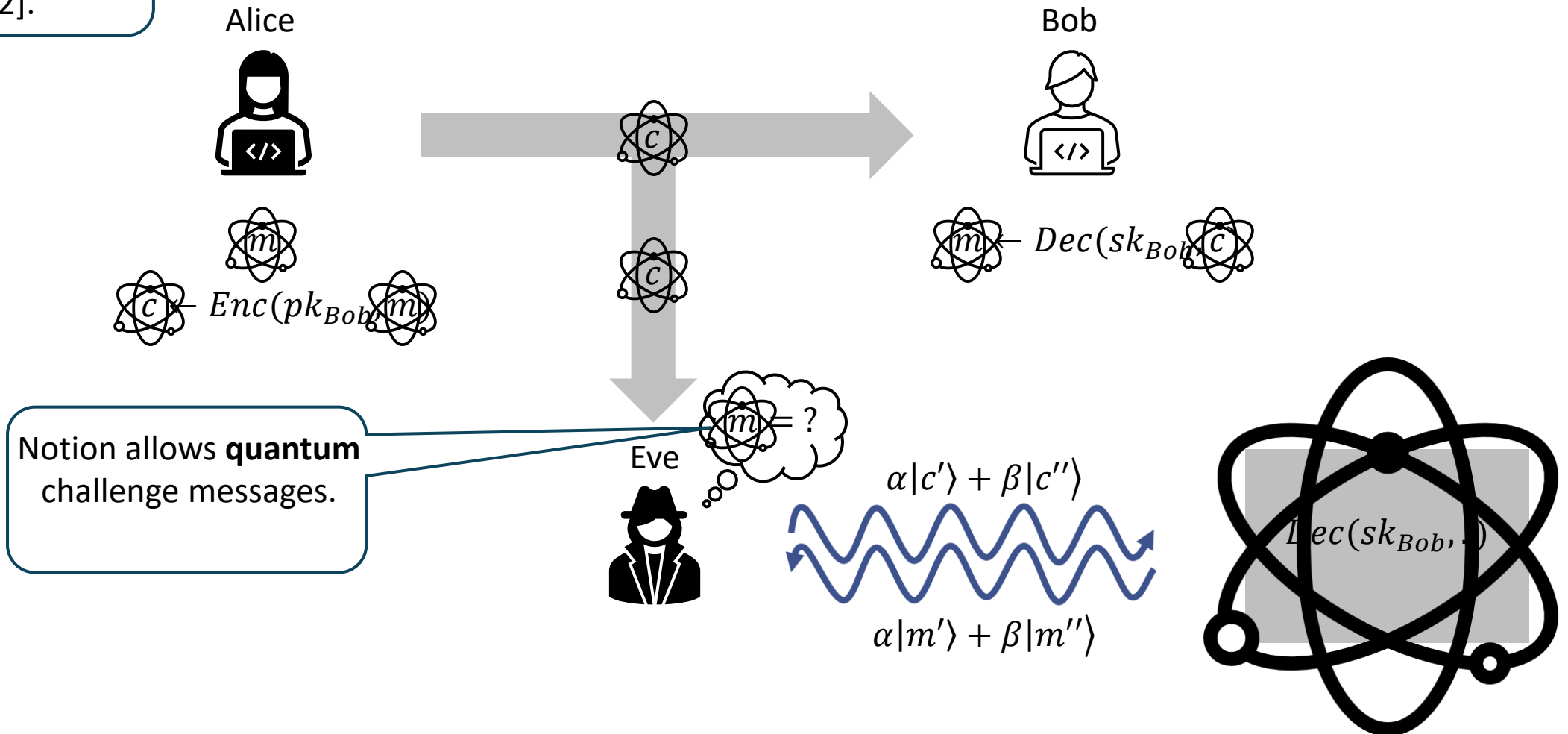
$$\alpha|c'\rangle + \beta|c''\rangle$$
$$\alpha|m'\rangle + \beta|m''\rangle$$



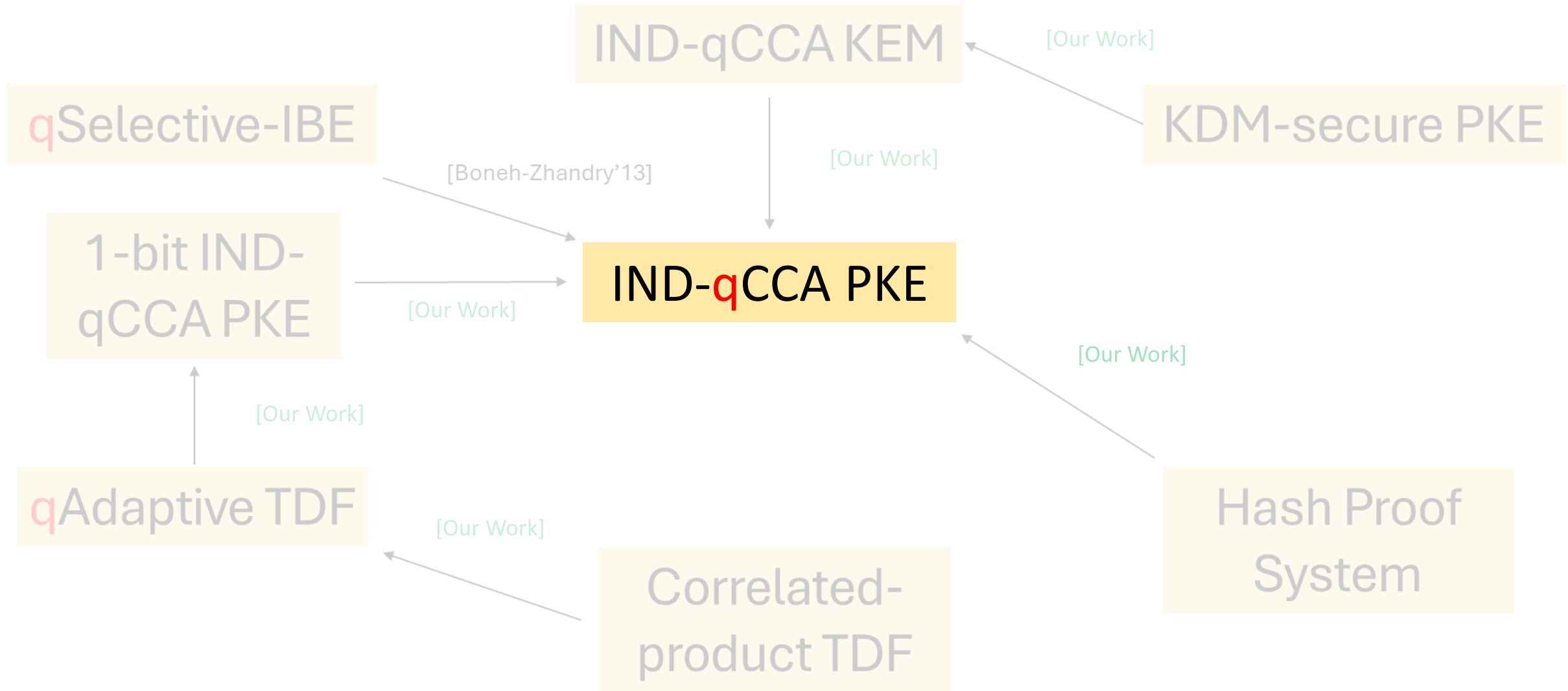
qIND-qCCA Security

Introduced by
[Chevalier-Ebrahimi-
Vu'22].

