

Post-quantum Cryptographic Analysis of SSH



SANDBOXAQ

Varun Maram
Cybersecurity Group
SandboxAQ



: <https://varun-maram.github.io/>



: varun-maram-pqc

Joint work with Benjamin Benčina, Benjamin Dowling, and Keita Xagawa



Why Analyze SSH?

Why Analyze SSH?

Top 10 real-world cryptographic protocols.

Here are ten widely used real-world cryptographic protocols that help secure communication, authentication, and data integrity across various applications:

1. Transport Layer Security (TLS) / Secure Sockets Layer (SSL)

- **Use Case:** Secure web browsing (HTTPS), email encryption, VoIP, and more.
- **Description:** TLS (successor to SSL) encrypts data transmission between clients and servers to ensure confidentiality and integrity.

2. Secure Shell (SSH)

- **Use Case:** Secure remote access to servers and encrypted file transfers (SCP, SFTP).
- **Description:** Uses public-key cryptography and symmetric encryption to secure login sessions over an unsecured network.

TLS @ RWC (≥ 27 talks)

- *A long answer to the simple question, "Is **TLS** provably secure?"* [2012]
- *On the Security of the **TLS** and DTLS Record Protocols* [2013]
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- *Provable security of advanced properties of **TLS** and SSH* [2014]
- *Post-quantum key exchange for the **TLS** protocol from the ring learning with errors problem* [2015]
- **TLS** 1.3 [2015]
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- *The OPTLS Protocol and **TLS** 1.3* [2016]
- *Automated Verification of **TLS** 1.3: 0-RTT, Resumption and Delayed Authentication* [2016]
- *PRNG Failures and **TLS** Vulnerabilities in the Wild* [2017]
- *Concerto: A Methodology Towards Reproducible Analyses of **TLS** Datasets* [2017]
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- *The Era of **TLS** 1.3: Measuring Deployment and Use with Active and Passive Methods* [2020]
- *The 9 Lives of Bleichenbacher's CAT: New Cache Attacks on **TLS** Implementations* [2020]
- *Deco: Liberating Web Data Using Decentralized Oracles for **TLS*** [2020]
- *Raccoon Attack: Finding and Exploiting Most-Significant-Bit-Oracles in **TLS**-DH(E)* [2021]
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- *TLS-Anvil: Adapting Combinatorial Testing for **TLS** Libraries* [2023]

Analyze SSH?

Log in

Sign up

Top 10 real-world cryptographic protocols.

protocols that help secure communication, applications:

The Sockets Layer (SSL)

Encryption, VoIP, and more.

transmission between clients and servers to

encrypted file transfers (SCP, SFTP).

symmetric encryption to secure login sessions over

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SSH @ RWC (2 talks)

- *Provable security of advanced properties of TLS and **SSH*** [2014]
- *Terrapin Attack: Breaking **SSH** Channel Integrity By Sequence Number Manipulation* [2024]

Sign up

Why Analyze SSH?



2. Secure Shell (SSH)

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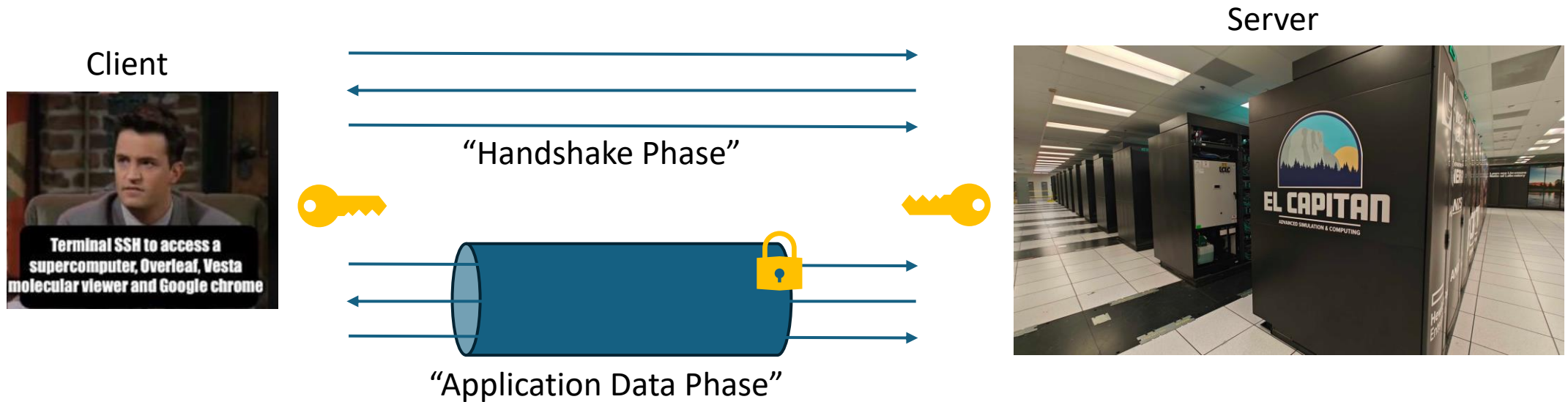
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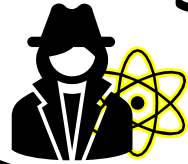
SSH Protocol



2. Secure Shell (SSH)

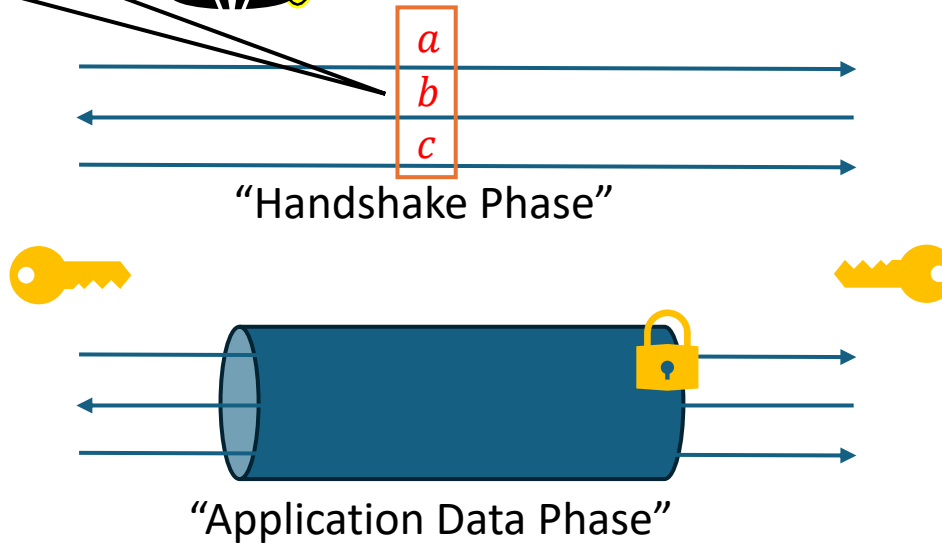
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Diffie-Hellman key-exchange,
not secure against quantum
attackers.



SSH Protocol

Client



Server

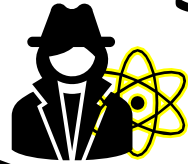


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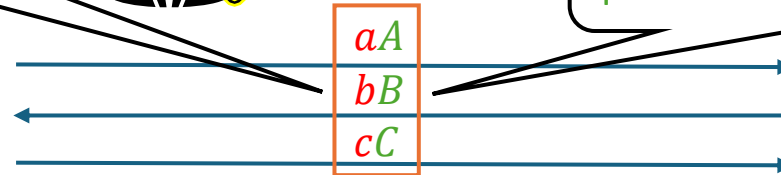
SSH Protocol

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“Hybridize” with (plausibly)
quantum-secure KEM.

