On Broadcast in Generalized Network and Adversarial Models

Varun Maram 7th December 2020

Joint work with Chen-Da Liu-Zhang and Ueli Maurer Department of Computer Science, ETH Zurich

Setting

- Parties: $P = \{P_1, P_2, ..., P_n\}$
- Synchronous
- No PKI setup



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Adversary

- Static
- Active (Byzantine)
- Unbounded



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Classical model

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Can we tolerate more?



3-minicast model

[FM00]:

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Trade-off b/w power of network and power of adversary?



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General: $A = \{A_1, A_2, \dots, A_k\} -$	A contains b-chain(s) and A is $(b + 1)$ -chain free	Some <i>b</i> -minicast channels	[LMM20]
$A_i \subseteq P$	$(A \in \mathfrak{A}^{(b)})$		

- 3-minicast model
- Six parties
- $t \leq 3$



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- Broadcast is impossible [CFF+05]



- 4-minicast model
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A is b-chain free \Rightarrow A is (b + 1)-chain free












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[LMM20]: For adversary structures $A \in \mathfrak{A}^{(b)}$, broadcast is achievable in general networks only if, for all k ($3 \le k \le b$):

- for every k-chain in A of the form $(S_1, S_2, ..., S_k)$,
- there is a k-minicast channel that has non-empty intersection with each of the sets S_1, S_2, \ldots, S_k .

Non-essential minicasts

- *n* parties
- adv. is (b + 1)-chain free
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<u>Idea</u>: Simulate *missing/non-essential b*-minicast channels with *local* executions of [Ray15]'s protocol.


- b parties
- adv. is *b*-chain free
- (b 1)-minicast model (complete)
- Broadcast is possible [Ray15]



- b parties
- adv. is *b*-chain free (locally restricted)
- (b 1)-minicast model (complete)
- Broadcast is possible [Ray15]















- b parties
- The projected adv. is *b*-chain free
- (b 1)-minicast model (complete)
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- there is a complete set of bilateral channels, and
- for each subset of parties ρ of size k ($3 \le k \le b$):
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Condition <u>non-trivial</u> for certain weak class of adversaries in $\mathfrak{A}^{(b)}$, namely <u>b</u>-chain adversaries.



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A *b*-chain adversary just contains a (single) *b*-chain, and nothing more

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 S_1 S_2 Parties: $P = \{P_1, P_2, ..., P_n\}$ • General: $A = \{A_1, A_2, ..., A_k\},\$ ٠ $A_i \subseteq P$ S_6 S_3 Partition: $S = (S_1, S_2, \dots, S_b)$ • • $\bigcup_{i=1}^{b} S_i = P$ A_1 • $S_i \cap S_j = \emptyset$ • non-empty S_i 's *b*-chain adversary corresponding to *S*: • S_5 S_4 $A^{S} = \{P \setminus (S_{i} \cup S_{i+1}) \mid 1 \leq i \leq b\}$

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 - it belongs to the class $\mathfrak{A}^{(b)}$ (i.e., is also (b + 1)-chain free). [LMM20]
 - there exist subsets of parties ρ ($|\rho| = b$) such that $A^{S}[\rho]$ is *b*-chain free (i.e., *b*-minicast channel among ρ is *non-essential*). [LMM20]

Other results

• Our conditions allow us to derive bounds on the no. of *b*-minicast channels that are necessary and that suffice in achieving global broadcast in general networks secure against general adversaries.

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- Our conditions allow us to derive bounds on the no. of *b*-minicast channels that are necessary and that suffice in achieving global broadcast in general networks secure against general adversaries.
 - Thereby providing a way to extend [JMS12]'s quantitative analysis in general 3-minicast networks to higher *b*-minicast networks.

Open problems

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 - We showed that a straightforward extension of a technique (so-called virtual party emulation) used by [RVS⁺04] in deriving such tight conditions in general 3-minicast networks does not generalize to higher b-minicast networks.

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- Providing tighter necessary and sufficient conditions on general networks for achieving broadcast while tolerating general adversaries.
 - We showed that a straightforward extension of a technique (so-called virtual party emulation) used by [RVS⁺04] in deriving such tight conditions in general 3-minicast networks does not generalize to higher b-minicast networks.
- Implications of such results on broadcast in general *b*-minicast networks, secure against general adversaries, in a realistic setting.