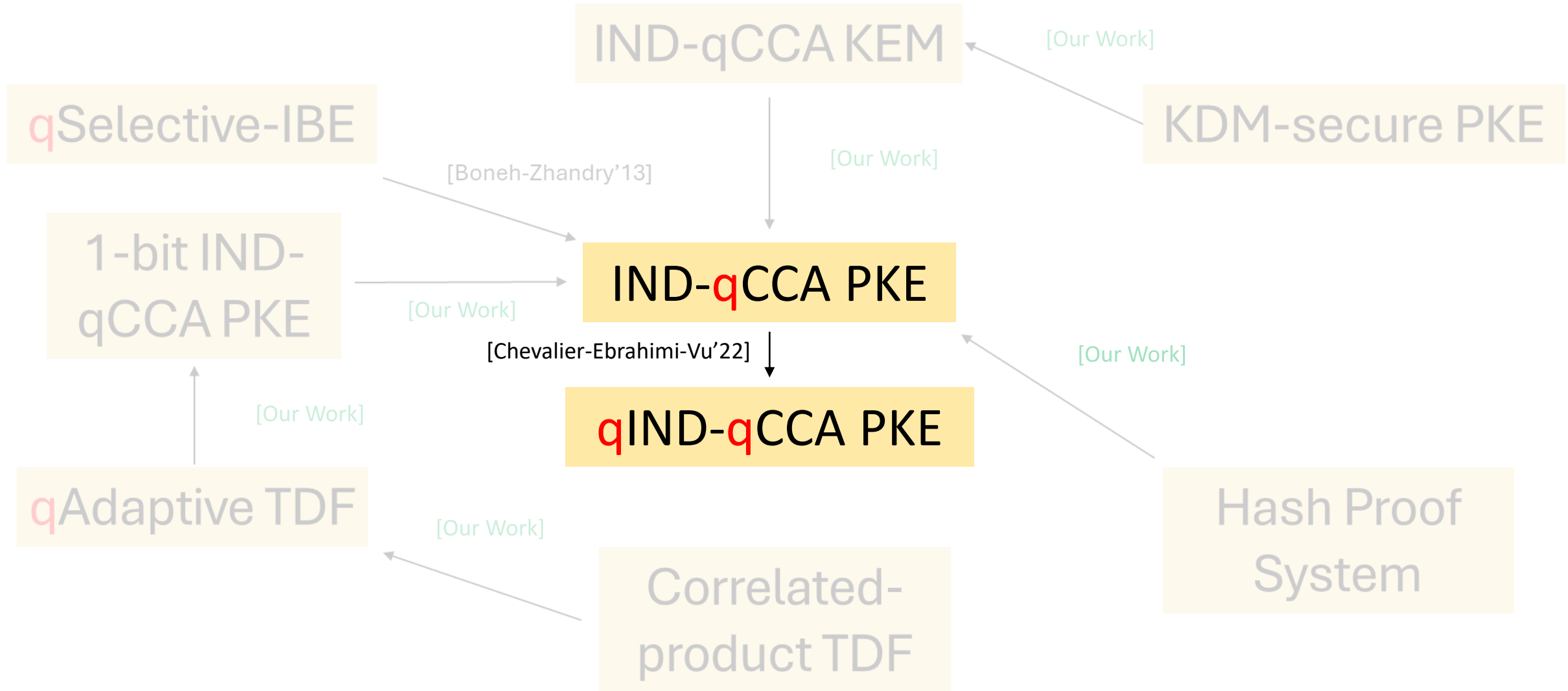
$$PKE = (KGen, Enc, Dec)$$



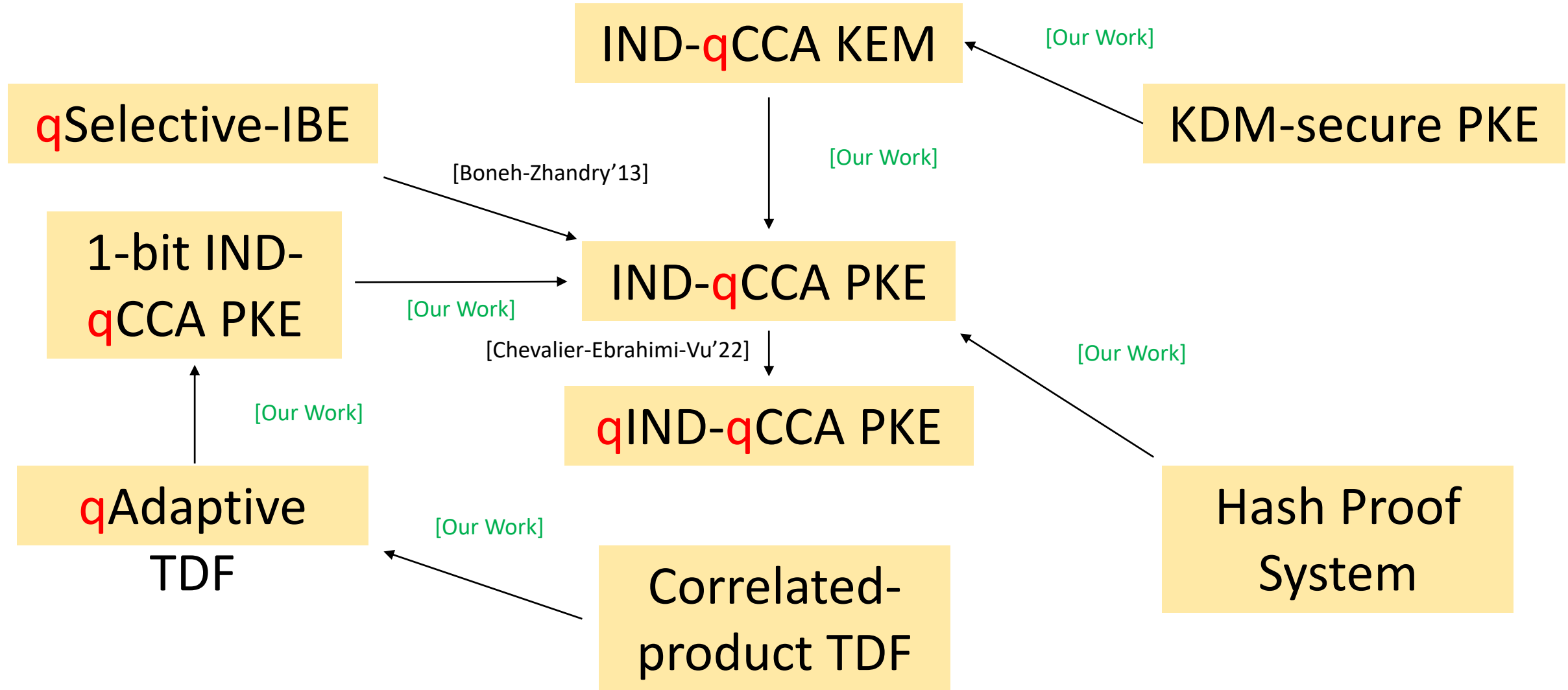
Overview: Results



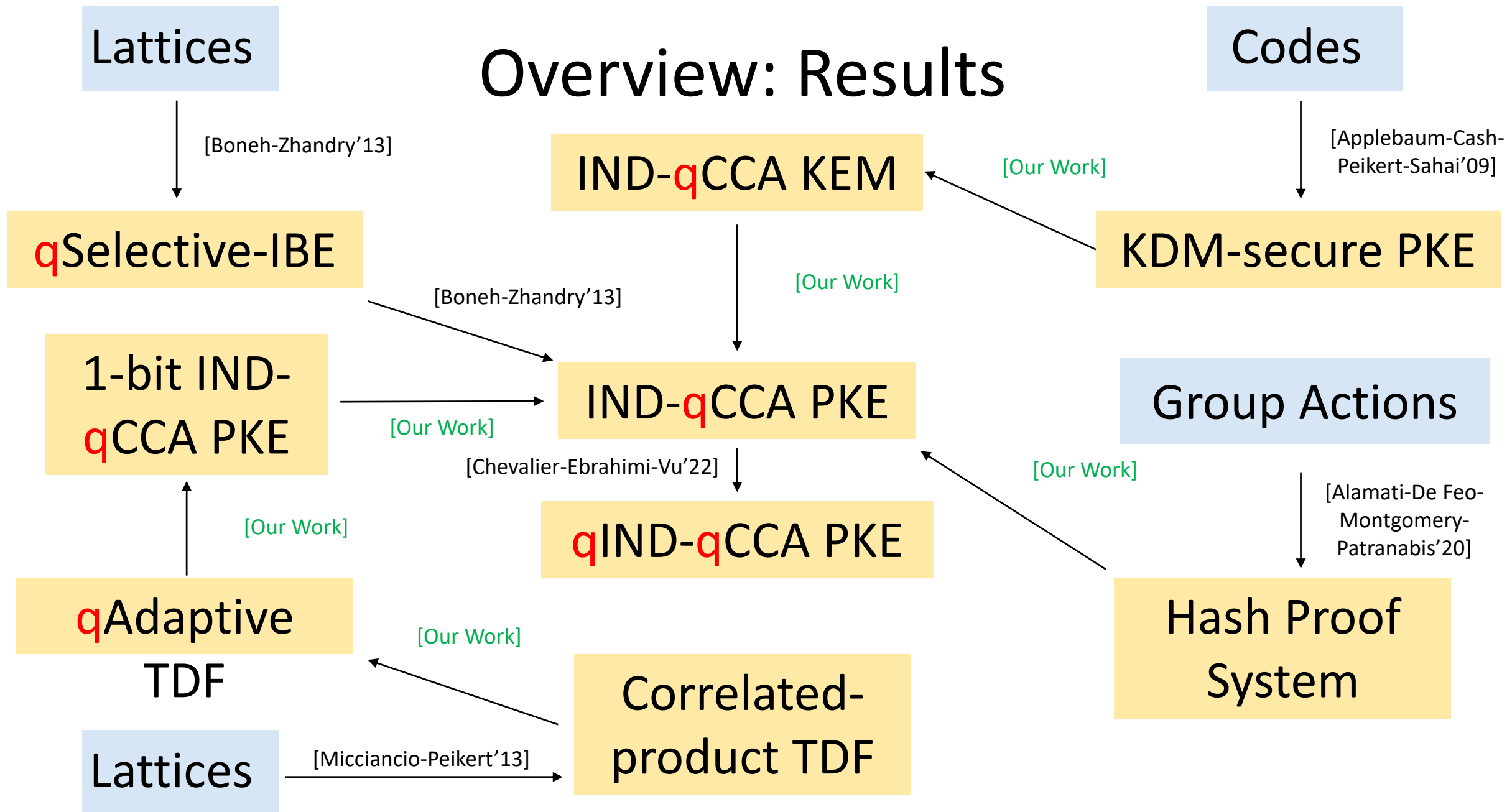
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All our analyzed PKE constructions are **classical**.

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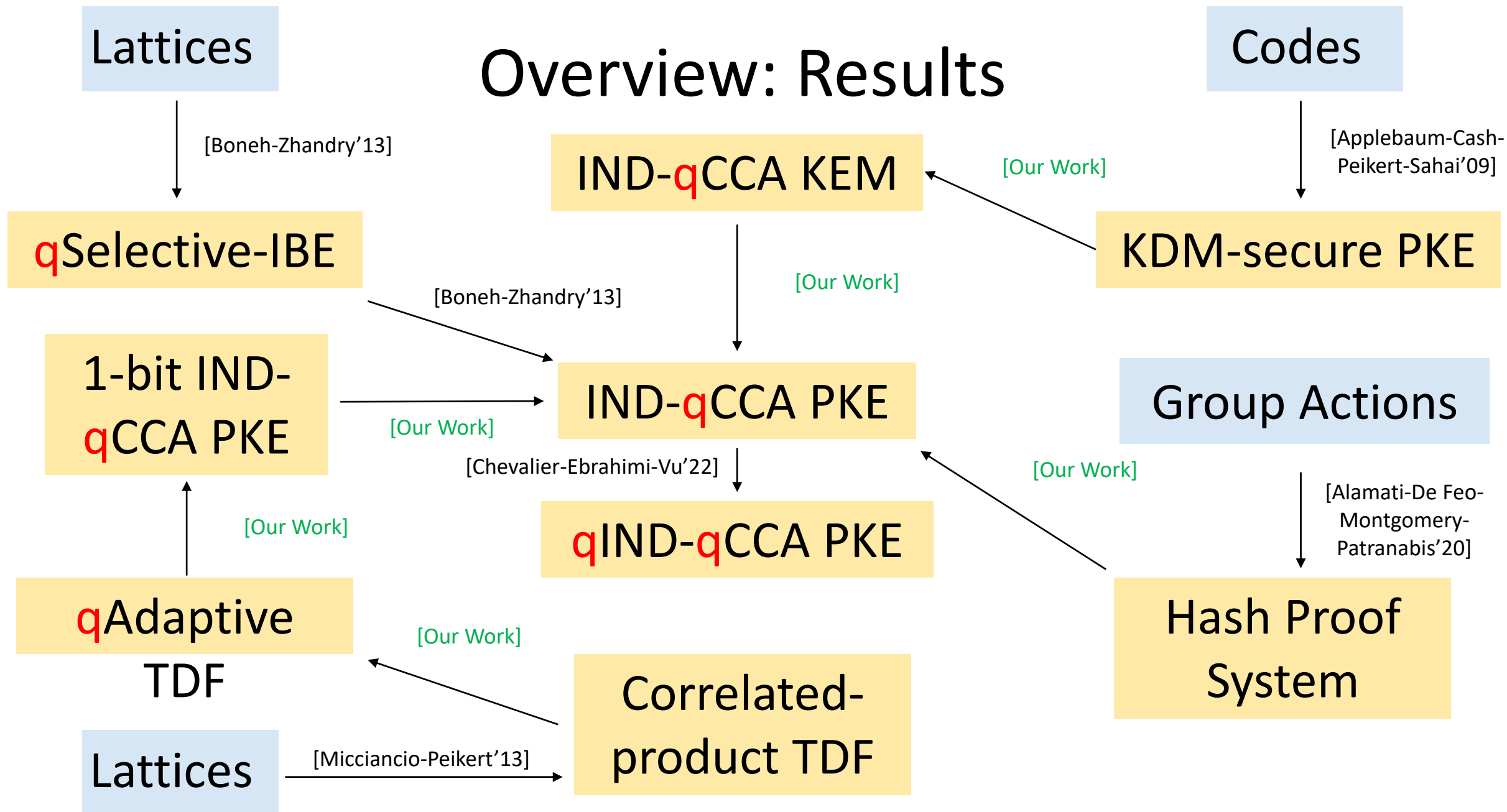
- They can be implemented on **classical computers**.

Overview: Results

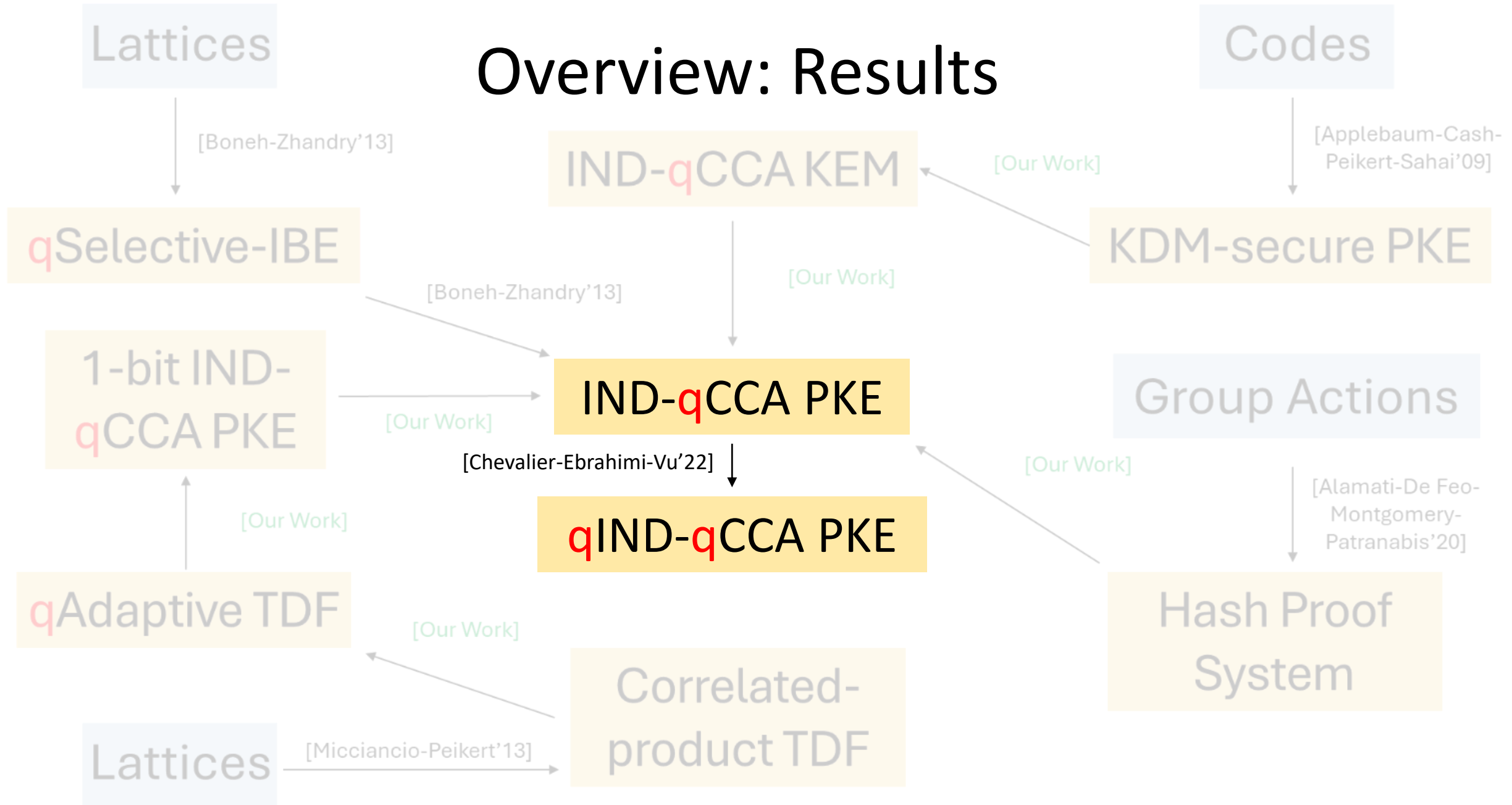
All our analyzed PKE constructions are **classical**.

- They can be implemented on **classical computers**.
- As opposed to **quantum PKE schemes** (e.g., in [Barooti-Grilo-Huguenin(-)Dumittan-Malavolta-Sattath-Vu-Walter'23]) which need inherent quantum components, such as **quantum public keys**.

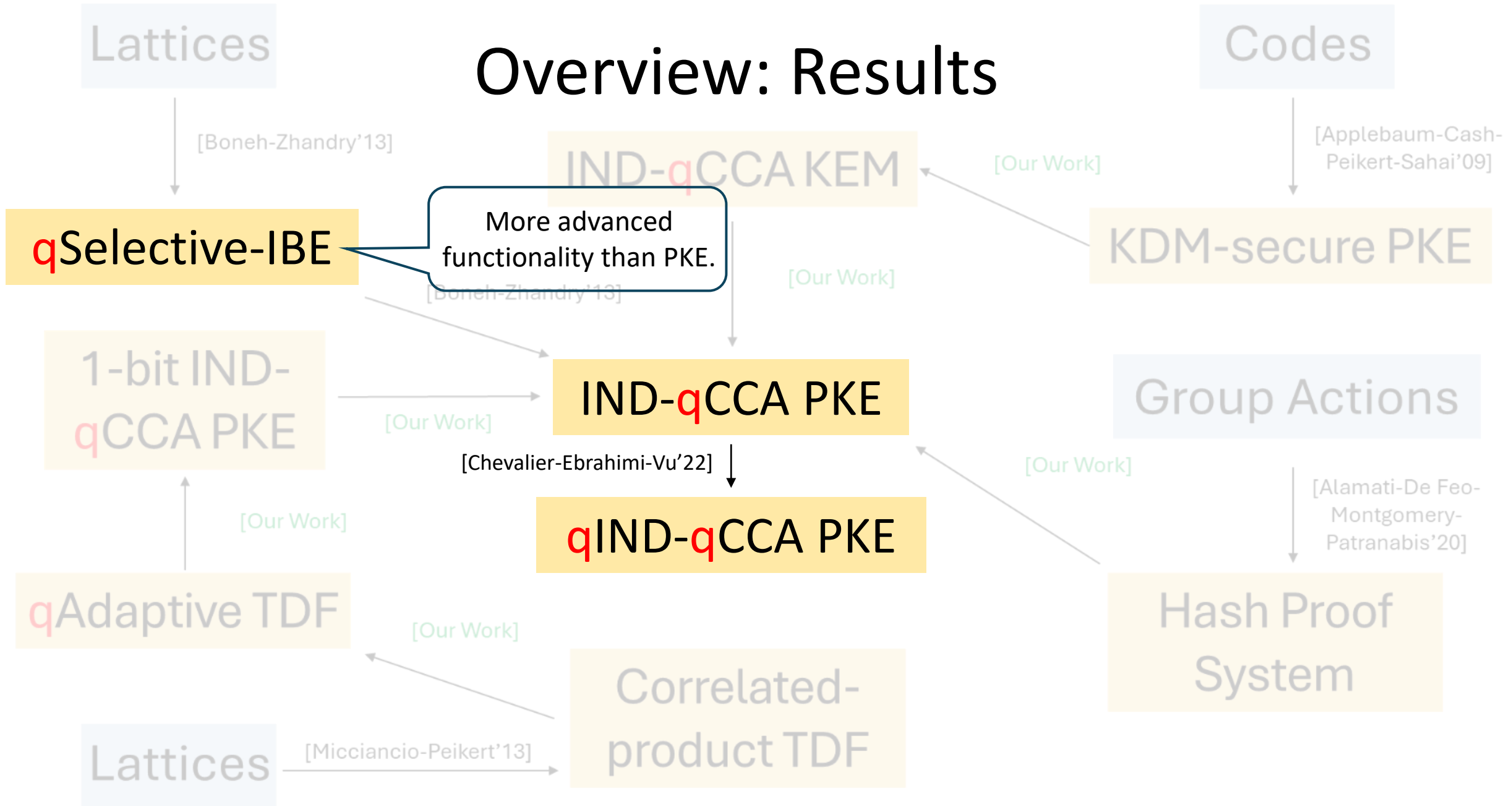
Overview: Results



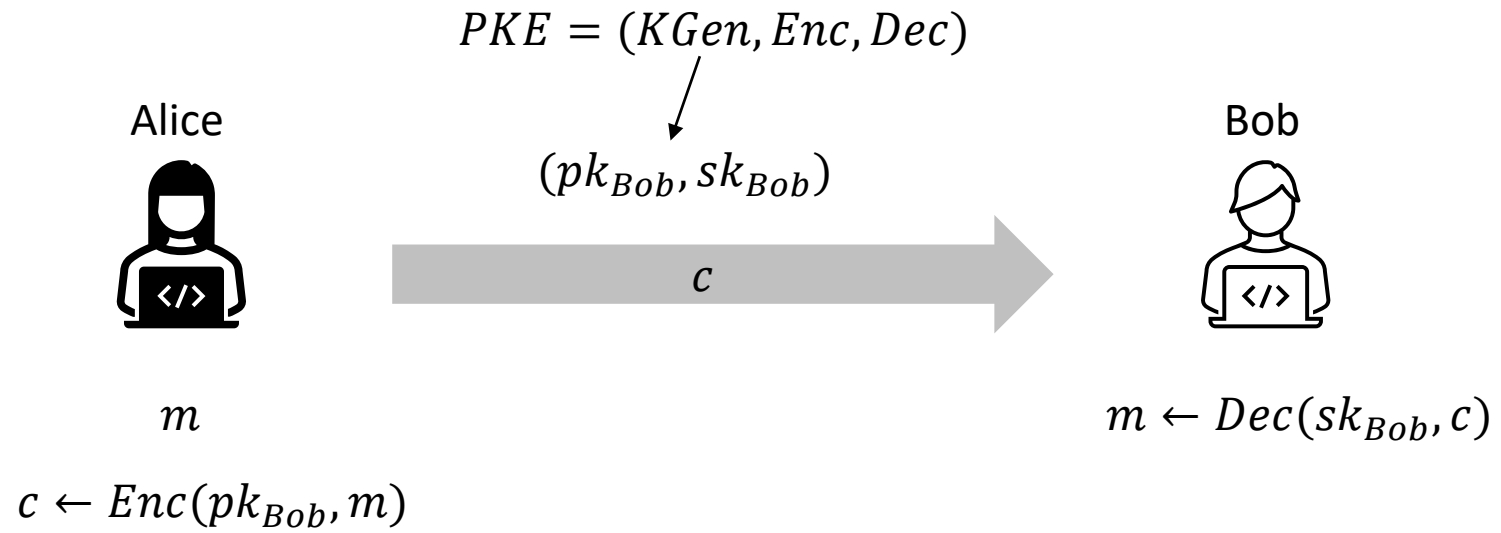
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Overview: Results