Client



“Handshake Phase”



“Application Data Phase”



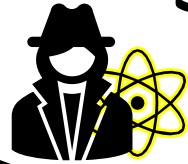
Server



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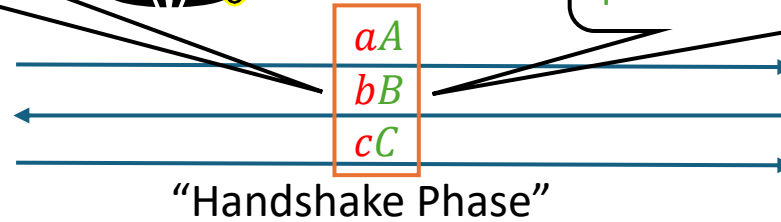


SSH Protocol

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Client

Server



History of PQC in SSH

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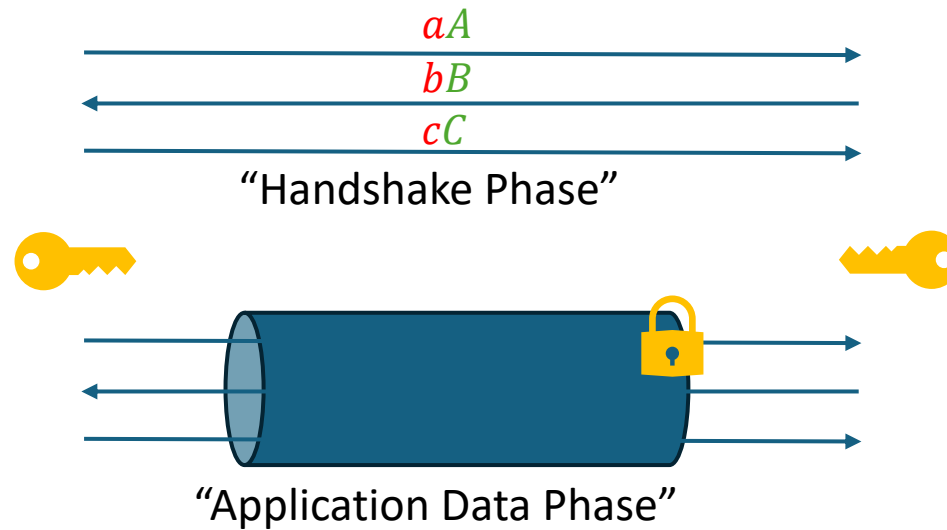
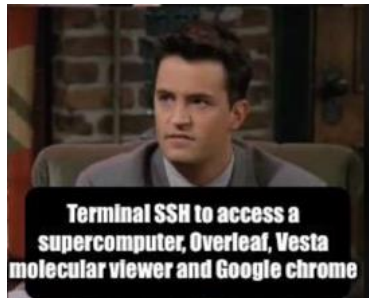
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- Description:
an unsecu

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Prior Analyses

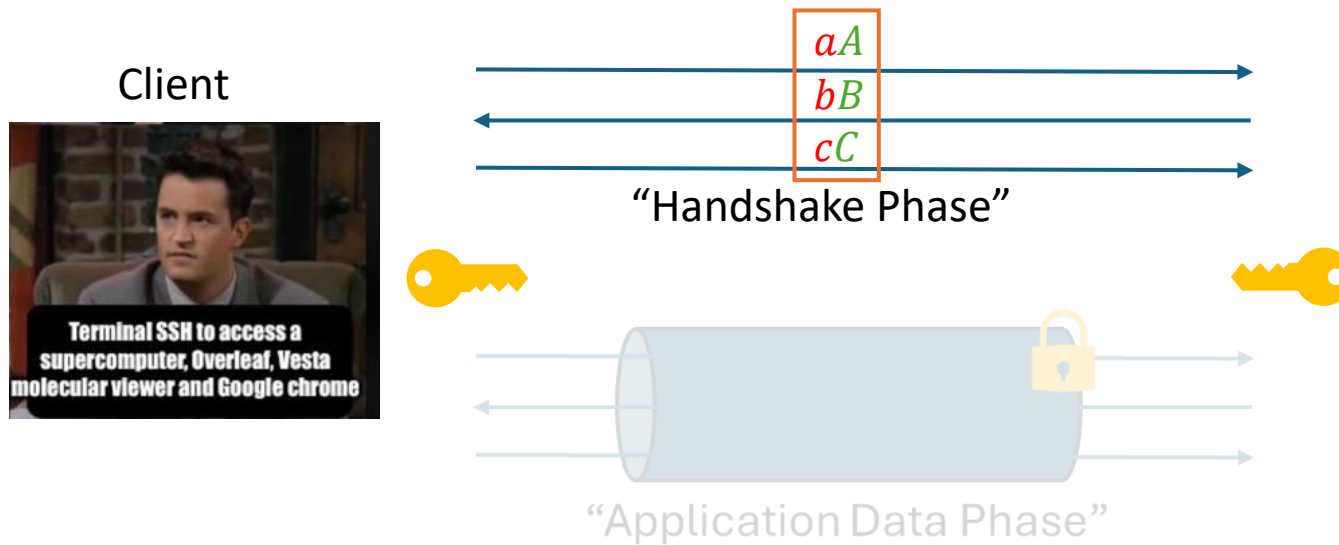
Client



Server



Prior Analyses



Server



Security of Hybrid Key Establishment using Concatenation

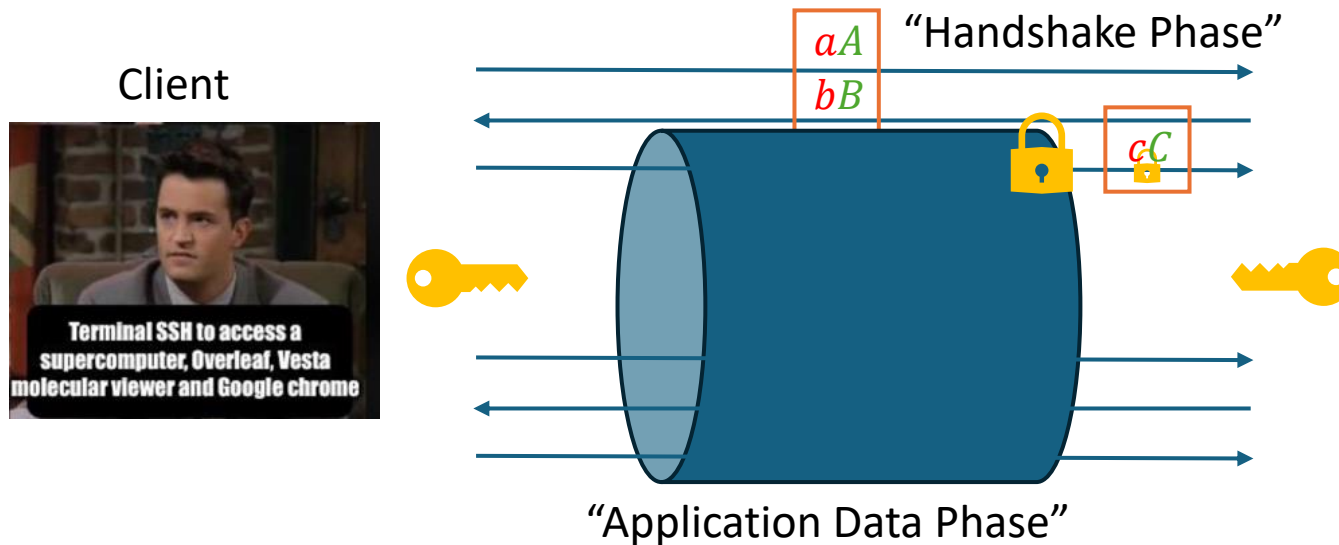
Adam Petcher and Matthew Campagna

Amazon Web Services

[ePrint '23]

- Prior works either analyze the hybrid key exchange **in isolation**, or...

Prior Analyses



Security of Hybrid Key Establishment using Post-quantum sound CRYPTOVERIF and verification of hybrid TLS and SSH key-exchanges

Bruno Blanchet
Inria, F-75012 Paris, France
bruno.blanchet@inria.fr

Charlie Jacomme
Université de Lorraine, CNRS, Inria, LORIA, F-54000 Nancy, France
charlie.jacomme@inria.fr

[CSF '24]

- Prior works either analyze the hybrid key exchange **in isolation**, or...
- ... analyze post-quantum SSH in **unsuitable protocol models**, such as “authenticated key exchange (AKE)”.

Our Contributions

Our Contributions

#1: “Post-quantum SSH is cryptographically secure in practice.”

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(*Analysis does not cover **low-level implementation details.**)

Our Contributions

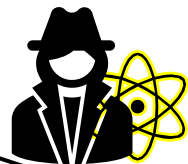
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cryptographic
practice.”*



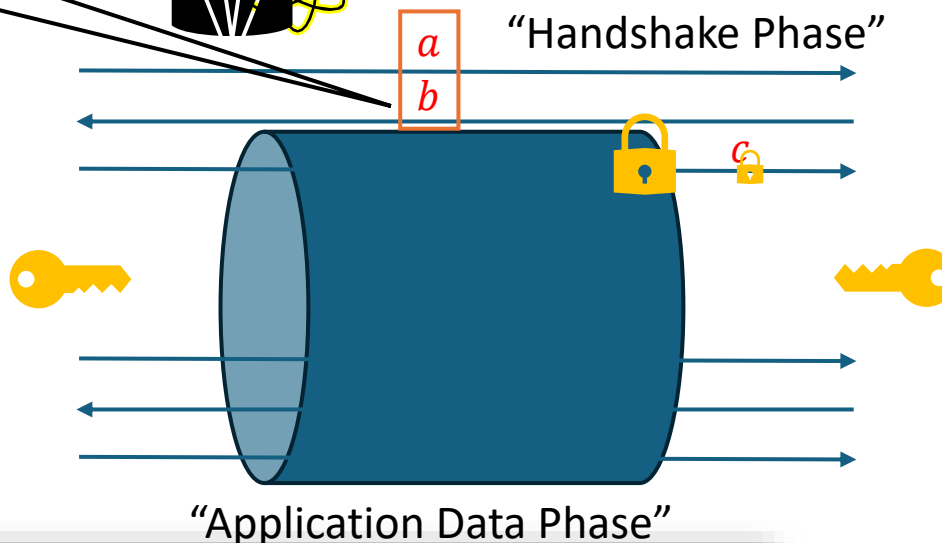
(*Analysis does not cover **low-level implementation details.**)

Diffie-Hellman key-exchange,
not secure against quantum
attackers.

Our Contributions



Client



Server



Multi-ciphersuite security of the Secure Shell (SSH) protocol

Florian Bergsma¹ Benjamin Dowling^{2a} Florian Kohlar¹ Jörg Schwenk¹
Douglas Stebila^{2a,2b}

¹ Horst Görtz Institute, Ruhr-Universität Bochum, Bochum, Germany
{florian.bergsma,florian.kohlar,joerg.schwenk}@rub.de

^{2a} School of Electrical Engineering and Computer Science

^{2b} School of Mathematical Sciences

^{2a,2b} Queensland University of Technology, Brisbane, Australia

{b1.dowling,stebila}@qut.edu.au

[CCS '14]

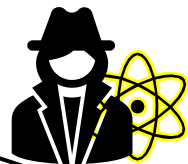
- [BDKSS14] proves security of SSH in the “authenticated and confidential channel establishment (ACCE)” model.
- However, analysis is in the **classical Diffie-Hellman setting**.

[BDKSS14]: F. Bergsma, B. Dowling, F. Kohlar, J. Schwenk, D. Stebila, “Multi-ciphersuite Security of the Secure Shell (SSH) Protocol”, CCS 2014

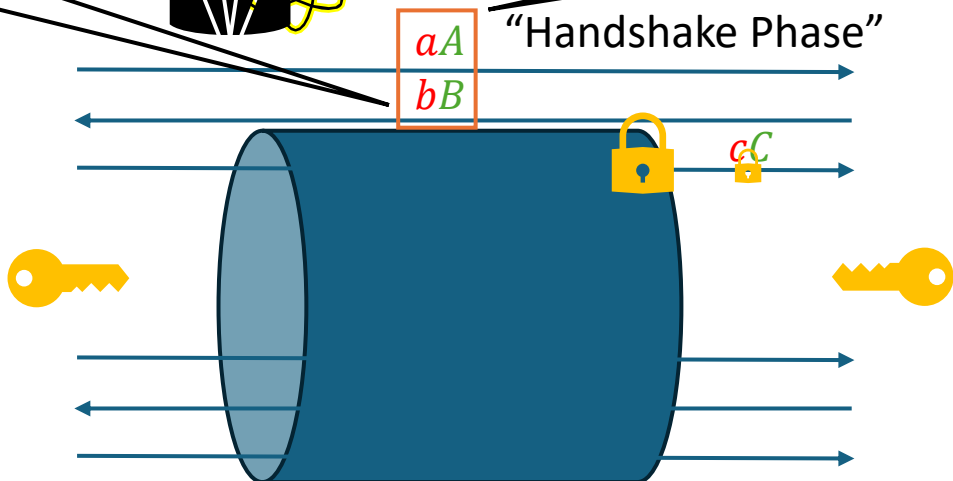
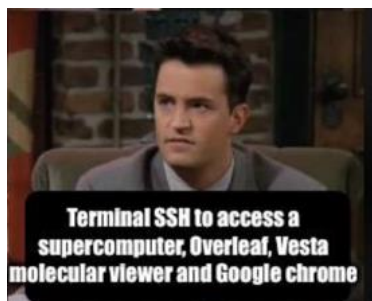
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Our Contributions

“Hybridize” with (plausibly)
quantum-secure KEM.