PKE

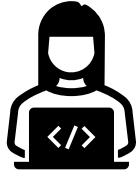


ABE

Attribute-Based
Encryption.

$$ABE = (KGen, Enc, Dec, SKGen)$$

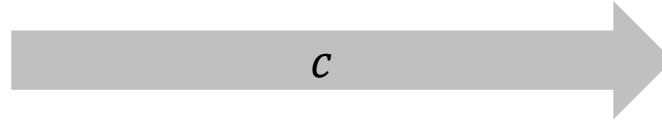
Alice



$m, attr$

$$c \leftarrow Enc(mp_k, m, attr)$$

(mp_k, msk)



c

Bob

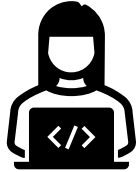


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Bob



msk



ABE

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Alice



$m, attr$

$$c \leftarrow Enc(mp_k, m, attr)$$

(mp_k, msk)

c

Bob



msk

$$sk \leftarrow SKGen(msk, pred)$$

“predicate”

ABE

Attribute-Based
Encryption.

$$ABE = (KGen, Enc, Dec, SKGen)$$

Alice



$m, attr$

$$c \leftarrow Enc(mp_k, m, attr)$$

(mp_k, msk)

c

Bob



$$m \leftarrow Dec(sk, c) \text{ if } pred(attr) = 1$$



msk

$$sk \leftarrow SKGen(msk, pred)$$

"predicate"

ABE

Attribute-Based
Encryption.

$$ABE = (KGen, Enc, Dec, SKGen)$$

Alice



$m, attr$

$$c \leftarrow Enc(mp_k, m, attr)$$

(mp_k, msk)

c

Bob



$$m \leftarrow Dec(sk, c) \text{ if } pred(attr) = 1$$

$$\perp \leftarrow Dec(sk, c) \text{ if } pred(attr) = 0$$

$$sk \leftarrow SKGen(msk, pred)$$



msk

“predicate”

ABE

Attribute-Based
Encryption.

$$ABE = (KGen, Enc, Dec, SKGen)$$

Alice



$m, attr$

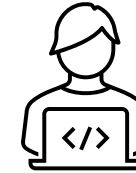
$$c \leftarrow Enc(mp_k, m, attr)$$

"Bob"

(mp_k, msk)

c

Bob



$$sk \leftarrow SKGen(msk, pred)$$

$$m \leftarrow Dec(sk, c) \text{ if } pred(attr) = 1$$
$$\perp \leftarrow Dec(sk, c) \text{ if } pred(attr) = 0$$

$$pred("Bob") = 1$$
$$pred(\cdot) = 0, \text{ otherwise}$$



ABE

Attribute-Based
Encryption.

$$ABE = (KGen, Enc, Dec, SKGen)$$

Alice



$m, attr$

$$c \leftarrow Enc(mp_k, m, attr)$$

"Dave"

(mp_k, msk)

c

Bob



$$m \leftarrow Dec(sk, c) \text{ if } pred(attr) = 1$$

$$\perp \leftarrow Dec(sk, c) \text{ if } pred(attr) = 0$$

$$sk \leftarrow SKGen(msk, pred)$$



$$pred("Bob") = 1$$

$$pred(\cdot) = 0, \text{ otherwise}$$

ABE

Attribute-Based
Encryption.

Implies Identity-Based
Encryption (IBE).

$$ABE = (KGen, Enc, Dec, SKGen)$$

Alice



$m, attr$

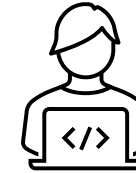
$$c \leftarrow Enc(mp_k, m, attr)$$

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$$m \leftarrow Dec(sk, c) \text{ if } pred(attr) = 1$$

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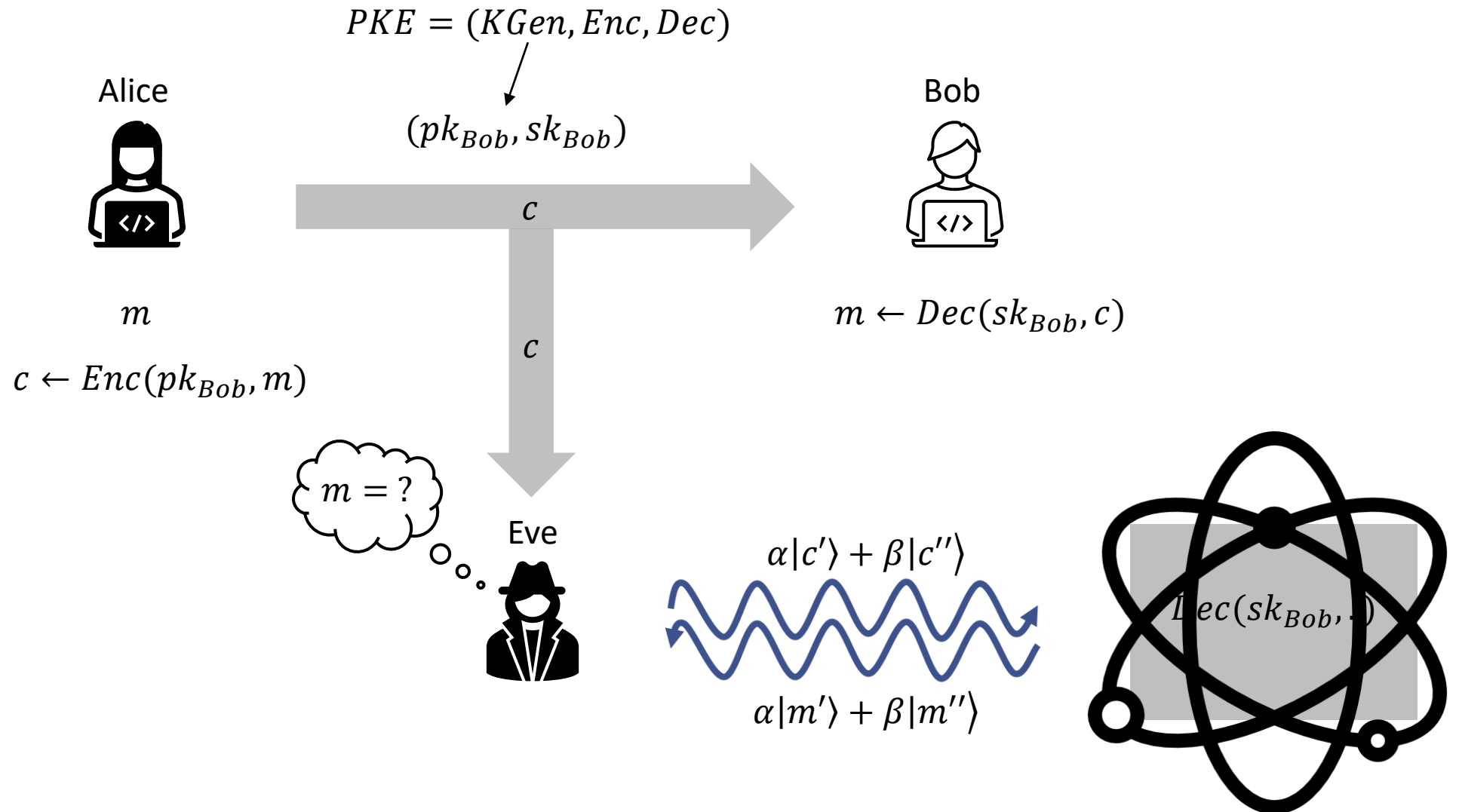
$$sk \leftarrow SKGen(msk, pred)$$