Client



Server



Multi-ciphersuite security

Post-quantum Cryptographic Analysis of SSH

Benjamin Benčina
Royal Holloway, University of London, UK
Email: benjamin.bencina.2022@live.rhul.ac.uk

Benjamin Dowling
King's College London, UK
Email: benjamin.dowling@kcl.ac.uk

Varun Maram
SandboxAQ, UK
Email: varun.maram@sandboxaq.com

Keita Xagawa
Technology Innovation Institute, UAE
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[S&P '25]

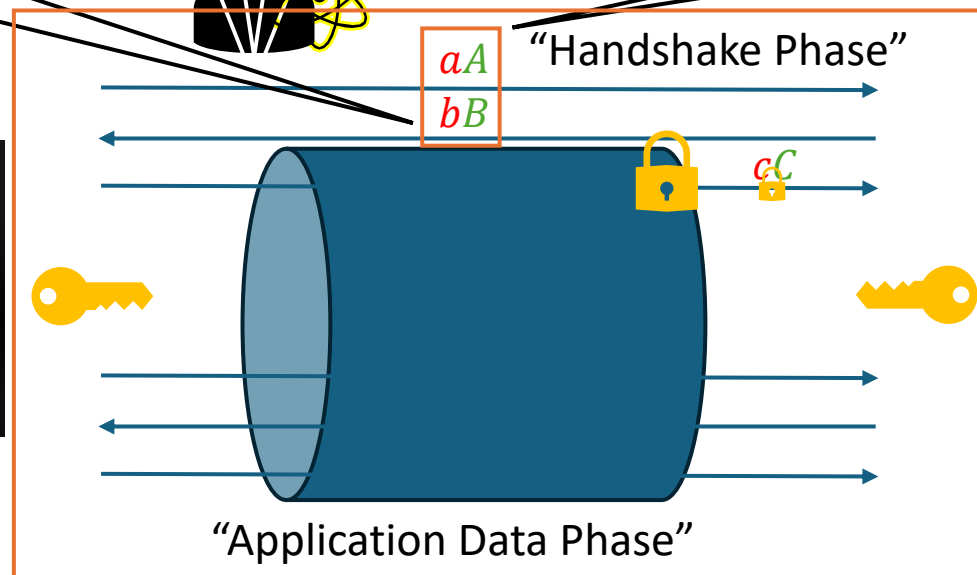
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- Analysis also captures **forward secrecy** (unlike the classical ACCE analysis in [BDKSS14]).

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“PQ ACCE Security”

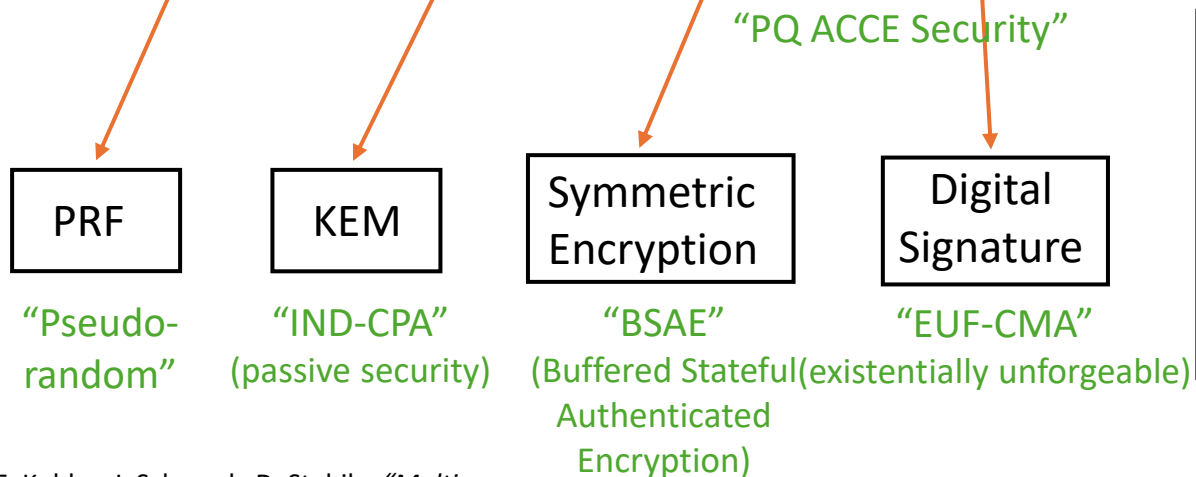
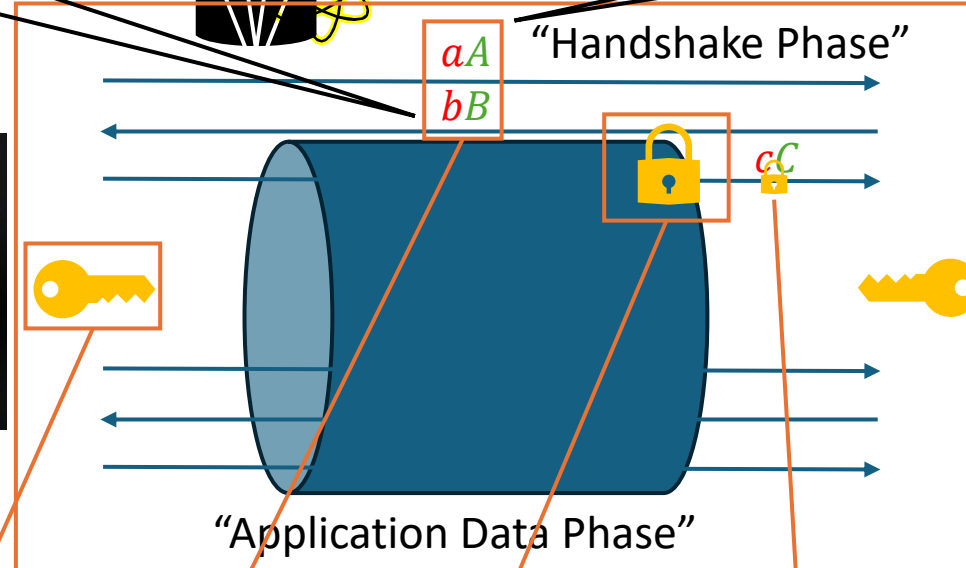


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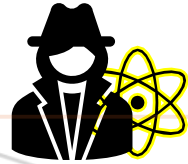


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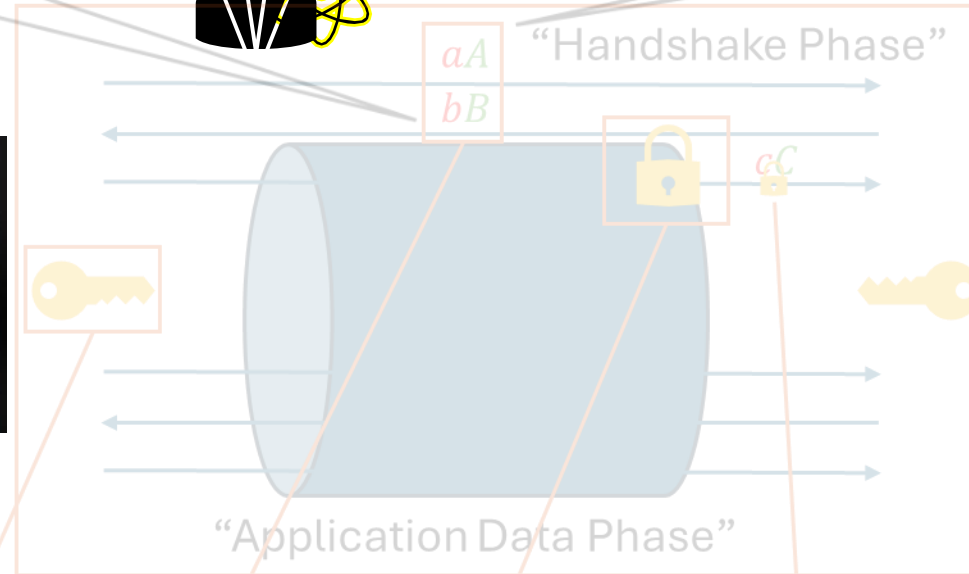
Our Contributions

Diffie-Hellman key-exchange, **not secure** against quantum attackers.