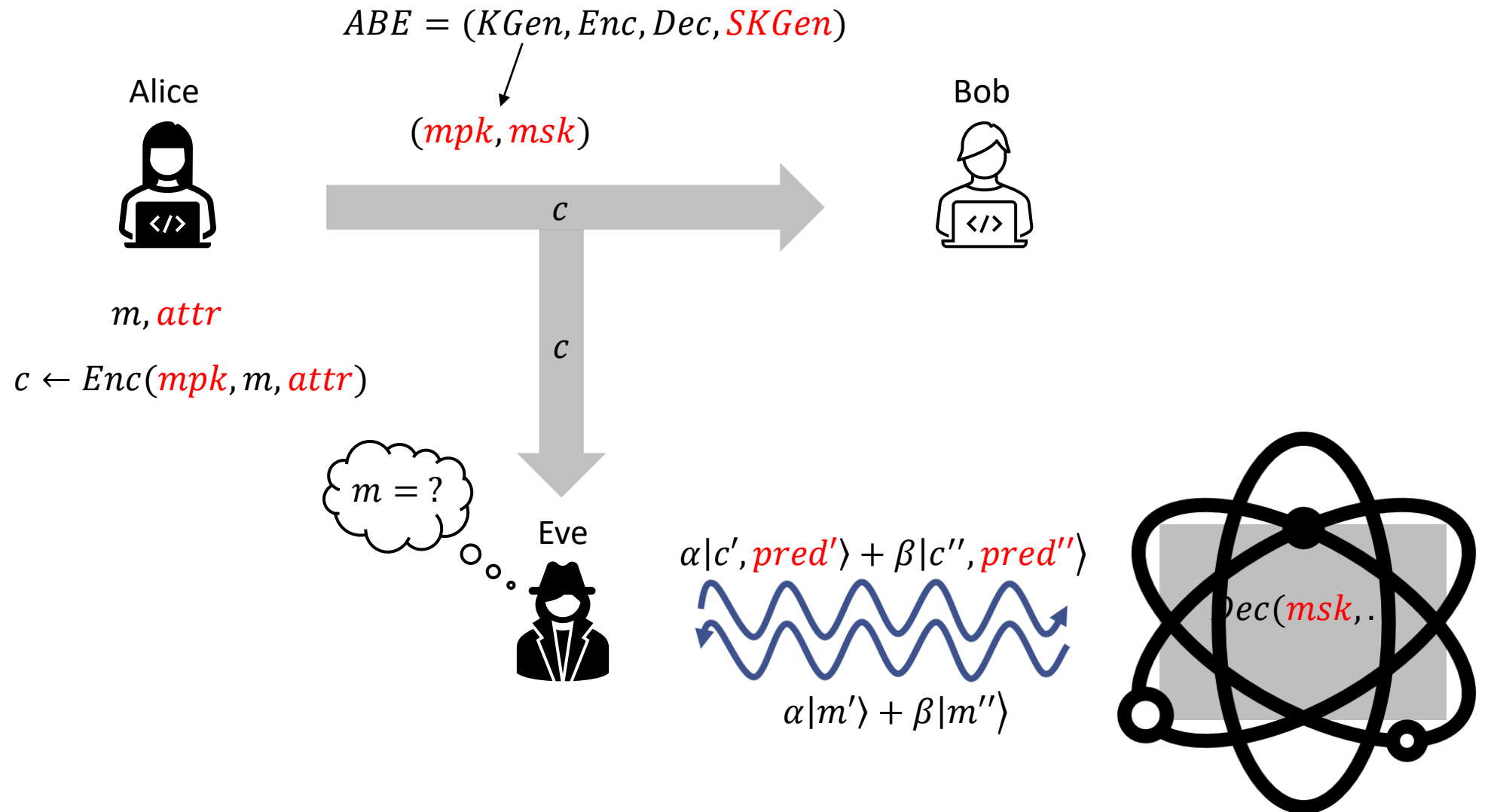
msk

$pred("Bob") = 1$
 $pred(\cdot) = 0, \text{ otherwise}$

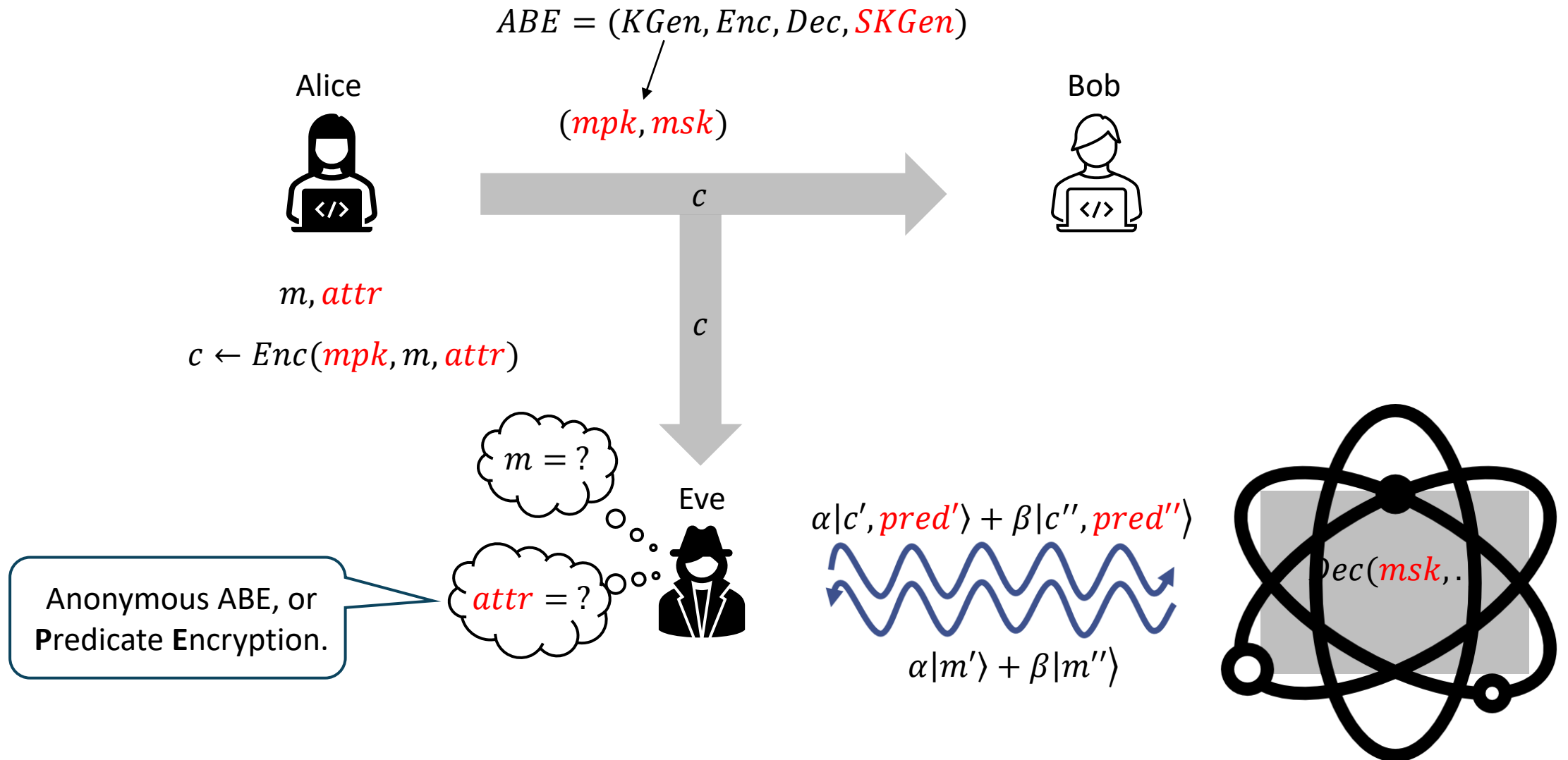
IND-**q**CCA Secure PKE



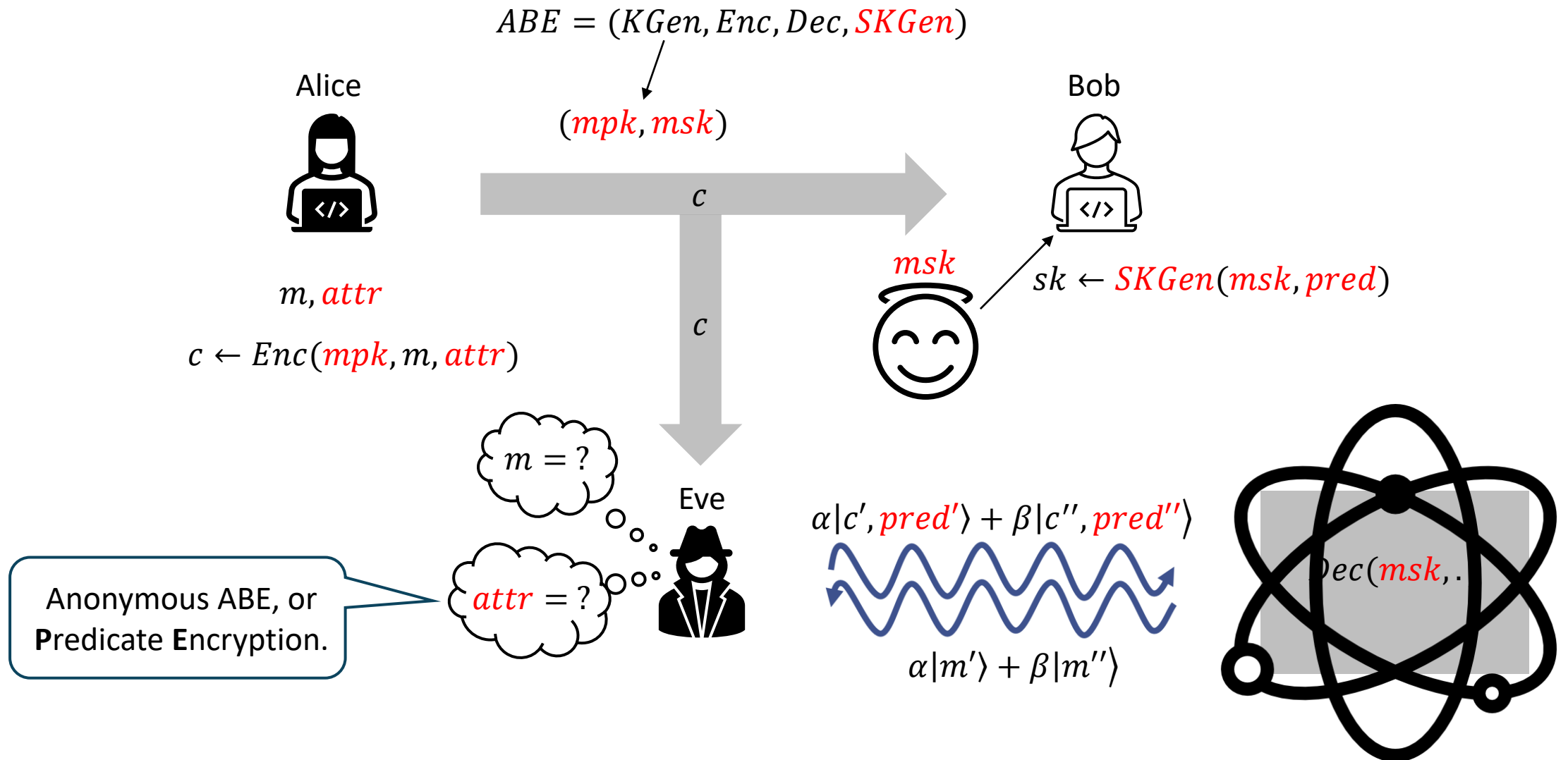
IND-**q**CCA Secure ABE



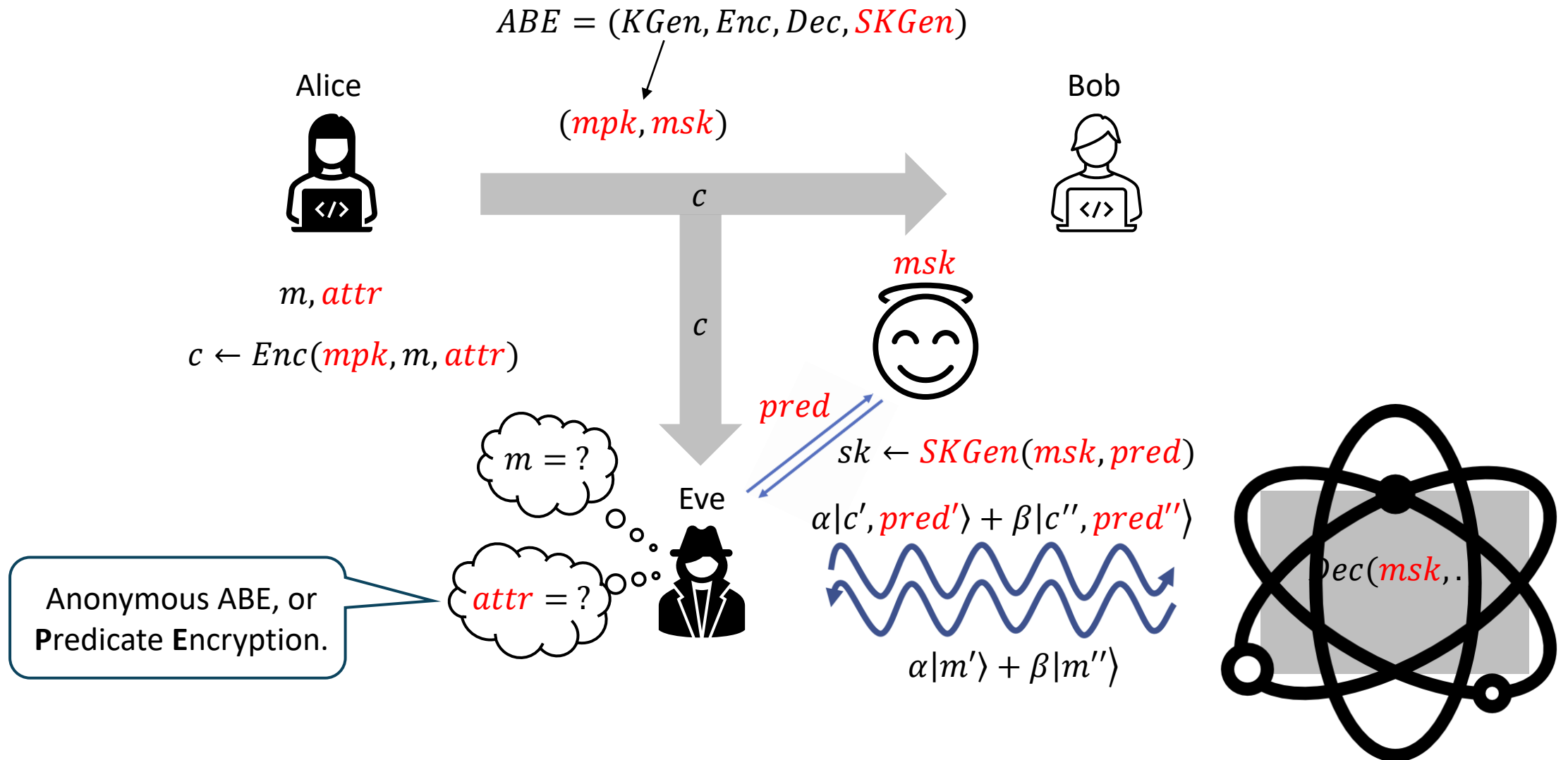
IND-**q**CCA Secure ABE



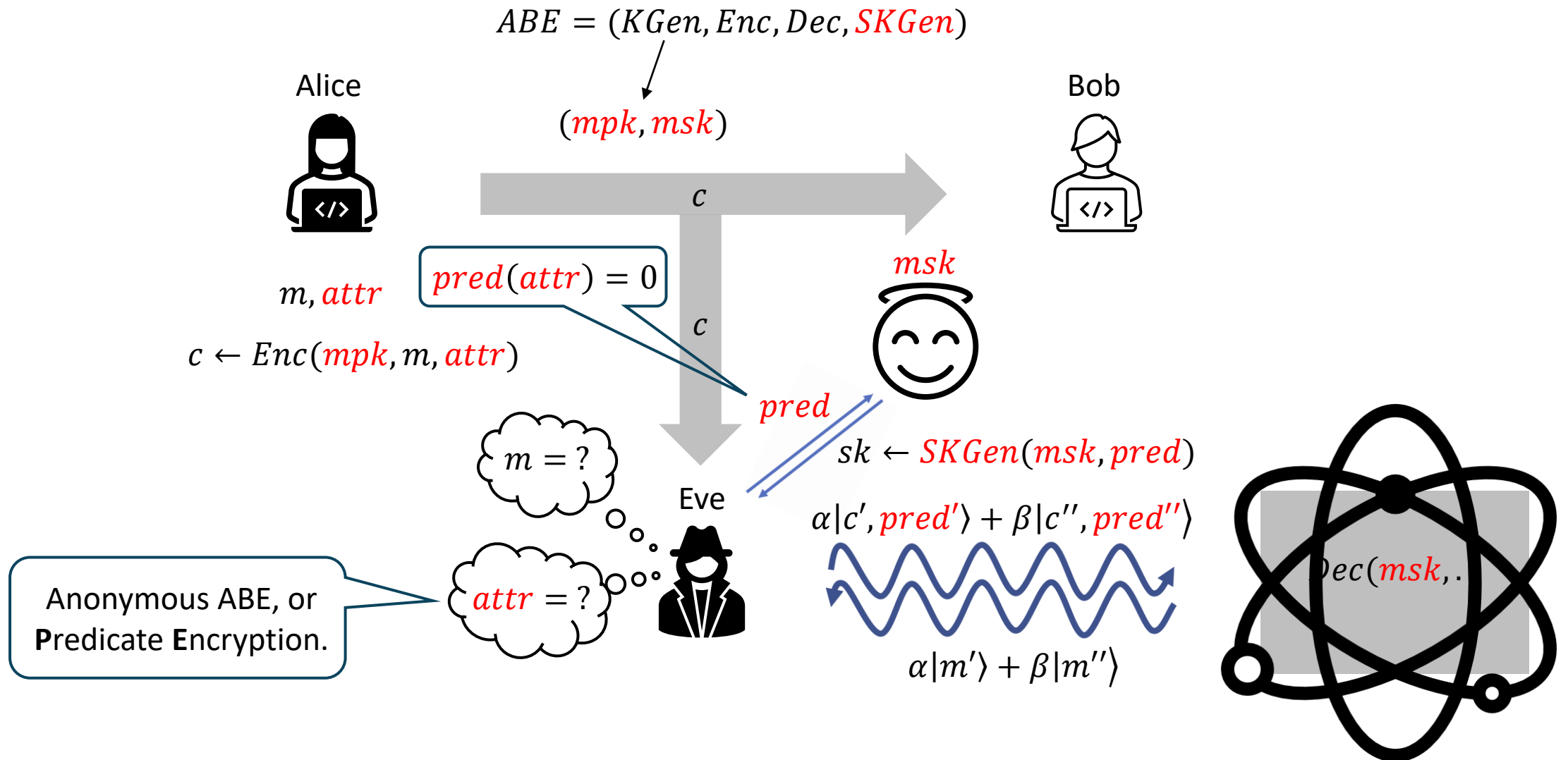
IND-**q**CCA Secure ABE



IND-**q**CCA Secure ABE



IND-**q**CCA Secure ABE



IND-**q**CCA-**c**KG Secure ABE

Adversary can only make “classical” **SKGen** queries.

$$ABE = (KGen, Enc, Dec, \textcolor{red}{SKGen})$$

Alice



$m, \textcolor{red}{attr}$

$$c \leftarrow Enc(\textcolor{red}{mpk}, m, \textcolor{red}{attr})$$

$$\textcolor{red}{pred(attr)} = 0$$

c

c

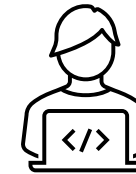
$\textcolor{red}{pred}$

$$sk \leftarrow \textcolor{red}{SKGen}(msk, \textcolor{red}{pred})$$

$$\alpha|c', \textcolor{red}{pred}'\rangle + \beta|c'', \textcolor{red}{pred}''\rangle$$

$$\alpha|m'\rangle + \beta|m''\rangle$$

Bob



$\textcolor{red}{msk}$

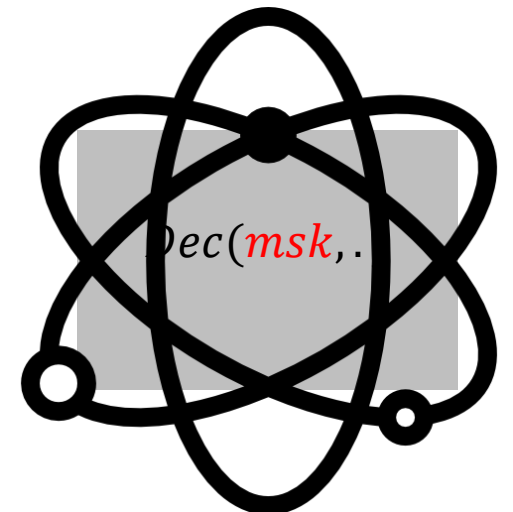


Anonymous ABE, or
Predicate Encryption.

$m = ?$

$\textcolor{red}{attr} = ?$

Eve

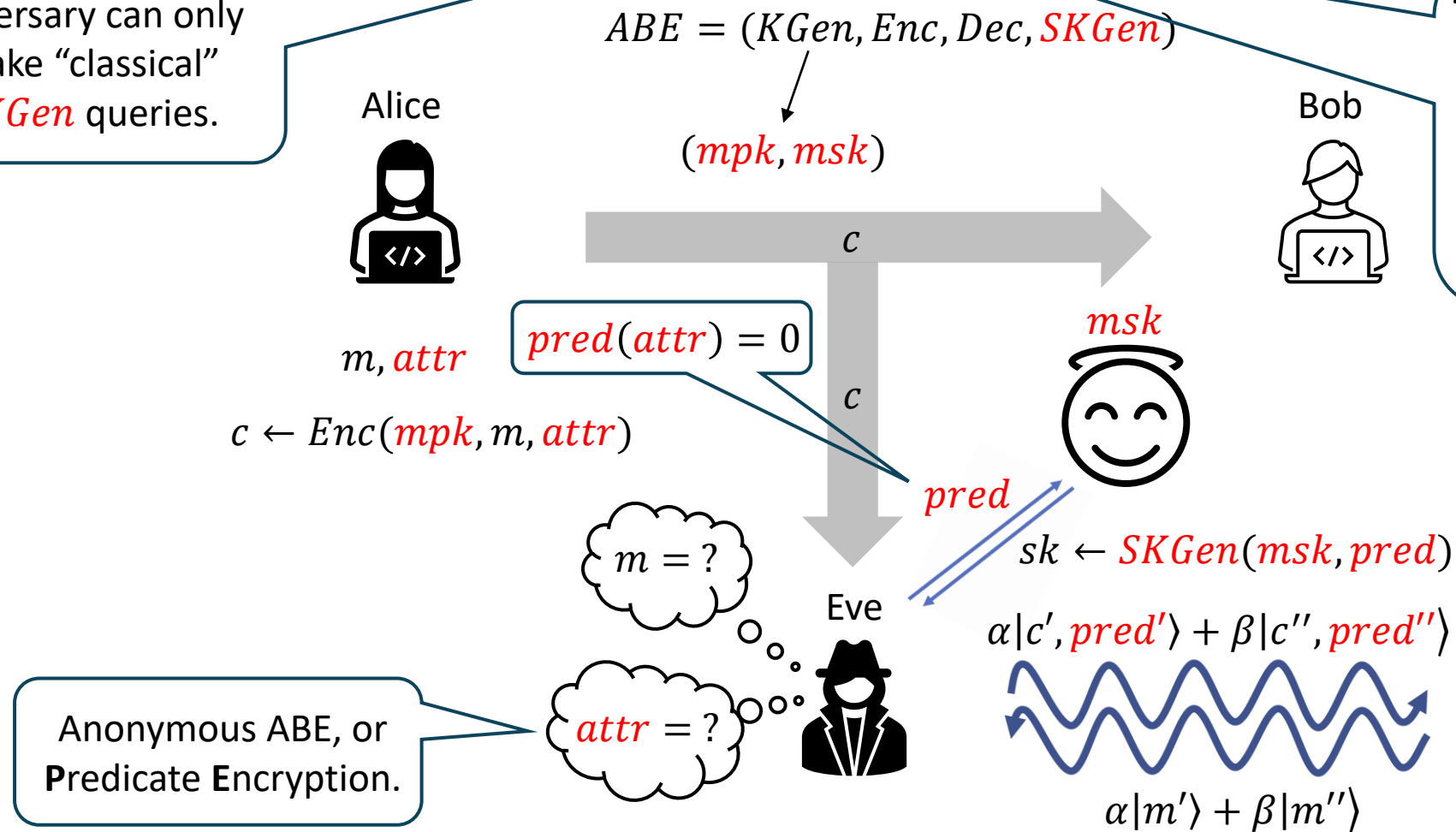


$$Dec(\textcolor{red}{msk}, .)$$

IND-**q**CCA-**c**KG Secure ABE

Adversary can only make “classical” **SKGen** queries.

“Quantum” **SKGen** queries come with definitional issues...



IND-**q**CCA-**c**KG Secure ABE

Adversary can only make “classical” **SKGen** queries.

“Quantum” **SKGen** queries come with definitional issues...

... but can be salvaged in a “semi-adaptive” setting.

$ABE = (KGen, Enc, Dec, \textcolor{red}{SKGen})$

Alice



$m, \textcolor{red}{attr}$

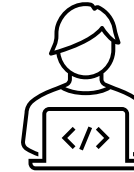
$\textcolor{red}{pred(attr)} = 0$

$c \leftarrow Enc(\textcolor{red}{mpk}, m, \textcolor{red}{attr})$

c

c

Bob



$\textcolor{red}{msk}$



$\textcolor{red}{pred}$

$sk \leftarrow \textcolor{red}{SKGen}(\textcolor{red}{msk}, \textcolor{red}{pred})$

$\alpha|c', \textcolor{red}{pred}'\rangle + \beta|c'', \textcolor{red}{pred}''\rangle$

$\alpha|m'\rangle + \beta|m''\rangle$

Eve

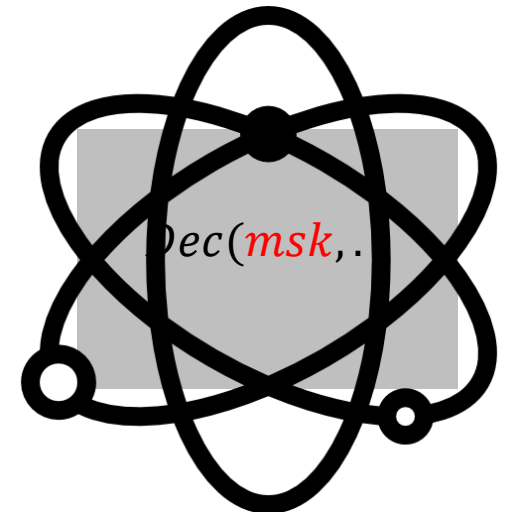


$m = ?$

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Anonymous ABE, or Predicate Encryption.

$Dec(\textcolor{red}{msk}, .)$



IND-**q**CCA-**c**KG Secure ABE

IND-CPA(-cKG) ABE

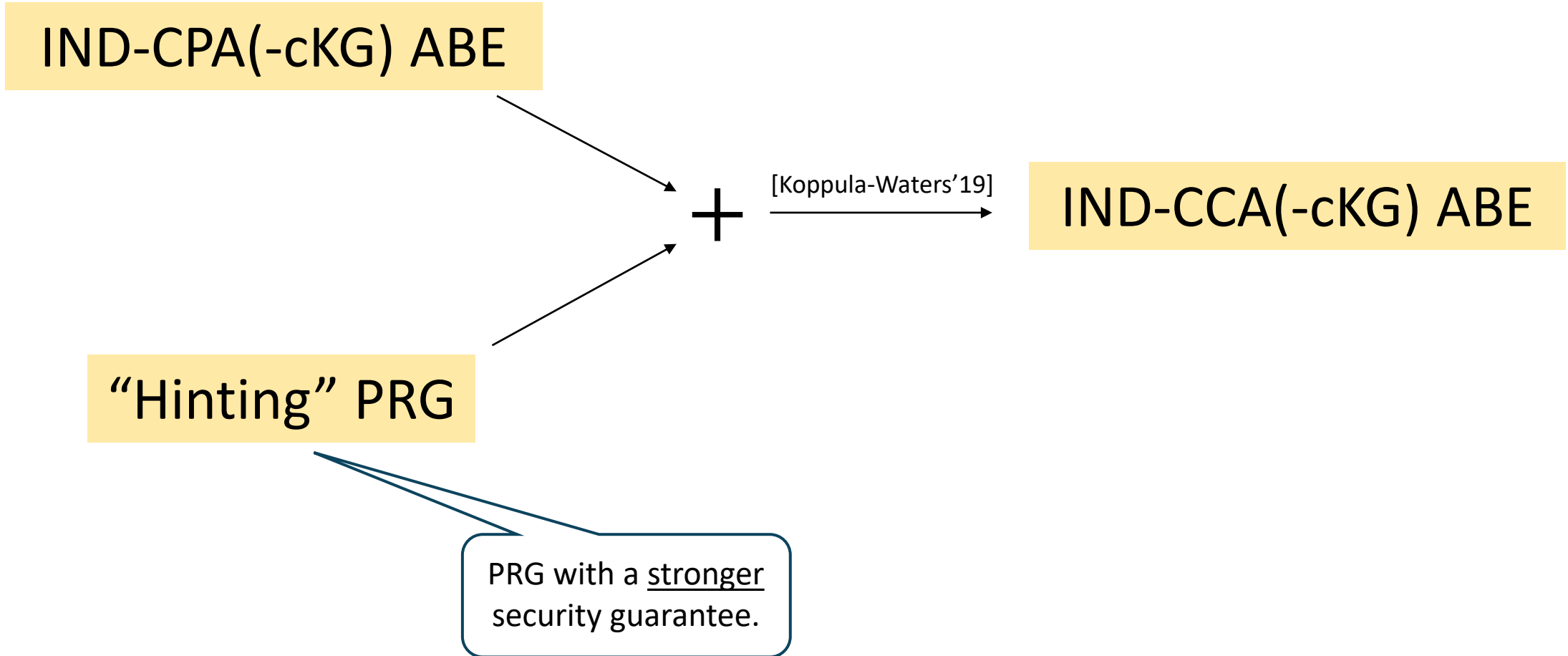
IND- q CCA- c KG Secure ABE

IND-CPA(-cKG) ABE

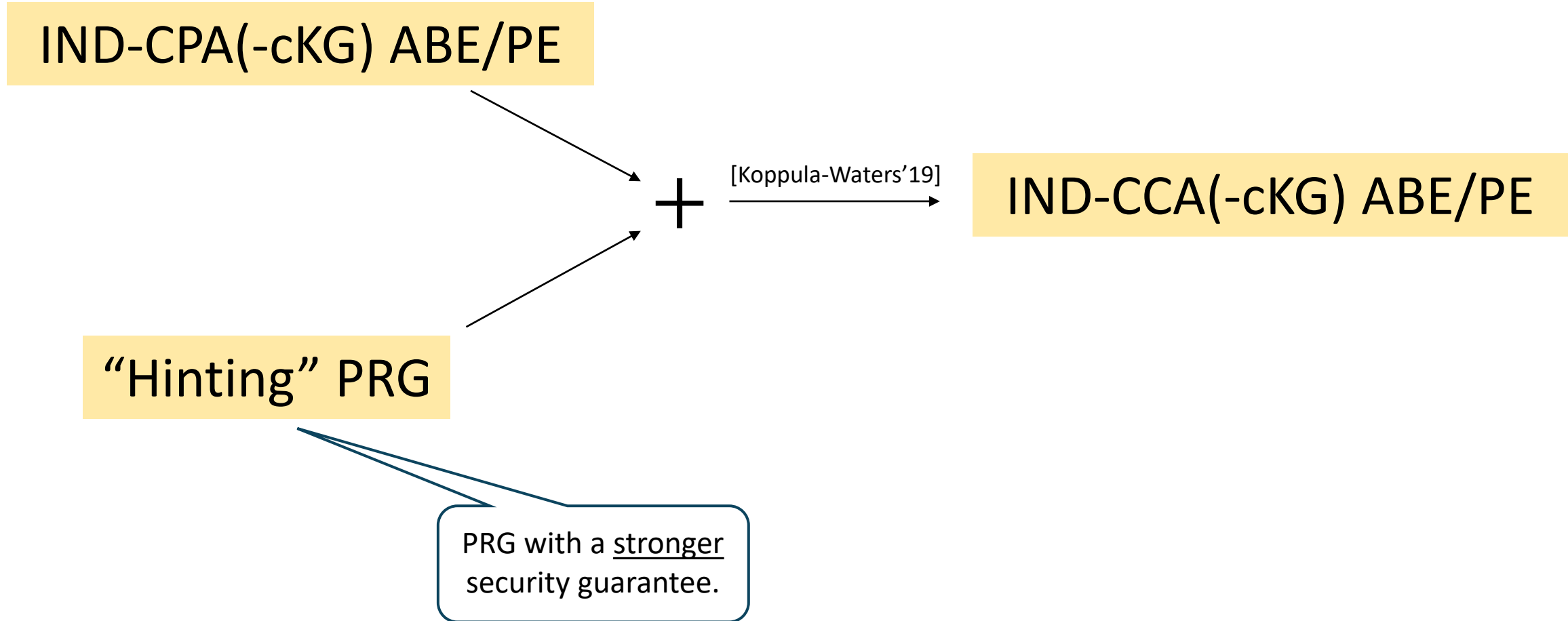
“Hinting” PRG

PRG with a stronger
security guarantee.