“Hybridize” with (plausibly) **quantum-secure** KEM.



Client



Server



PRF

“Pseudo-random”

KEM

“IND-CPA”
(passive security)

Symmetric
Encryption

“BSAE”
(Buffered Stateful Authenticated Encryption)

Digital
Signature

“EUF-CMA”

- We establish security of SSH in a **PQ-extension of ACCE model** that accounts for “harvest now, decrypt later” attacks.
- Analysis also captures **forward secrecy** (unlike the classical ACCE analysis in [BDKSS14]).
- We then prove corresponding PQ security properties of **SSH primitives**.

Our Contributions

#2: “Post-quantum SSH
can be more efficient, and
eco(log|nom)ical.”

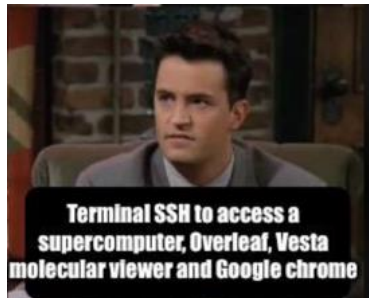
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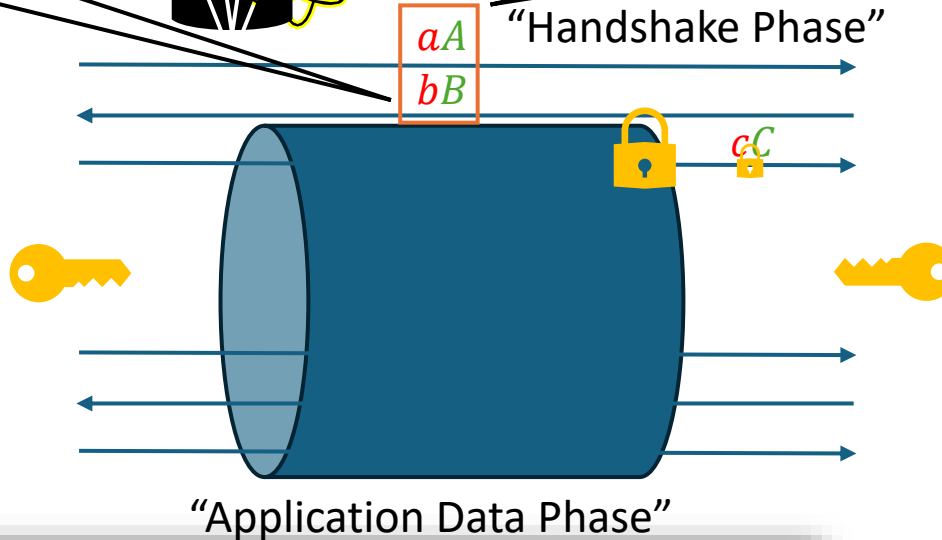
“Hybridize” with (plausibly)
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Client



Terminal SSH to access a
supercomputer, Overleaf, Vesta
molecular viewer and Google chrome



Server



Post-quantum sound CryptoVerif and verification of
hybrid TLS and SSH key-exchanges

Theorem 5 (Post-quantum SSH security, simplified)

Under the post-quantum **IND-CCA2** assumption for the KEM, the post-quantum PRF assumption for dPRF and PRF_c, the CR assumption for h , the classical EUF-CMA assumption for signatures, the hybrid SSH key exchange ensures **forward secrecy** even against quantum attackers, provided quantum attackers do not exist yet when the handshake is performed.

- PQ-SSH analysis of [BlaJac24] relies on **IND-CCA** secure (ephemeral) KEMs.

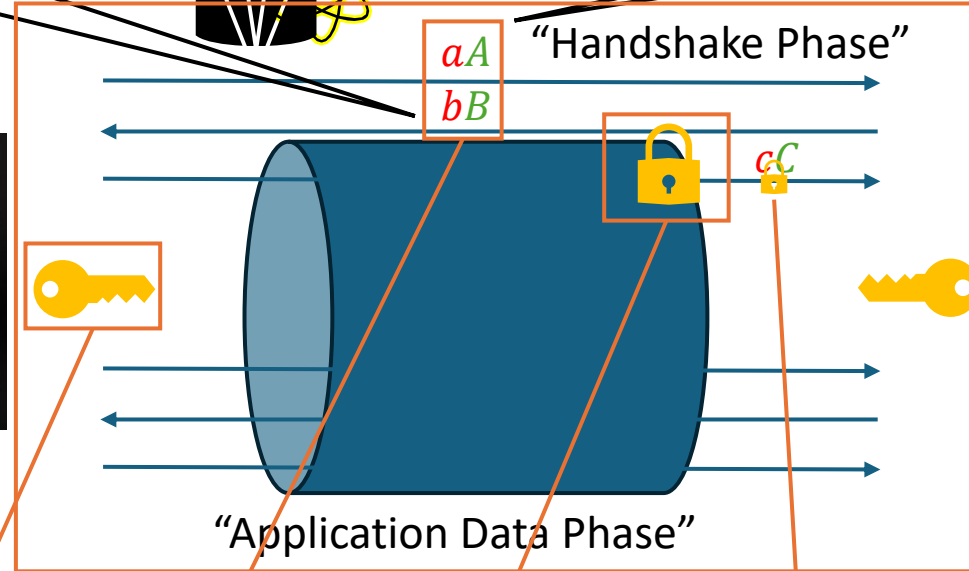
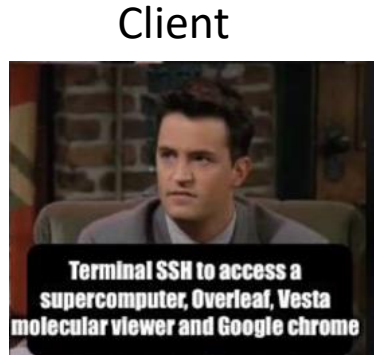


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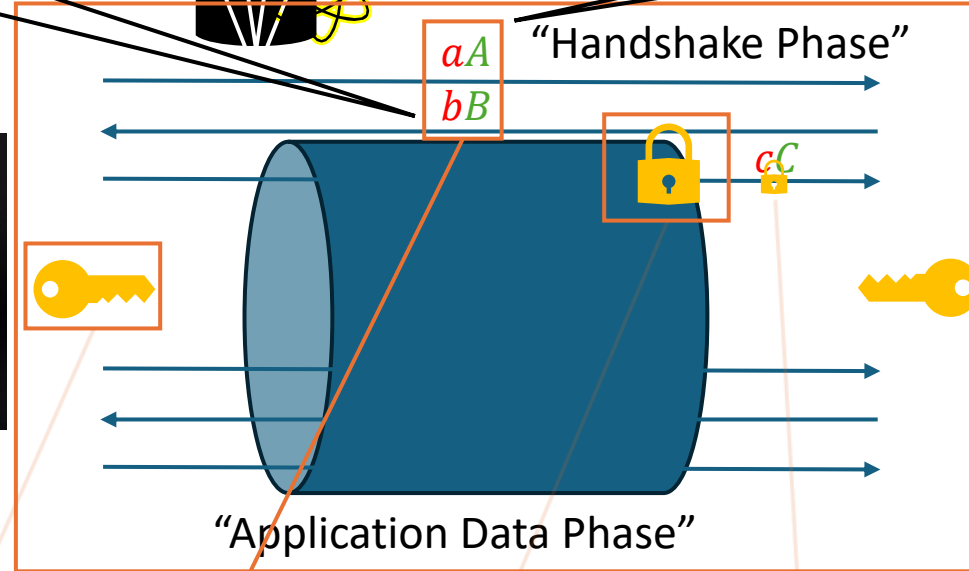


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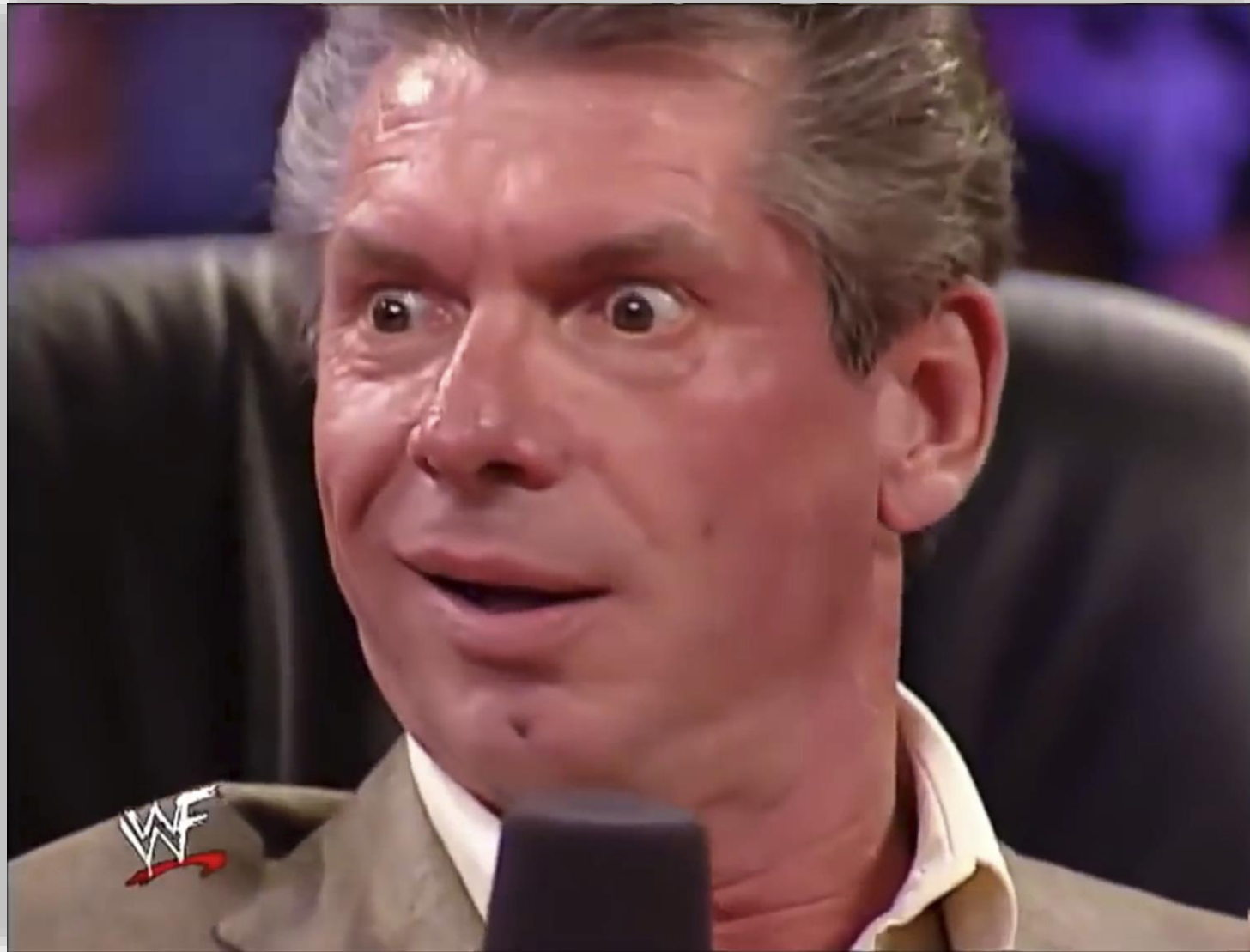
Symmetric
Encryption

“BSAE”
(Buffered Stateful
Authenticated
Encryption)

Digital
Signature

“EUF-CMA”
(existentially
unforgeable)

- PQ-SSH analysis of [BlaJac24] relies on **IND-CCA** secure (ephemeral) KEMs.
- Whereas our analysis relies on the weaker property of **IND-CPA** security.



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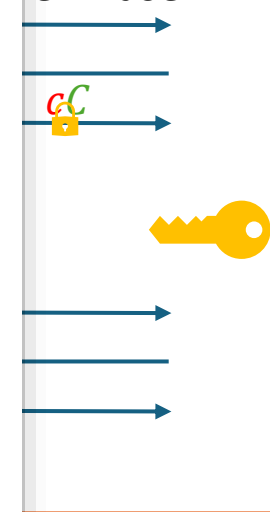
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History of PQC in SSH

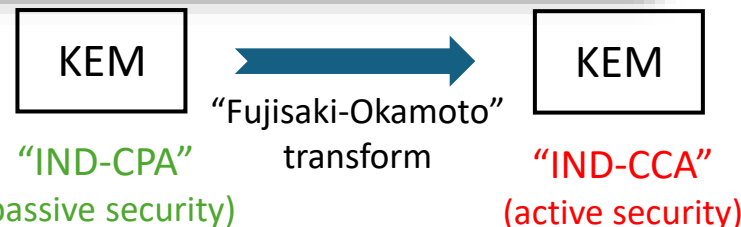
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ake Phase”



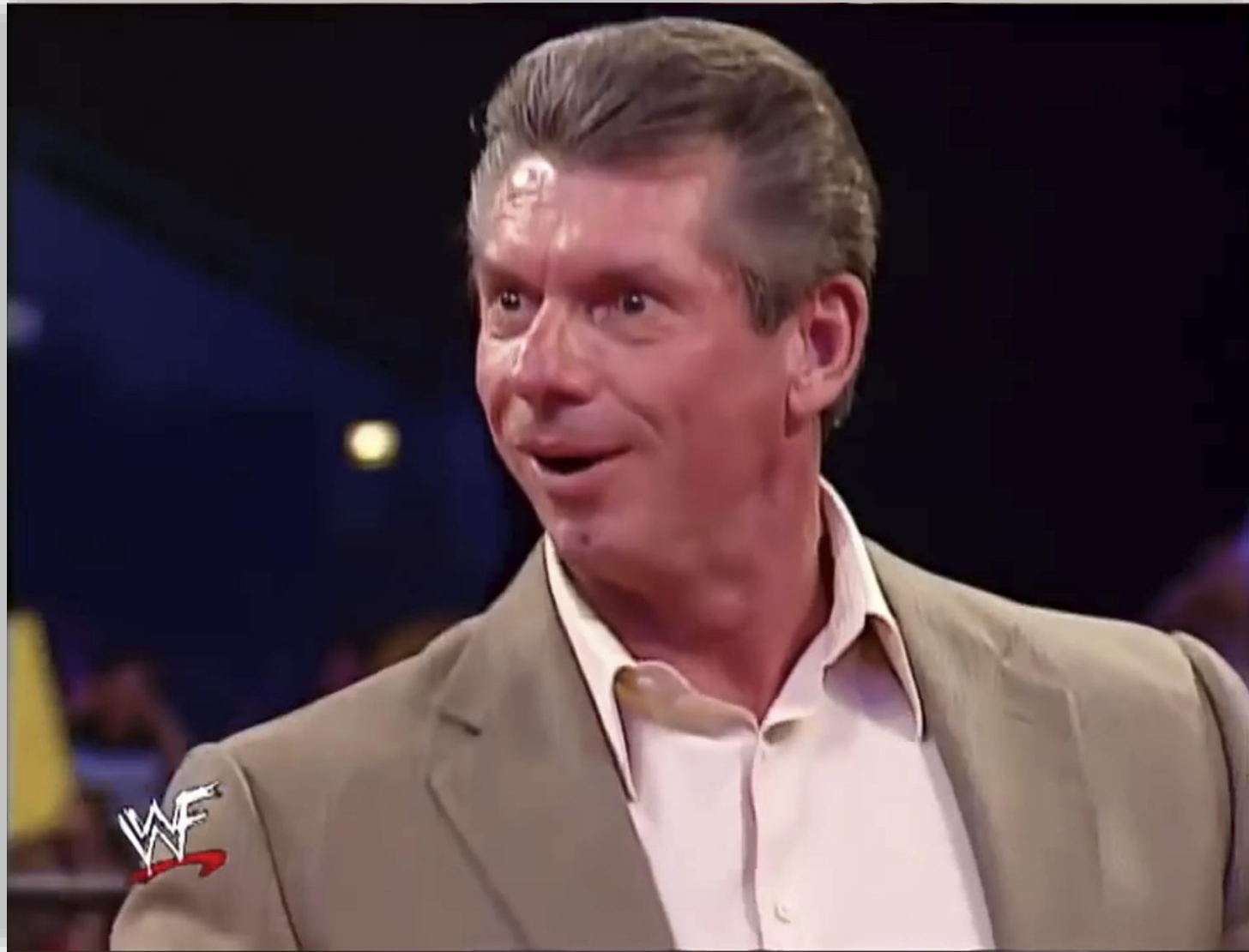
Security”