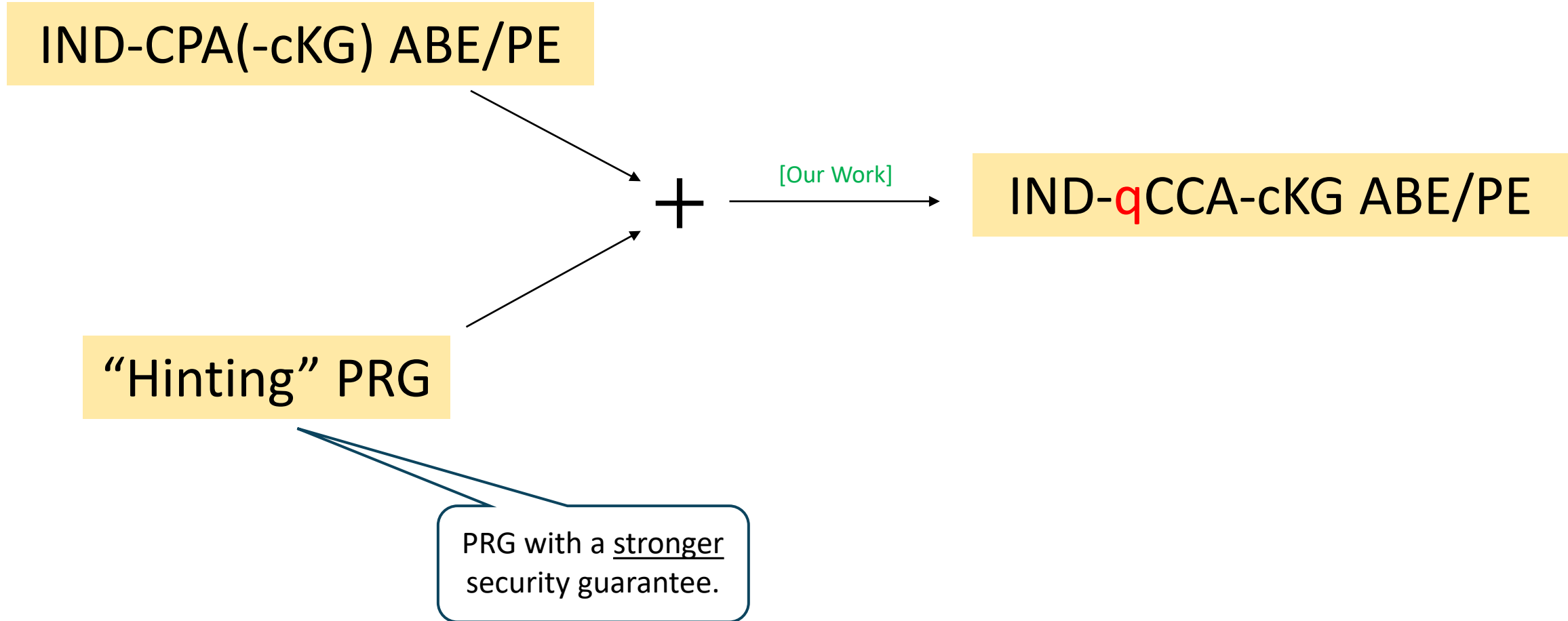
IND-**q**CCA-**c**KG Secure ABE



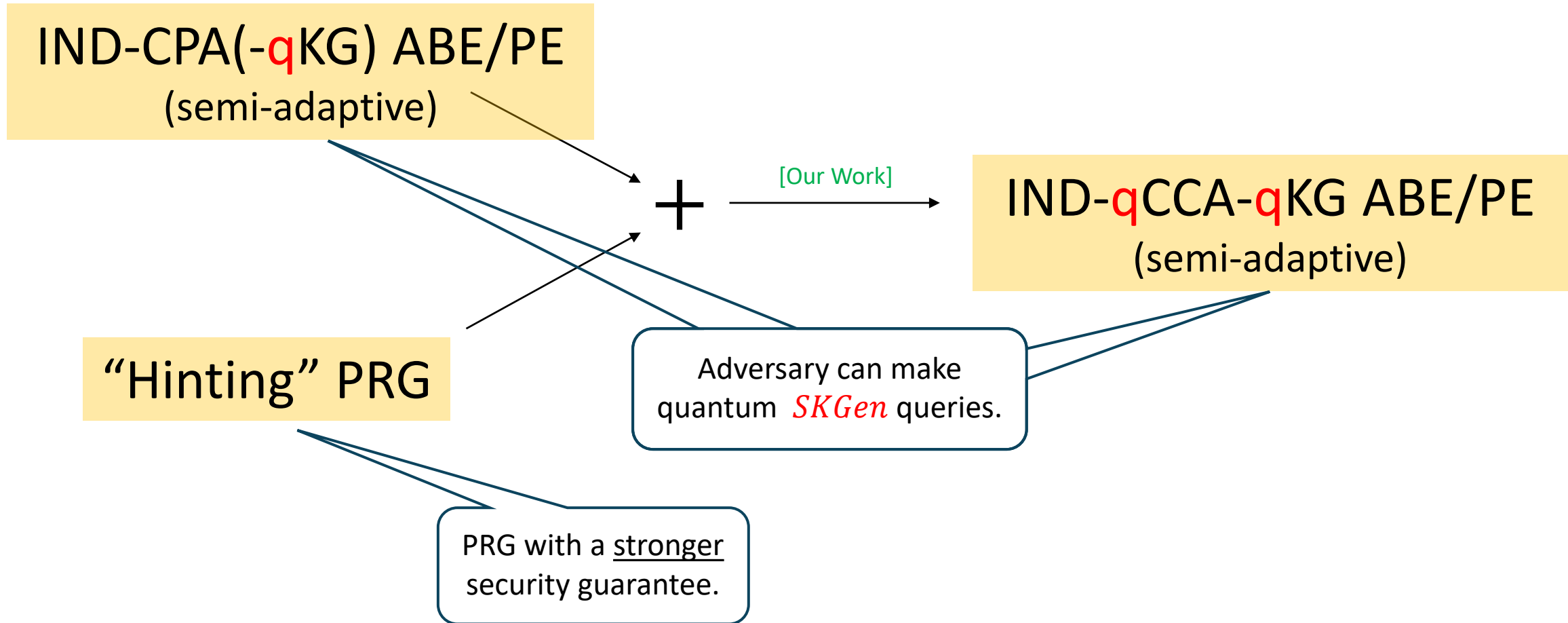
IND-**q**CCA-**c**KG Secure ABE



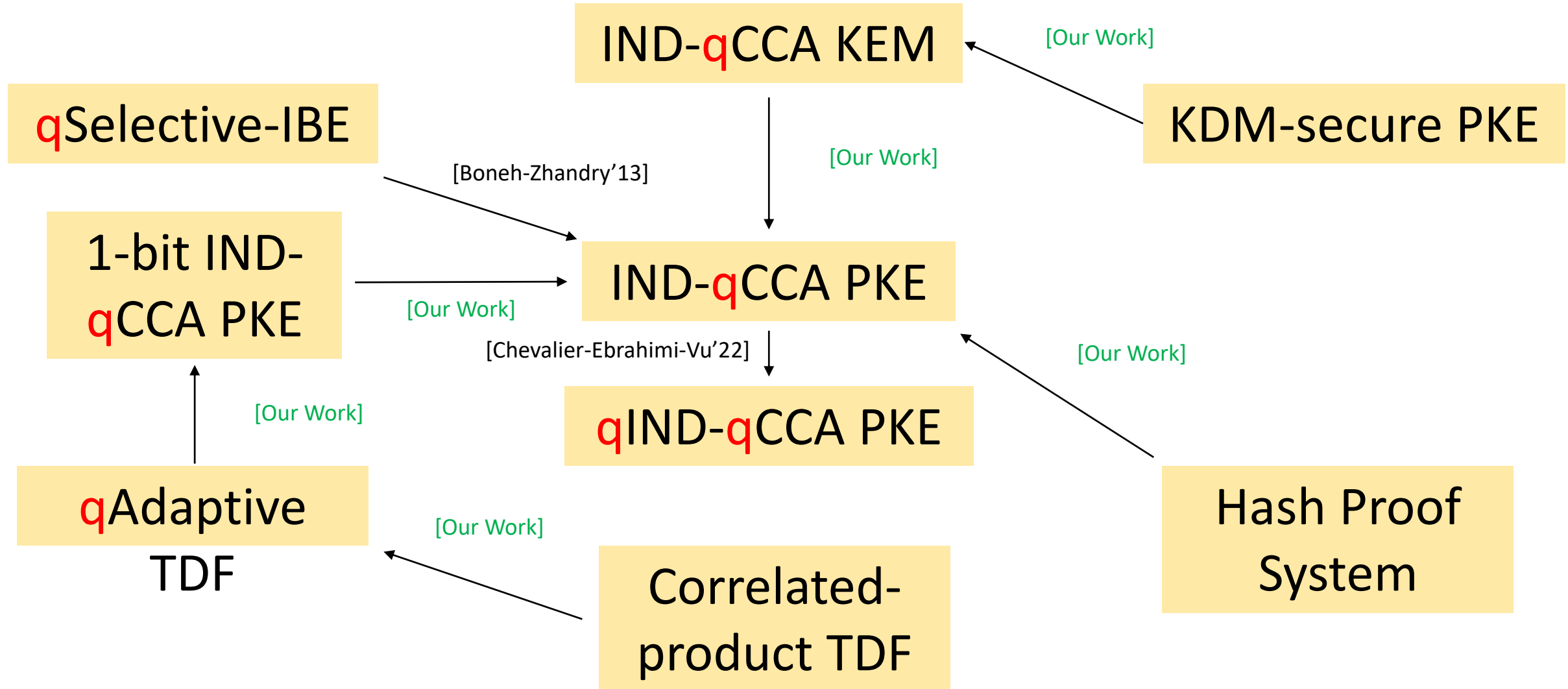
IND-**q**CCA-**c**KG Secure ABE



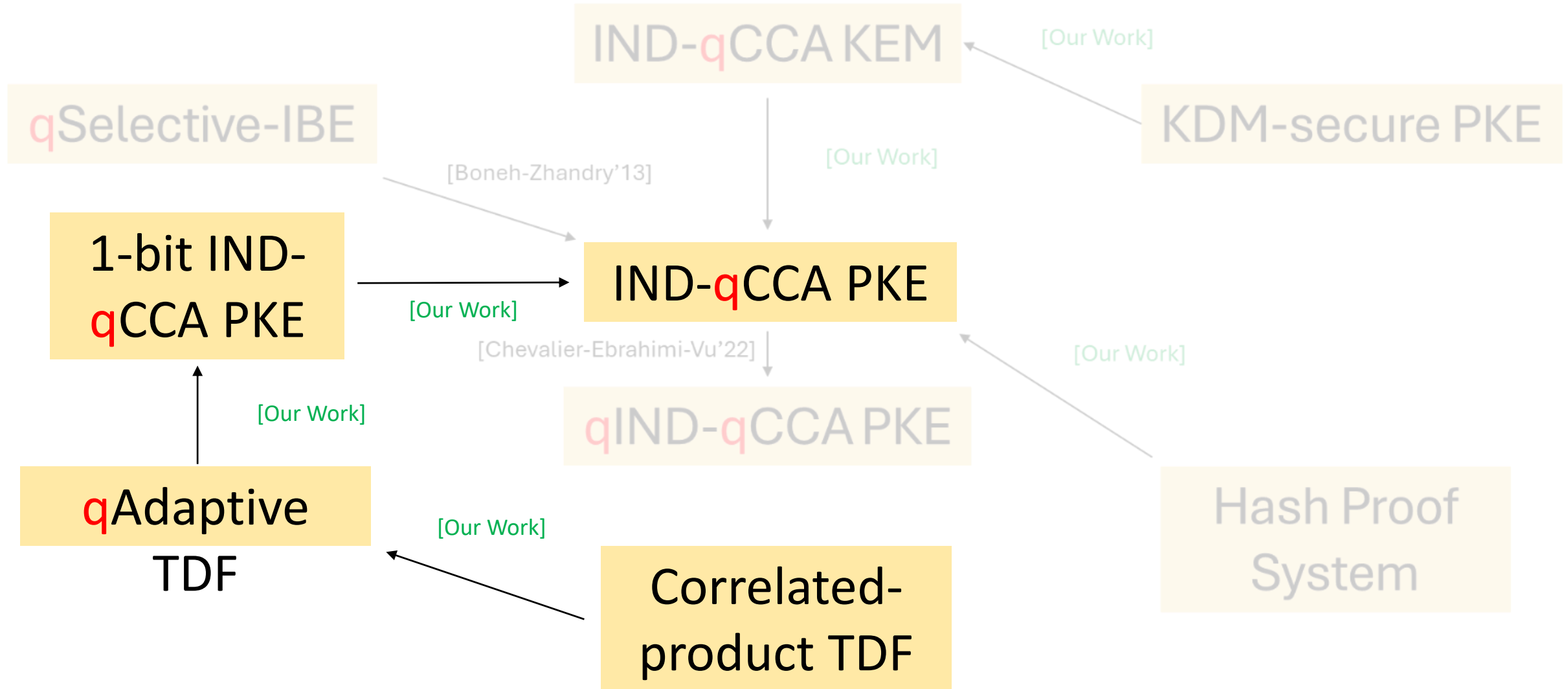
IND- q CCA- c KG Secure ABE



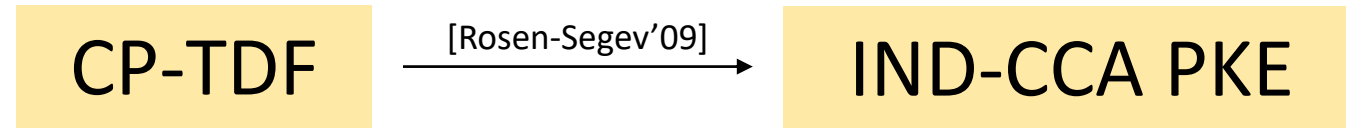
Overview: Techniques



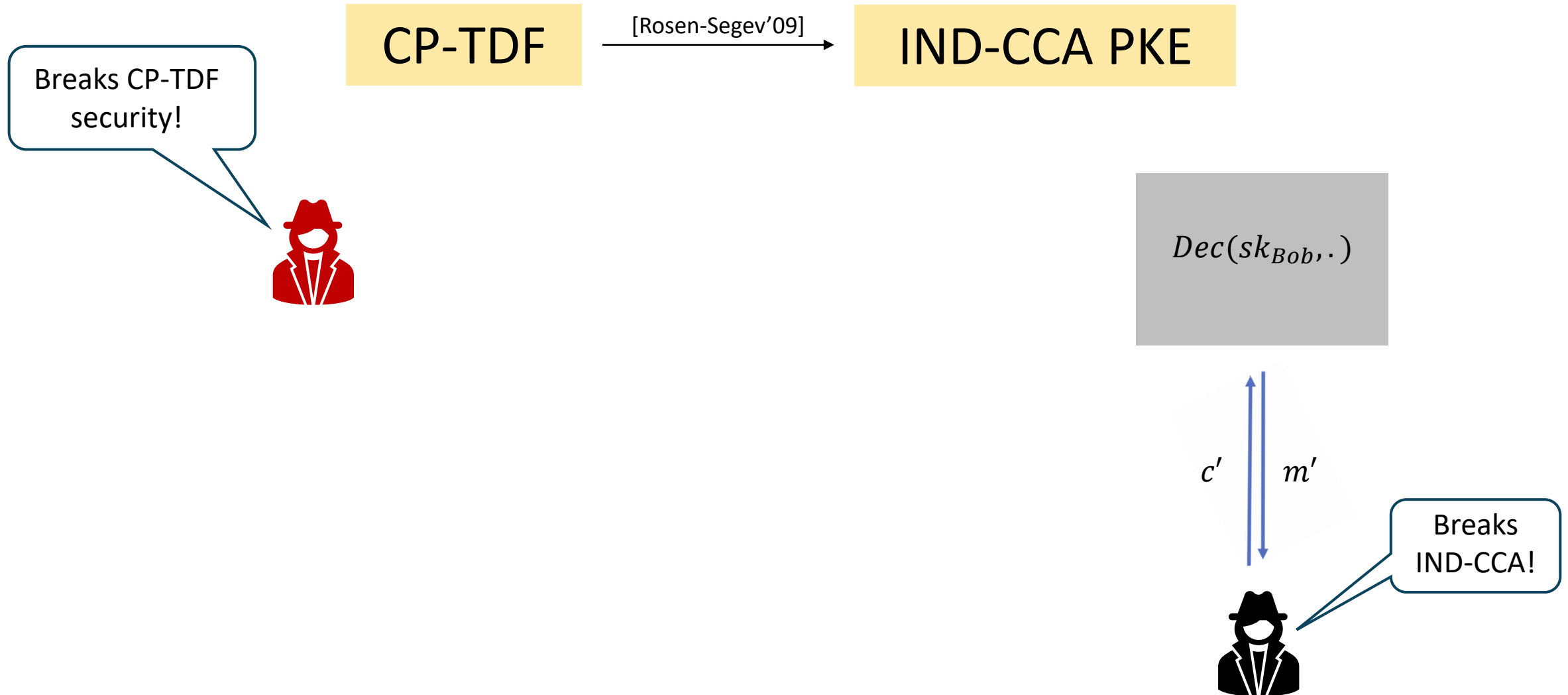
Overview: Techniques



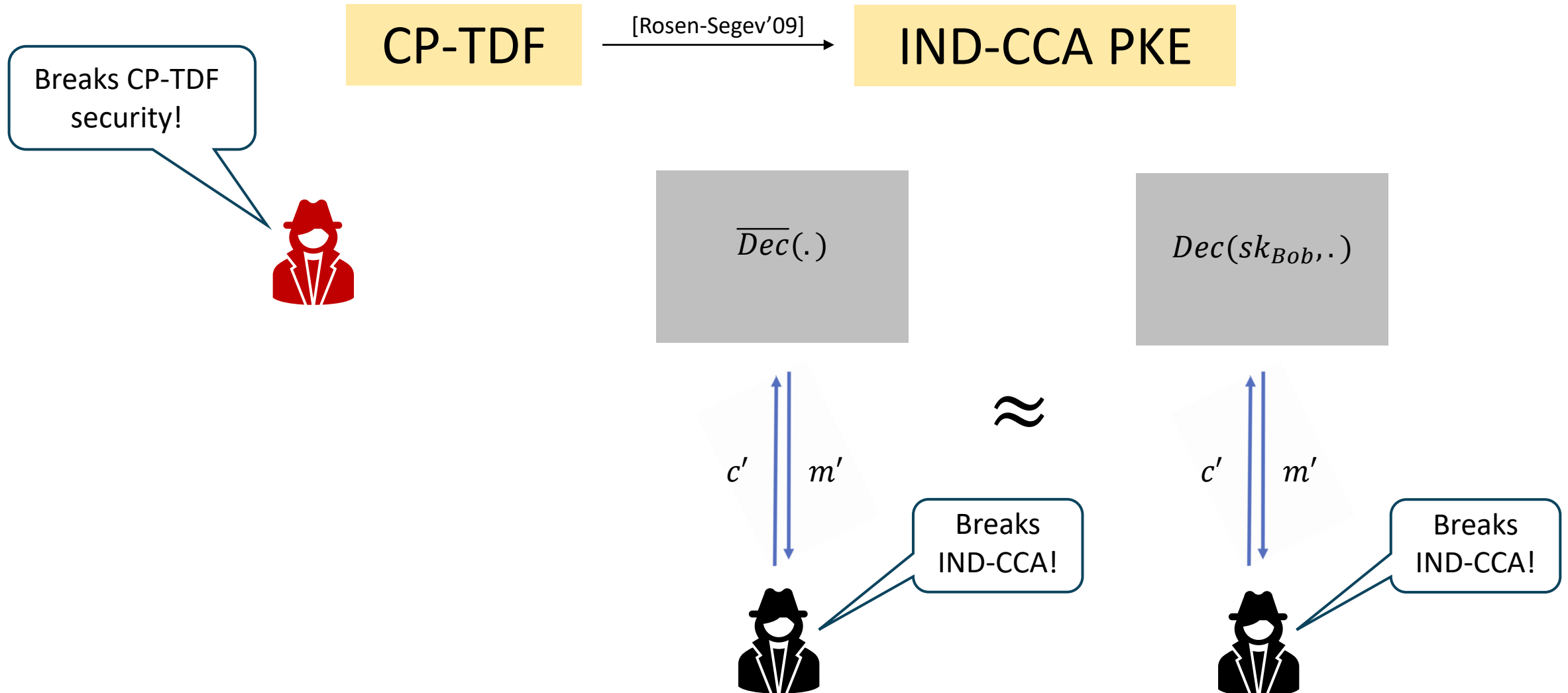
Overview: Techniques



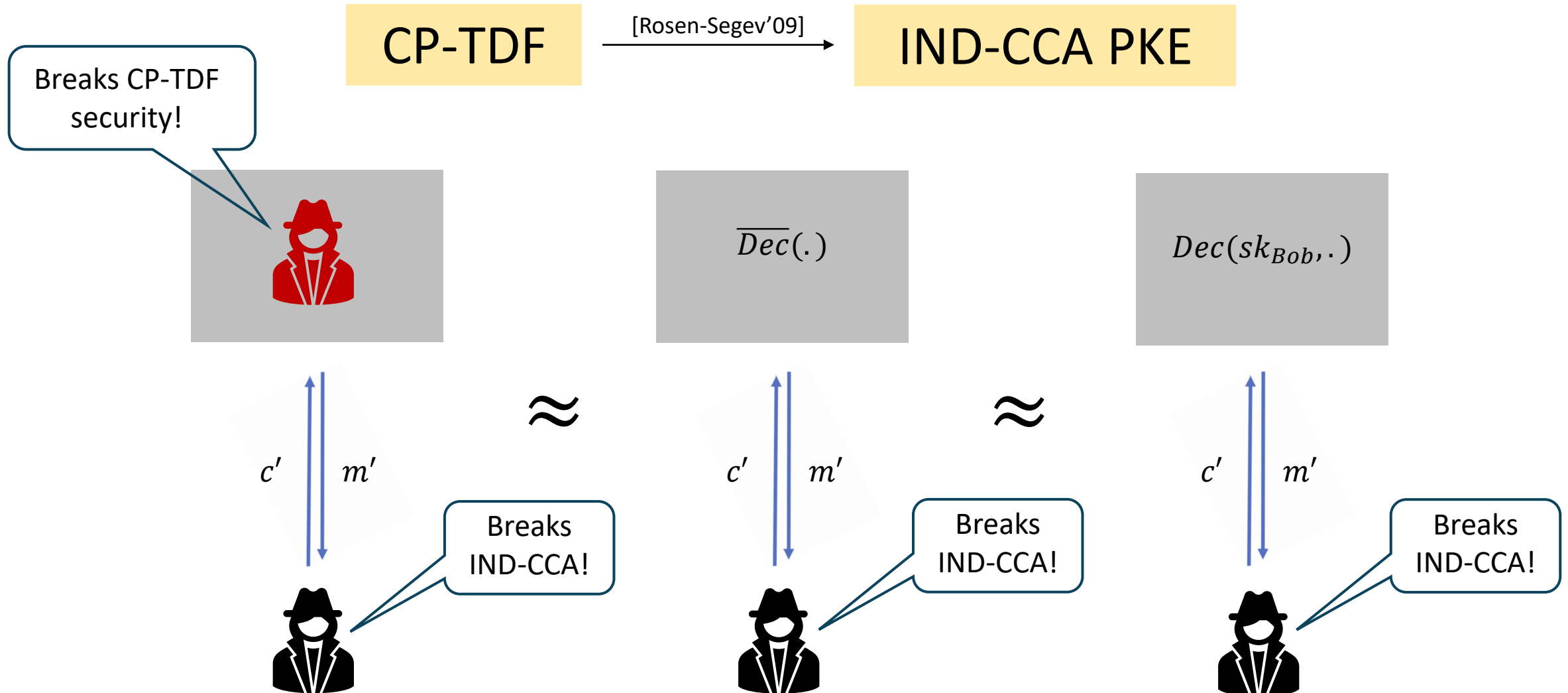
Overview: Techniques



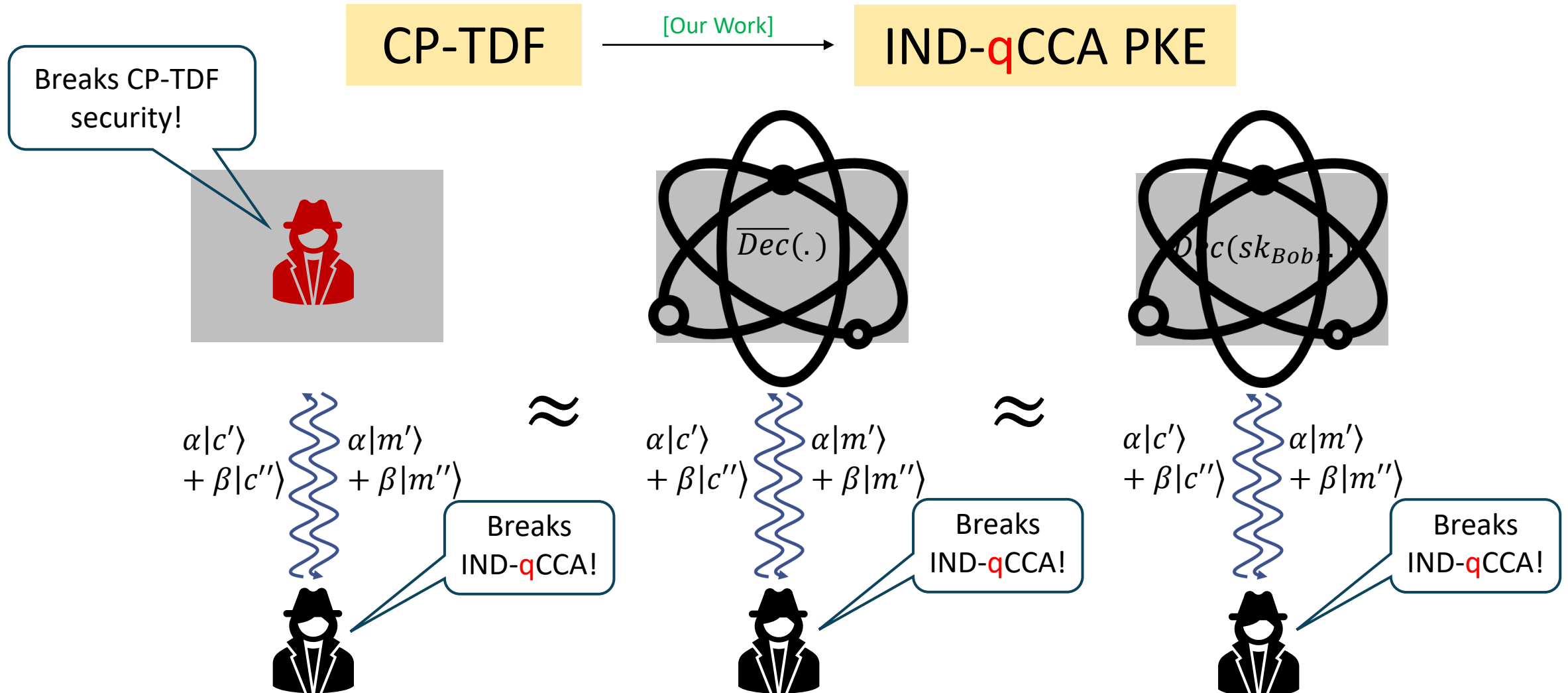
Overview: Techniques



Overview: Techniques



Overview: Techniques



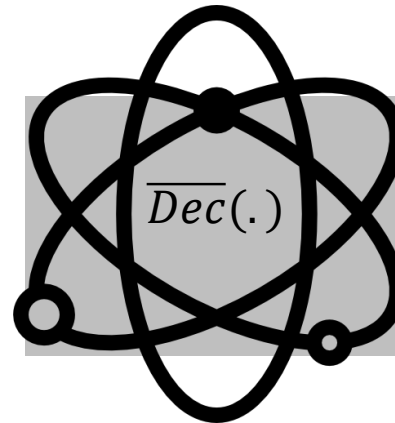
Overview: Techniques

CP-TDF

[Our Work]

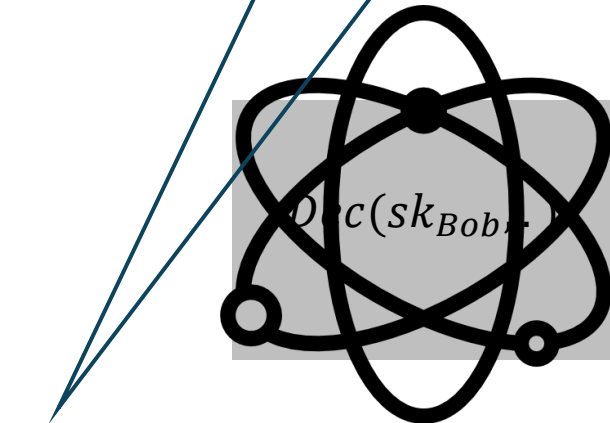
IND-**q**CCA PKE

Breaks CP-TDF
security!



$$\alpha|c'\rangle + \beta|c''\rangle \quad \alpha|m'\rangle + \beta|m''\rangle$$

Breaks
IND-**q**CCA!



$$\alpha|c'\rangle + \beta|c''\rangle \quad \alpha|m'\rangle + \beta|m''\rangle$$

Breaks
IND-**q**CCA!



Breaks
IND-**q**CCA!



Shown using the **generalized OW2H lemma** of [Ambainis-Hamburg-Unruh'19].

Overview: Techniques

Showed using the generalized OW2H lemma of [Ambainis-Hamburg-Unruh'19].

Breaks
security

- The original “**One-Way To Hiding**” (**OW2H**) lemma of [Unruh'14] was used to argue indistinguishability of **quantum (uniformly) random oracles**.

Breaks
OW-CCA!

Overview: Techniques

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Breaks
security

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- Our work involves the first application of the (generalized) OW2H

Breaks
D-CCA

Overview: Techniques

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Breaks
security

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- Our work involves the first application of the generalized OW2H lemma w.r.t. qCCA decryption oracles in the **standard model** – as opposed to the **QROM**.

Breaks
qCCA

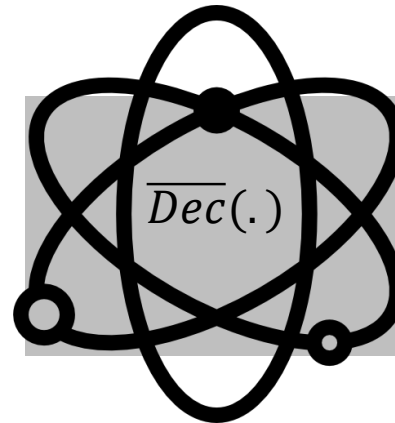
Overview: Techniques

CP-TDF

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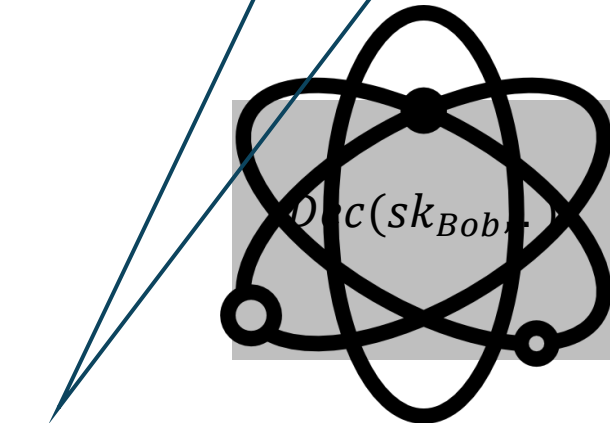
IND-**q**CCA PKE

Breaks CP-TDF
security!



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Breaks
IND-**q**CCA!



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Breaks
IND-**q**CCA!