Server



- PQ-SSH analysis of [BlaJac24] relies on **IND-CCA** secure (ephemeral) KEMs.
- Whereas our analysis relies on the weaker property of **IND-CPA** security.
- Our analysis suggests FO transform is not needed, which in turn can lead to **performance improvements** in PQ-SSH.



(*honest reaction)

Our Contributions

Table 1. IND-CCA vs IND-CPA benchmarks w.r.t. ephemeral KEMs in post-quantum SSH.

Scope	KEM	IND-CCA[s]	IND-CPA[s]	Speedup[%]
Primitive-level (CPU timings of all KEM operations)	sntrup761	$2.8164 \cdot 10^{-2}$	$2.7056 \cdot 10^{-2}$	3.93
	mlkem768	$2.7853 \cdot 10^{-5}$	$1.3242 \cdot 10^{-5}$	52.46
	sntrup761x25519-sha512	$3.0290 \cdot 10^{-2}$	$2.9105 \cdot 10^{-2}$	3.91
	mlkem768x25519-sha256	$3.1412 \cdot 10^{-3}$	$3.1015 \cdot 10^{-3}$	1.26
Protocol-level (Networks timings of an SSH connection)	sntrup761x25519-sha512	0.1565	0.1534	1.98
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Measurement w.r.t. a
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Small performance gains
should accumulate in large
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	mlkem768x25519-sha256	0.1325	0.1316	0.68

Measurement w.r.t. a
single SSH connection.

Small performance gains
should accumulate in large
scale SSH deployments.

IND-CCA → IND-CPA leads
to $\approx 80\%$ reduction in
hash computations.

Our Contributions

How bad an extra hash can be?

By Sasha Frolov and Rafael Misoczki

- Key exchange is a (very) commonly performed operation at Meta
 - **Currently, ~0.05% of CPU cycles in Meta's data centers are spent doing X25519 key exchange**
 - We hope this data point is useful for making cost estimates while defining PQC standards specs
- This means
 - Deploying post-quantum key exchange has a non-negligible capacity cost
 - Apparently innocuous steps can cost hundreds of thousands or even millions of dollars a year
 - e.g. extra hashing steps, like hashing randomness or hashing parts of the transcript, which are being discussed as part of finalizing Kyber specification
 - Even if an extra step does not affect latency, the extra power usage/consumption of shared resources on highly parallel servers still has costs

Feedback? Write to sashafrolov@meta.com or rafam@meta.com.

KEMs in post-quantum SSH.

D-CCA[s]	IND-CPA[s]	Speedup[%]
$164 \cdot 10^{-2}$	$2.7056 \cdot 10^{-2}$	3.93
$853 \cdot 10^{-5}$	$1.3242 \cdot 10^{-5}$	52.46
$290 \cdot 10^{-2}$	$2.9105 \cdot 10^{-2}$	3.91
$412 \cdot 10^{-3}$	$3.1015 \cdot 10^{-3}$	1.26
0.1565	0.1534	1.98
0.1325	0.1316	0.68

IND-CCA → IND-CPA leads to **≈80% reduction** in hash computations.



(*honest reaction)

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*“**TLS**, and not SSH, represents the major workload for Meta.”*

- Reviewer

2

IND-CCA → IND-CPA leads to **≈80% reduction** in hash computations.

Our Contributions

Post-quantum Cryptographic Analysis of SSH

Benjamin Benčina

Royal Holloway, University of London, UK
Email: benjamin.bencina.2022@live.rhul.ac.uk

Varun Maram

SandboxAQ, UK
Email: varun.maram@sandboxaq.com

Benjamin Dowling

King's College London, UK
Email: benjamin.dowling@kcl.ac.uk

Keita Xagawa

Technology Innovation Institute, UAE
Email: keita.xagawa@tii.ae

[S&P '25]

IND-CPA secure ephemeral
KEMs suffice for PQ-SSH!

Our Contributions

Post-quantum Cryptographic Analysis of SSH

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Royal Holloway, University of London, UK
Email: benjamin.bencina.2022@live.rhul.ac.uk

Benjamin Dowling

King's College London, UK
Email: benjamin.dowling@kcl.ac.uk

Varun Maram

SandboxAQ, UK
Email: varun.maram@sandboxaq.com

Keita Xagawa

Technology Innovation Institute, UAE
Email: keita.xagawa@tii.ae

[S&P '25]

IND-CPA secure ephemeral
KEMs suffice for PQ-SSH!

IND-CPA secure ephemeral
KEMs suffice for PQ-TLS too!

On IND-qCCA security in the ROM and its applications

CPA security is sufficient for TLS 1.3

Loïs Huguenin-Dumittan, Serge Vaudenay

EPFL, Lausanne, Switzerland
{lois.huguenin-dumittan, serge.vaudenay}@epfl.ch

[Eurocrypt '22]

CPA-secure KEMs are also sufficient for Post-Quantum TLS 1.3

Biming Zhou^{1,3}, Haodong Jiang², and Yunlei Zhao^{1,3}

¹ College of Computer Science and Technology, Fudan University, Shanghai 200433, China

² Henan Key Laboratory of Network Cryptography Technology, Zhengzhou, 450001, Henan, China

³ State Key Laboratory of Cryptology, Beijing 100878, China
bmzhou22@m.fudan.edu.cn, hdjiang13@163.com, ylzhao@fudan.edu.cn

[Asiacrypt '24]

Discussion

Q: “Should post-quantum SSH/TLS be made efficient, and eco(log|nom)ical?”

Discussion

How bad an extra hash can be?