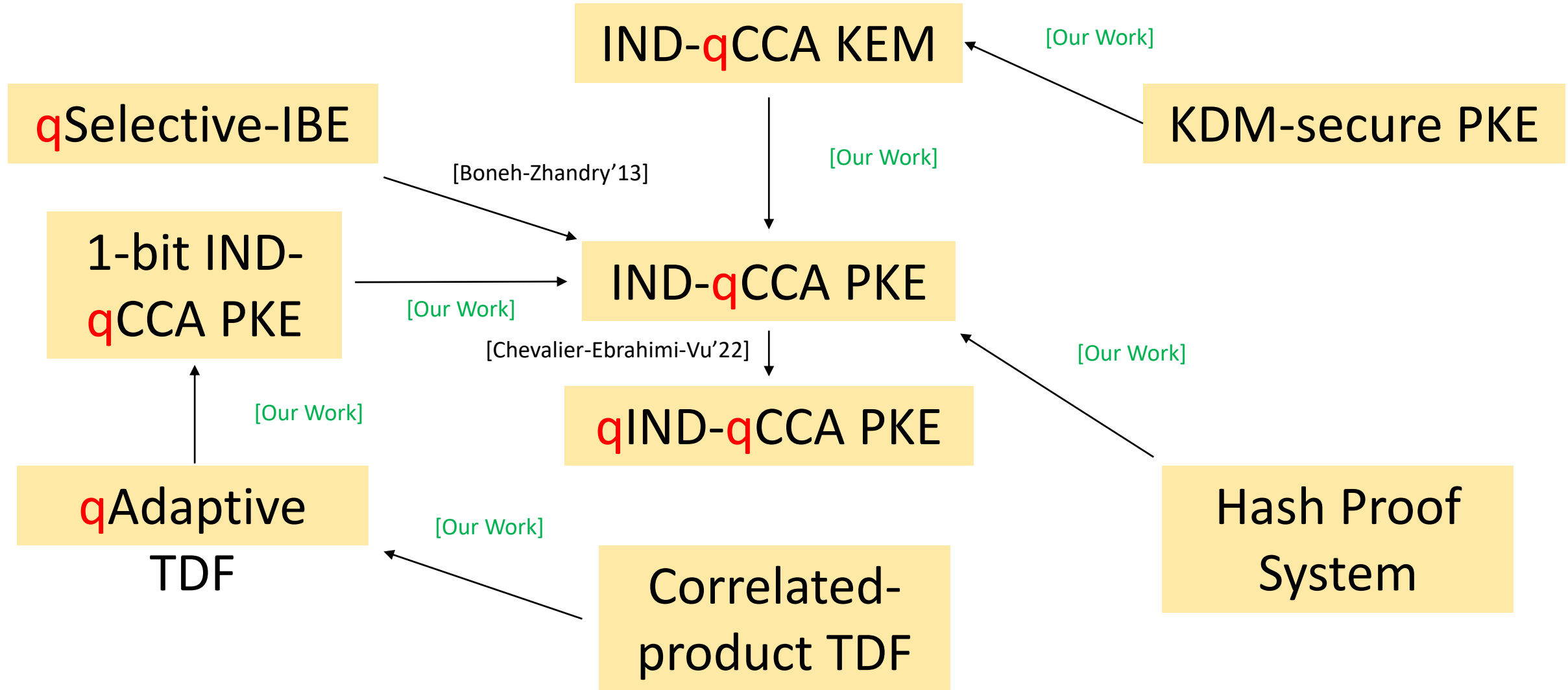


Breaks
IND-**q**CCA!



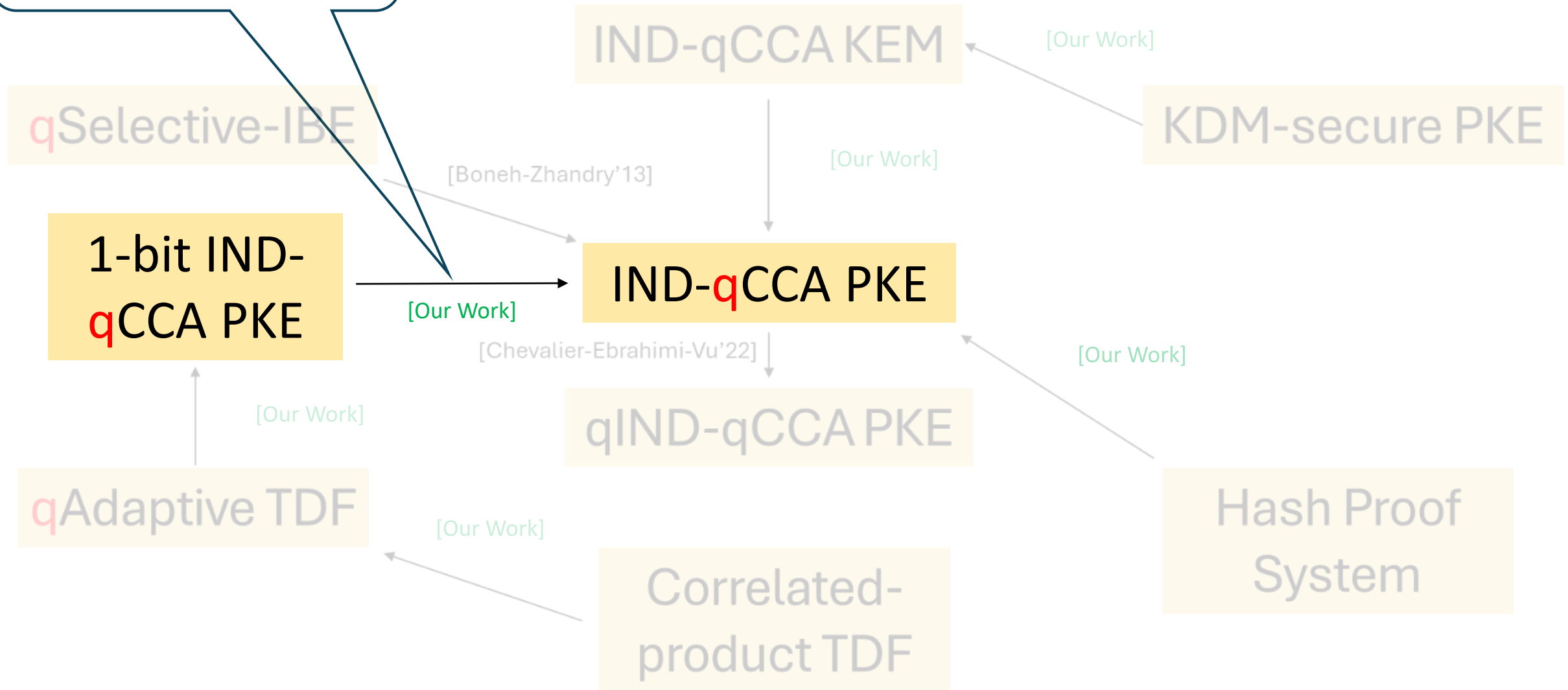
Shown using the **generalized OW2H lemma** of [Ambainis-Hamburg-Unruh'19].

Overview: Techniques



Required a “**nested**”
application of the generalized
OW2H lemma.

Overview: Techniques



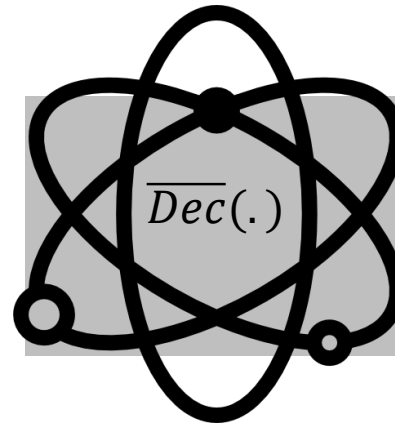
Overview: Techniques

CP-TDF

[Our Work]

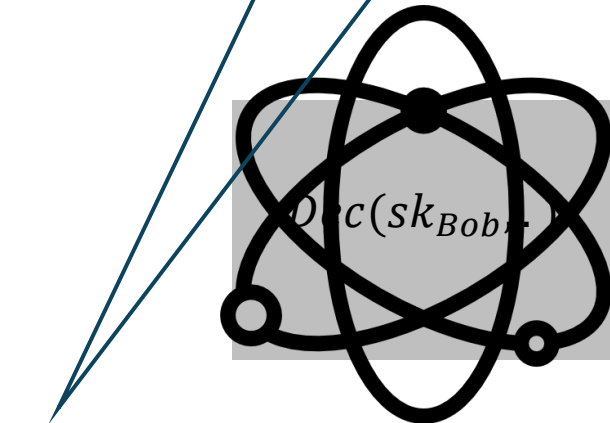
IND-**q**CCA PKE

Breaks CP-TDF
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Breaks
IND-**q**CCA!



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Breaks
IND-**q**CCA!



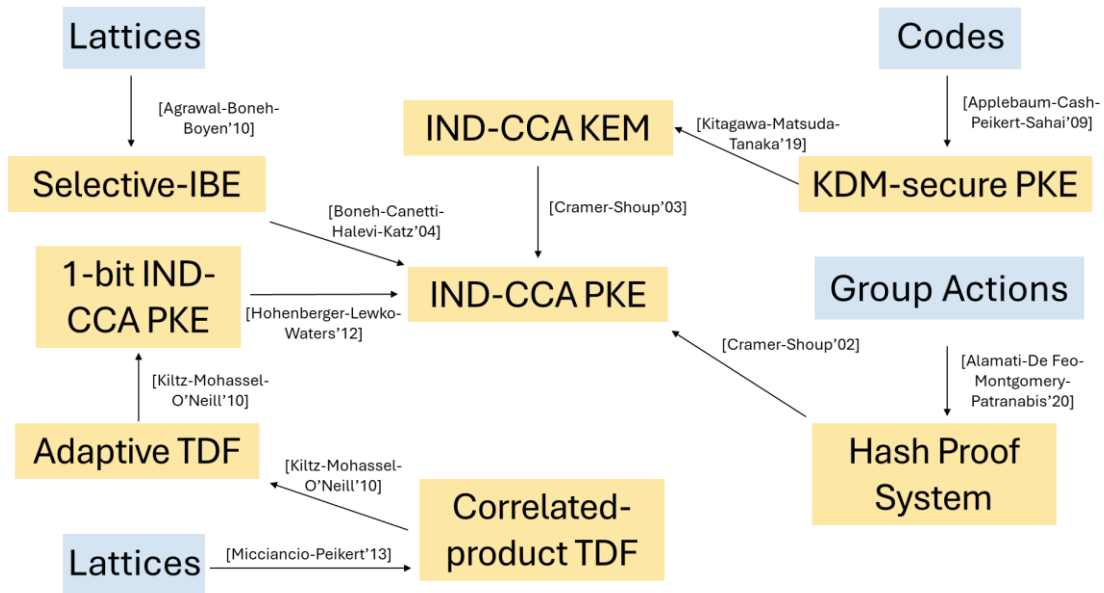
Breaks
IND-**q**CCA!



Shown using the **generalized OW2H lemma** of [Ambainis-Hamburg-Unruh'19].

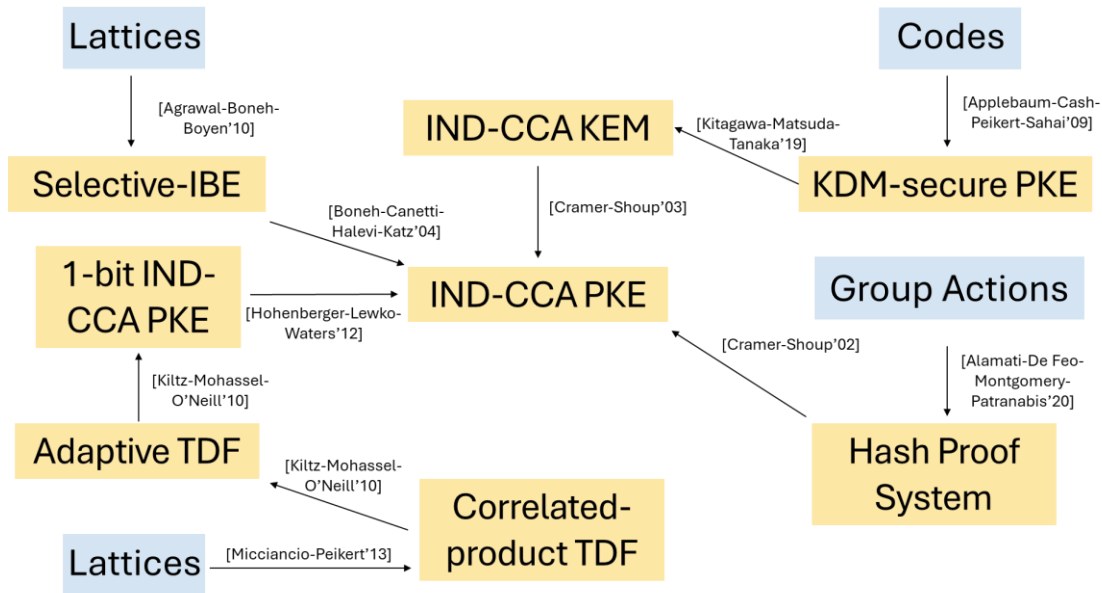
Conclusion

Conclusion

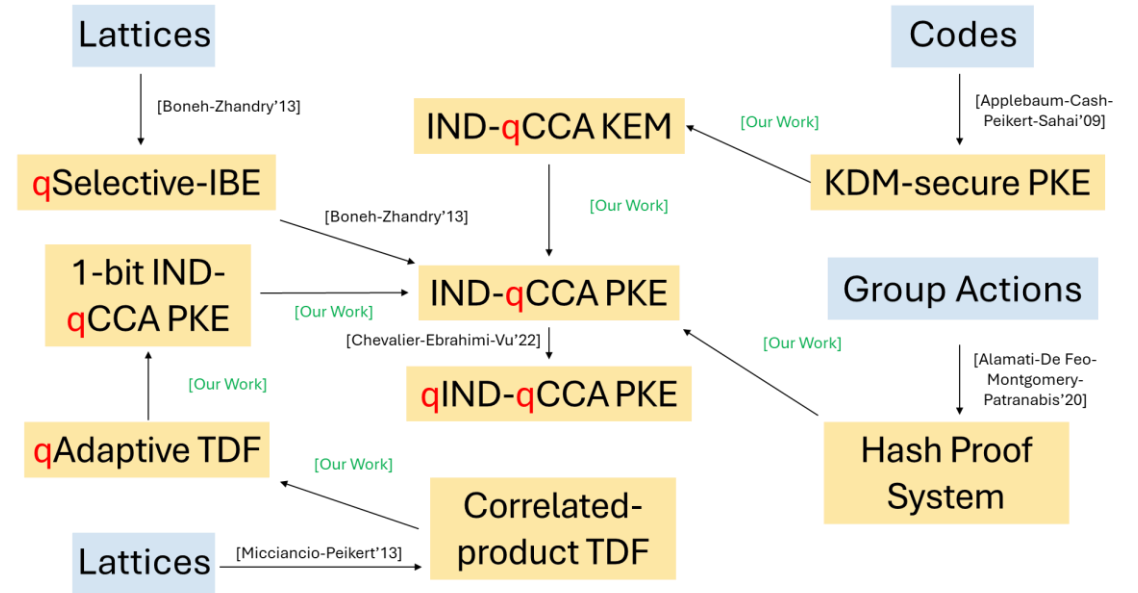


IND-CCA PKE

Conclusion



IND-CCA PKE



IND- q CCA PKE

Conclusion