By Sasha Frolov and Rafael Misoczki

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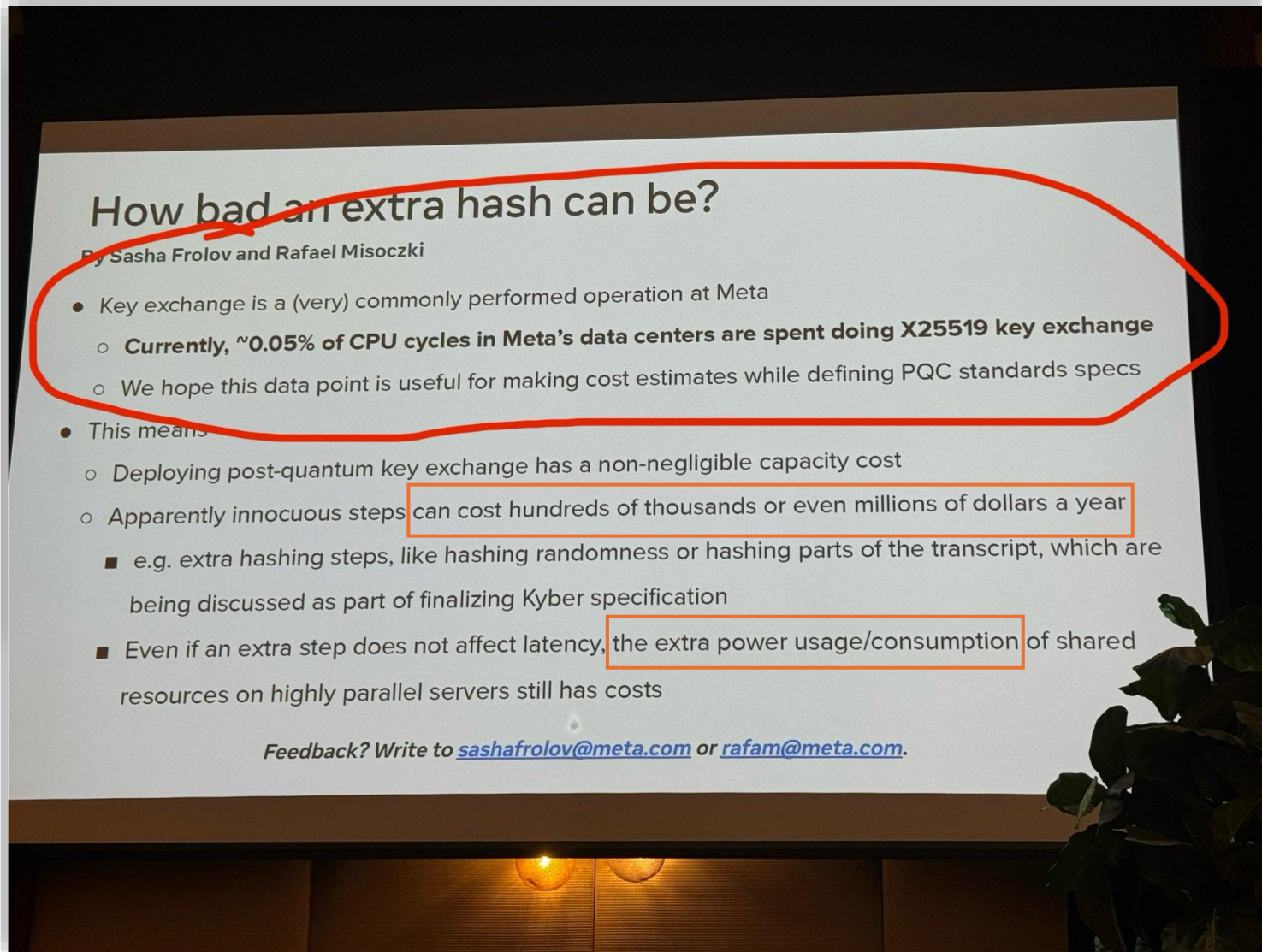
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*“**TLS**, and not SSH, represents the major workload for Meta.”*

- Reviewer

2

Discussion



*“**TLS**, and not SSH, represents the major workload for Meta.”*

- Reviewer

2

*“~~TLS~~, and not ~~SSH~~, **AI** represents the major workload for Meta.”*

- Reviewer

3

Discussion

Stebila, et al.
Internet-Draft

Expires 18 July 2025
ietf-tls-hybrid-design

[Page 15]
January 2025

that the hash function is a dual-PRF.

As noted in Section 2, KEMs used in the manner described in this document MUST explicitly be designed to be secure in the event that the public key is reused, such as achieving IND-CCA2 security or having a transform like the Fujisaki-Okamoto transform applied. ML-KEM has such security properties. However, some other post-quantum KEMs designed to be IND-CPA-secure (i.e., without countermeasures such as the FO transform) are completely insecure under public key reuse; for example, some lattice-based IND-CPA-secure KEMs are vulnerable to attacks that recover the private key after just a few thousand samples [FLUHRER].

- IND-CCA secure KEMs defend against potential **public key-reuse attacks** in faulty TLS/SSH implementations.
- Is there a way to retain **IND-CPA efficiency**, while preventing such **implementation errors** from happening?

Discussion

Stebila, et al.
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Our Contributions

#1: “Post-quantum cryptography in practice*.”



(*Analysis does not cover low-level implementation details.)

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Our Contributions

#1: “Post-quantum cryptography in practice*.”



(*Analysis does not cover **low-level implementation details**.)

- IND-CCA secure KEMs defend against potential **public key-reuse attacks** in faulty TLS/SSH implementations.
- Is there a way to retain **IND-CPA efficiency**, while preventing such **implementation errors** from happening?
- Or, should we formally analyze TLS/SSH in a “faulty/adversarial implementation” (i.e., **kleptographic**) model?
 - We might then have to rely on “stronger-than-IND-CCA” security of KEMs.

Discussion

Efficient IND-CPA secure KEMs
suffice for post-quantum SSH!



Discussion

Efficient IND-CPA secure KEMs suffice for post-quantum SSH!



Cool! Are such IND-CPA secure KEMs FIPS-compliant?



Damien Miller
(OpenSSH Maintainer)

Discussion

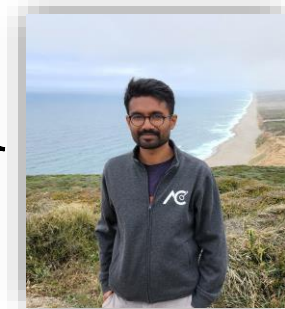
Efficient IND-CPA secure KEMs suffice for post-quantum SSH!



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Discussion

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4.A.3 Security Definition for Ephemeral-Only Encryption/Key-Establishment While chosen ciphertext security is necessary for many existing applications (for example, nominally ephemeral key exchange protocols that allow key caching), it is possible to implement a purely ephemeral key exchange protocol in such a way that only passive security is required from the encryption or KEM primitive.

For these applications, NIST will consider standardizing an encryption or KEM scheme which provides semantic security with respect to chosen plaintext attack. This property is generally denoted *IND-CPA security* in academic literature.

Should NIST have an “on-ramp” process for IND-CPA secure KEMs w.r.t. key exchange protocols?

Post-quantum Cryptographic Analysis of SSH



SANDBOXAQ

Varun Maram
Cybersecurity Group
SandboxAQ



: <https://varun-maram.github.io/>



: varun-maram-pqc

Joint work with Benjamin Benčina, Benjamin Dowling, and Keita Xagawa

